

A NOTE ON CONTINUOUS FUNCTIONS ON METRIC SPACES

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Abstract. Continuous functions on the unit interval are relatively *tame* from the logical and computational point of view. A similar behaviour is exhibited by continuous functions on compact metric spaces *equipped with a countable dense subset*. It is then a natural question what happens if we omit the latter ‘extra data’, i.e., work with ‘unrepresented’ compact metric spaces. In this paper, we study basic third-order statements about continuous functions on such unrepresented compact metric spaces in Kohlenbach’s higher-order Reverse Mathematics. We establish that some (very specific) statements are classified in the (second-order) Big Five of Reverse Mathematics, while most variations/generalisations are not provable from the latter, and much stronger systems. Thus, continuous functions on unrepresented metric spaces are ‘wild’, though ‘more tame’ than (slightly) discontinuous functions on the reals.

§1. Introduction. In a nutshell, we study basic third-order statements about continuous functions on ‘unrepresented’ metric spaces, i.e., the latter come *without* second-order representation, working in Kohlenbach’s higher-order Reverse Mathematics, as introduced in [19] and Section 1.1.2. We establish that certain (very specific) such statements are classified in the second-order Big Five of Reverse Mathematics, while most variations/generalisations are not provable from the latter, and much stronger systems. Thus, we generalise the results in [32] to metric spaces, but restrict ourselves to continuous functions.

We believe these results to be of broad interest as the logic (and even mathematics) community should be aware of the influence representations have on some of the most basic objects, like continuous functions on compact metric spaces, that feature in undergraduate curricula in mathematics and physics.

Moreover, our results also shed new light on Kohlenbach’s *proof mining* program: as stated in [20, Section 17.1] or [21, Section 1], the success of proof mining often crucially depends on *avoiding* the use of separability

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conditions. By the results in this paper, avoiding such conditions seems to be a highly non-trivial affair.

We provide some background and motivation for these results in Section 1.1, including a gentle introduction to higher-order Reverse Mathematics. We formulate necessary definitions and axioms in Section 1.2 and prove our main results in Section 2. Finally, some foundational implications of our results, including related to the coding practise of Reverse Mathematics, are discussed in Section 3.

1.1. Motivation and background.

1.1.1. Introduction. In a nutshell, the topic of this paper is the study of compact metric spaces in higher-order arithmetic; this section provides ample motivation for this study, as well as a detailed overview of our results.

Now, we assume familiarity with the program *Reverse Mathematics*, abbreviated ‘RM’ in the below. An introduction to RM for the mathematician-in-the-street may be found in [44], while [10, 42] are textbooks on RM. We shall work in Kohlenbach’s higher-order RM, introduced in Section 1.1.2. In Section 1.1.3, we provide some motivation for higher-order RM (and this paper), based on the following items.

- (a) *Simplicity*: the coding of continuous functions and other higher-order objects in second-order RM is generally not needed in higher-order RM.
- (b) *Scope*: discontinuous $\mathbb{R} \rightarrow \mathbb{R}$ -functions have been studied by Euler, Abel, Riemann, Fourier, and Dirichlet, i.e., the former are definitely part and parcel of ordinary mathematics. Discontinuous functions are studied directly in higher-order RM; the second-order approach via codes has its problems.
- (c) *Generality*: do the results of RM, like the Big Five phenomenon, depend on the coding practise of RM? Higher-order RM provides a (much needed, in our opinion) negative answer in the case of continuous functions.

Having introduced and motivated higher-order RM in Sections 1.1.2 and 1.1.3, we discuss the results of this paper in some detail in Section 1.1.4. As we will see, the motivation for the study of compact metric spaces in this paper is provided by items (a) and (c) above. In particular, we investigate whether the representation of metric spaces in second-order RM has an influence on the logical properties of third-order theorems about compact metric spaces. The answer turns out to be positive, for all but one very specific choice of definitions.

1.1.2. Higher-order Reverse Mathematics. We provide a gentle introduction to Kohlenbach’s *higher-order* RM, including the base theory RCA_0^ω . Our focus is on intuitive understanding rather than full technical detail.

First of all, the language of RCA_0^ω includes all finite types. In particular, the collection of *all finite types* \mathbf{T} is defined by the two clauses:

- (i) $0 \in \mathbf{T}$ and (ii) If $\sigma, \tau \in \mathbf{T}$ then $(\sigma \rightarrow \tau) \in \mathbf{T}$,

where 0 is the type of natural numbers, and $\sigma \rightarrow \tau$ is the type of mappings from objects of type σ to objects of type τ . The following table provides an overview of the most common objects, their types, and their order.

Symbol	Order	Type	Name
$n \in \mathbb{N}$ or n^0	First	0	Natural number
$X \subset \mathbb{N}$	Second	1	Subset of \mathbb{N}
$f \in 2^{\mathbb{N}}$	Second	1	Element of Cantor space $2^{\mathbb{N}}$
$f \in \mathbb{N}^{\mathbb{N}}$ or f^1	Second	1	Element of Baire space $\mathbb{N}^{\mathbb{N}}$
$Y : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}$ or Y^2	Third	2	Mapping of Baire space to \mathbb{N}
$Y : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}$ or $Y^{1 \rightarrow 1}$	Third	2	Mapping of Baire space to Baire space

One often identifies elements of Cantor space $2^{\mathbb{N}}$ and subsets of \mathbb{N} , as the former can be viewed as characteristic functions for the latter. Similarly, a subset $X \subset \mathbb{N}^{\mathbb{N}}$ is given by the associated characteristic function $\mathbb{1}_X : \mathbb{N}^{\mathbb{N}} \rightarrow \{0, 1\}$. In this paper, we shall mostly restrict ourselves to third-order objects.

Secondly, a basic axiom of mathematics is that functions *map equal inputs to equal outputs*. The *axiom of function extensionality* guarantees this behaviour and is included in RCA_0^ω . As an example, we write $f =_1 g$ in case $(\forall n \in \mathbb{N})(f(n) = g(n))$ and say that ‘ f and g are equal elements of Baire space’. Any third-order $Y^{1 \rightarrow 1}$ thus satisfies the following instance of the axiom of function extensionality:

$$(\forall f, g \in \mathbb{N}^{\mathbb{N}})(f =_1 g \rightarrow Y(f) =_1 Y(g)). \tag{E_{1 \rightarrow 1}}$$

Now, the real number field \mathbb{R} is central to analysis and other parts of mathematics. The real numbers are defined in RCA_0^ω in *exactly the same way* as in RCA_0 , namely as fast-converging Cauchy sequences. Hence, the formulas ‘ $x \in \mathbb{R}$ ’, ‘ $x <_{\mathbb{R}} y$ ’, and ‘ $x =_{\mathbb{R}} y$ ’ have their usual well-known meaning. To define functions on the reals, we let an ‘ $\mathbb{R} \rightarrow \mathbb{R}$ -function’ be any $Y^{1 \rightarrow 1}$ that satisfies

$$(\forall x, y \in \mathbb{R})(x =_{\mathbb{R}} y \rightarrow Y(x) =_{\mathbb{R}} Y(y)), \tag{E_{\mathbb{R} \rightarrow \mathbb{R}}}$$

which is the axiom of function extensionality relative to the (defined) equality ‘ $=_{\mathbb{R}}$ ’. We stress that all symbols pertaining to the real numbers in $(E_{\mathbb{R} \rightarrow \mathbb{R}})$ have their usual second-order meaning.

Thirdly, RCA_0 is a system of ‘computable mathematics’ that includes comprehension for Δ_1^0 -formulas and induction for Σ_1^0 -formulas. The former allows one to build algorithms, e.g., via primitive recursion, while the latter certifies their correctness. The base theory RCA_0^ω includes axioms that prove

Δ_1^0 -comprehension and Σ_1^0 -induction, as expected. Moreover, RCA_0^ω includes the defining axiom of the recursor constant \mathbf{R}_0 , namely that for $m \in \mathbb{N}$ and $f \in \mathbb{N}^{\mathbb{N}}$, we have

$$\mathbf{R}_0(f, m, 0) := m \text{ and } \mathbf{R}_0(f, m, n + 1) := f(n, \mathbf{R}_0(f, m, n)), \tag{1.1}$$

which defines primitive recursion with second-order functions. We hasten to add that higher-order parameters are allowed; as an example, we could use in (1.1) the function $f \in \mathbb{N}^{\mathbb{N}}$ defined as $f(n) := Y(q_n)$ for any $Y : \mathbb{R} \rightarrow \mathbb{N}$ and where $(q_n)_{n \in \mathbb{N}}$ lists all rational numbers. The previous example also illustrates—to our mind—the need for *lambda calculus* notation, where we would simply define $f \in \mathbb{N}^{\mathbb{N}}$ as $\lambda n^0. Y(q_n)$, underscoring that n is the (only) variable and Y a parameter. The system RCA_0^ω includes the defining axioms for λ -abstraction via combinators.

Fourth, second-order RM includes many results about codes for continuous functions, and we would like to ‘upgrade’ these results to third-order functions that satisfy the usual ‘epsilon-delta’ definition of continuity. To this end, RCA_0^ω includes the following fragment of the Axiom of Choice, provable in ZF:

$$(\forall f \in \mathbb{N}^{\mathbb{N}})(\exists n^0 \in \mathbb{N})A(f, n) \rightarrow (\exists Y : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N})(\forall f \in \mathbb{N}^{\mathbb{N}})A(f, Y(f)), \tag{QF-AC^{1,0}}$$

for any quantifier-free formula A . Now, the following formula:

$$\Phi \subset \mathbb{N} \text{ is a total code for a continuous } \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}^{\mathbb{N}}\text{-function} \tag{1.2}$$

has exactly the form as in the antecedent of QF-AC^{0,1}. Applying the latter to (1.2), we readily obtain a function $Z^{1 \rightarrow 1}$ such that $Z(f)$ equals the value of Φ at any $f \in \mathbb{N}^{\mathbb{N}}$. A similar argument goes through for codes of continuous $\mathbb{R} \rightarrow \mathbb{R}$ -functions.

Finally, RCA_0^ω and RCA_0 prove the same second-order sentences, a fact that is proved via the so-called ECF-translation (see [19, Section 2]). In a nutshell, the latter replaces third- and higher-order objects by second-order codes for continuous functions, a concept not alien to second-order RM. We will discuss the coding of continuous functions in more detail in Section 1.1.3.

1.1.3. The coding practise of Reverse Mathematics. We discuss the coding practise of RM, which will be seen to provide motivation for Kohlenbach’s higher-order RM.

First of all, second-order RM makes use of the rather frugal language of second-order arithmetic, which only includes variables for natural numbers $n \in \mathbb{N}$ and sets of natural numbers $X \subset \mathbb{N}$. As a result, higher-order objects like functions on the reals and metric spaces, are unavailable and need to be ‘represented’ or ‘coded’ by second-order objects. This coding practise

can complicate basic definitions: the reader need only compare the one-line ‘epsilon-delta’ definition of continuity to the second-order definition in [42, II.6.1]; the latter takes the better half of a page. Similar complications arise for metric spaces, where the reader can compare Definition 1.3 to [42, I.8.6]. Hence, a framework that includes third- and higher-order objects provides a *simpler* approach.

Secondly, discontinuous functions have been studied in mathematics long before the advent of set theory, by ‘big name’ mathematicians like Euler, Dirichlet, Riemann, Abel, and Fourier, as discussed in [37, Section 5.2]. Hence, basic properties of discontinuous functions are part of ordinary mathematics and should be studied in RM. Higher-order RM provides a natural framework for the study of discontinuous functions, as explored in detail in [19, 32]. By contrast, the study of discontinuous functions via codes is problematic¹, as discussed in detail in [37, Section 6.2.2].

Thirdly, a central objective of mathematical logic is the classification of mathematical statements in hierarchies according to their logical strength. An example due to Simpson is the *Gödel hierarchy* from [43]. The goal of RM, namely finding the minimal axioms that prove a theorem of ordinary mathematics, fits squarely into this objective. Ideally, the place of a given statement in the hierarchy at hand does not depend greatly on the representation used. In particular, RM seeks to analyse theorems of ordinary mathematics ‘as they stand’. The following quotes from [42, pages 32 and 137] illustrate this claim.

The typical constructivist response to a nonconstructive mathematical theorem is to modify the theorem by adding hypotheses or “extra data”. In contrast, our approach in this book is to analyze the provability of mathematical theorems as they stand [...]

[...] we seek to draw out the set existence assumptions which are implicit in the ordinary mathematical theorems *as they stand*.

Essentially the same claim may be found in [10, Section 10.5.2] and in many parts of the RM literature. The main point is always the same: RM ideally studies mathematics ‘as is’ without any logical additions.

Since the coding of continuous functions is conspicuously absent from mainstream mathematics, it is then a natural question whether the aforementioned coding has an influence on the classification of theorems in

¹In a nutshell, RCA_0 and RCA_0^ω are consistent with *Church’s thesis* CT, i.e., the statement *all real numbers are computable*. In fact, the recursive sets form the minimal ω -model of RCA_0 by [42, I.7.5]. Now, $RCA_0^\omega + CT$ proves that all functions on the reals are continuous, while in $RCA_0 + CT$, there are plenty of codes for discontinuous functions, e.g., for the Heaviside function. Thus, RCA_0^ω exhibits a connection between second- and third-order objects and theorems, which is seemingly obliterated by coding functions.

RM. The following theorem implies that at least the Big Five phenomenon of RM does not really depend on the coding of continuous functions on various spaces.

THEOREM 1.1. *The system RCA_0^ω proves the following for $\mathbb{X} = \mathbb{N}^{\mathbb{N}}$ or $\mathbb{X} = \mathbb{R}$.*

Let Φ be a code for a continuous $\mathbb{X} \rightarrow \mathbb{X}$ -function. There is a third-order $F : \mathbb{X} \rightarrow \mathbb{X}$ such that $F(x)$ equals the value of Φ at x for any $x \in \mathbb{X}$.

The system $\text{RCA}_0^\omega + \text{WKL}_0$ proves the following for $\mathbb{X} = 2^{\mathbb{N}}$ or $\mathbb{X} = [0, 1]$.

Let the third-order function $F : \mathbb{X} \rightarrow \mathbb{X}$ be continuous on \mathbb{X} . Then there is a code Φ such that $F(x)$ equals the value of Φ at any $x \in \mathbb{X}$.

PROOF. See [18, Section 4] and [32, Section 2] for proofs. ◻

As a corollary, we observe that over RCA_0^ω , the *second-order* axiom WKL_0 is equivalent to *third-order* theorems like for any third-order $F : [0, 1] \rightarrow \mathbb{R}$, *continuity implies boundedness*. One could argue that continuous functions are ‘really’ second-order, but the reader should keep the previous sentence in mind nonetheless.

1.1.4. Metric spaces and higher-order Reverse Mathematics. Having introduced and motivated higher-order RM in the previous two sections, we can now discuss and motivate the results in this paper in some detail.

First of all, an immediate corollary of Theorem 1.1 is that over RCA_0^ω , the *second-order* axiom WKL_0 is equivalent to many basic *third-order* theorems from real analysis about continuous functions. Motivated by this observation, Dag Normann and the author show in [32, 40, 41] that *many third-order* theorems from real analysis about (possibly) discontinuous functions on the reals, are equivalent to the *second-order* Big Five, over RCA_0^ω . Moreover, *slight* variations/generalisations of the function class at hand yield third-order theorems that are not provable from the Big Five and the same for much stronger systems like $Z_2^\omega + \text{QF-AC}^{0,1}$ from Section 1.2. It is then a natural question whether a similar phenomenon can be found in other parts of mathematics and RM.

Secondly, in this paper, we study a different kind of generalisation than the one in [32]: rather than going beyond the continuous functions, we study properties of the latter on compact metric spaces. Now, the study of the latter in second-order RM proceeds via codes: a *complete separable metric space* is represented via a countable and dense subset, as can be gleaned from [42, II.5.1] or [5]. By contrast, we use the standard textbook definition of metric space as in Definition 1.3 without any additional data except that we are dealing with sets of reals. This study is not just *spielerei* as *avoiding* separability is, e.g., important in proof mining, as follows.

[...] it is crucial to exploit the fact that the proof to be analyzed does not use any separability assumption on the underlying spaces [...]. [21, Section 1]

It will turn out that for [the aforementioned uniformity conditions] to hold we—in particular—must not use any separability assumptions on the spaces. [20, page 377]

Thirdly, in light of the previous two paragraphs, it is then a natural question whether basic properties of compact metric spaces *without separability conditions* are provable from second-order (comprehension) axioms or not. Theorem 2.2 provides a (rather) negative answer: well-known theorems due to Ascoli, Arzelà, Dini, Heine, and Pincherle, formulated for metric spaces, are not provable in Z_2^ω , a conservative extension of Z_2 introduced in Section 1.2. We only study metric spaces (M, d) where M is a subset of the reals or Baire space, i.e., the metric $d : M^2 \rightarrow \mathbb{R}$ is just a third-order mapping. By contrast, some (very specific) basic properties of metric spaces are provable from the Big Five and related systems by Theorem 2.3.

Fourth, the negative results in this paper are established using the uncountability of \mathbb{R} as formalised by the following principles (see Section 1.2 for details).

- $\text{NIN}_{[0,1]}$: there is no injection from $[0, 1]$ to \mathbb{N} .
- $\text{NBI}_{[0,1]}$: there is no bijection from $[0, 1]$ to \mathbb{N} .

In particular, these principles are not provable in relatively strong systems, like Z_2^ω from Section 1.2. In Section 2.1, we identify a long and robust list of theorems that imply $\text{NBI}_{[0,1]}$ or $\text{NIN}_{[0,1]}$. We have shown in [31, 32, 36] that many third-order theorems imply $\text{NIN}_{[0,1]}$ while we only know few theorems that only imply $\text{NBI}_{[0,1]}$. As will become clear in Section 2.2, metric spaces provide (many) elegant examples of the latter. We also refine our results in Section 2.2, including connections to the RM of weak König's lemma and the Jordan decomposition theorem.

In conclusion, we show that many basic (third-order) properties of continuous functions on metric spaces cannot be proved from second-order (comprehension) axioms when we omit the second-order representation of these spaces. A central principle is the uncountability of the reals as formalised by $\text{NBI}_{[0,1]}$ introduced above. These results carry foundational implications, as discussed in Section 3.

1.2. Preliminaries and definitions. We introduce some definitions, like the notion of open set or metric space in RM, and axioms that cannot be found in [19]. We emphasise that we only study metric spaces (M, d) where M is a subset of $\mathbb{N}^{\mathbb{N}}$ or \mathbb{R} , modulo the coding of finite sequences² of reals. Thus,

²We use w^{1^*} to denote finite sequences of elements of $\mathbb{N}^{\mathbb{N}}$ and $|w|$ as the length of w^{1^*} .

everything can be formalised in the language of third-order arithmetic, i.e., we do not really go much beyond analysis on the reals.

Zerth of all, we need to define the notion of (open) set. Now, open sets are represented in second-order RM by *countable unions of basic open balls*, namely as in [42, II.5.6]. In light of [42, II.7.1], (codes for) *continuous functions* provide an equivalent representation over RCA_0 . In particular, the latter second-order representation is exactly the following definition restricted to (codes for) continuous functions, as can be found in [42, II.6.1].

DEFINITION 1.2.

- A set $U \subset \mathbb{R}$ (and its complement U^c) is given by $h_U : \mathbb{R} \rightarrow [0, 1]$ where we say ‘ $x \in U$ ’ if and only if $h_U(x) > 0$.
- A set $U \subset \mathbb{R}$ is *open* if $y \in U$ implies $(\exists N \in \mathbb{N})(\forall z \in B(y, \frac{1}{2N}))(z \in U)$. A set is closed if the complement is open.
- A set $U \subset \mathbb{R}$ is *finite* if there is $N \in \mathbb{N}$ such that for any finite sequence (x_0, \dots, x_N) , there is $i \leq N$ with $x_i \notin U$. We sometimes write ‘ $|U| \leq N$ ’.

Now, codes for continuous functions denote third-order functions in RCA_0^ω by [32, Section 2], i.e., Definition 1.2 thus includes the second-order definition of open set. To be absolutely clear, combining [32, Theorem 2.2] and [42, II.7.1], RCA_0^ω proves

[a second-order code U for an open set] represents an open set as in Definition 1.2.

Assuming Kleene’s quantifier (\exists^2) defined below, Definition 1.2 is equivalent to the existence of a characteristic function for U ; the latter definition is used in, e.g., [27, 33]. The interested reader can verify that over RCA_0^ω , a set U as in Definition 1.2 is open if and only if h_U is lower semi-continuous.

First of all, we shall study metric spaces (M, d) as in Definition 1.3, where M comes with its own equivalence relation ‘ $=_M$ ’ and the metric d satisfies the axiom of extensionality on M as follows:

$$(\forall x, y, v, w \in M)([x =_M y \wedge v =_M w] \rightarrow d(x, v) =_{\mathbb{R}} d(y, w)).$$

Similarly, we use $F : M \rightarrow \mathbb{R}$ to denote functions from M to \mathbb{R} ; the latter satisfy

$$(\forall x, y \in M)(x =_M y \rightarrow F(x) =_{\mathbb{R}} F(y)), \tag{E_M}$$

i.e., the axiom of function extensionality relative to M .

DEFINITION 1.3. A functional $d : M^2 \rightarrow \mathbb{R}$ is a *metric on M* if it satisfies the following properties for $x, y, z \in M$:

- (a) $d(x, y) =_{\mathbb{R}} 0 \leftrightarrow x =_M y$,
- (b) $0 \leq_{\mathbb{R}} d(x, y) =_{\mathbb{R}} d(y, x)$,
- (c) $d(x, y) \leq_{\mathbb{R}} d(x, z) + d(z, y)$.

We use standard notation like $B_d^M(x, r)$ to denote $\{y \in M : d(x, y) < r\}$.

To be absolutely clear, quantifying over M amounts to quantifying over $\mathbb{N}^{\mathbb{N}}$ or \mathbb{R} , perhaps modulo coding, i.e., the previous definition can be made in third-order arithmetic for the intents and purposes of this paper. The definitions of ‘open set in a metric space’ and related constructs are now clear *mutatis mutandis*.

Secondly, the following definitions are now standard, where we note that the first item is called ‘Heine–Borel compact’ in, e.g., [3, 5]. Moreover, coded complete separable metric spaces as in [42, I.8.2] are only *weakly complete* over RCA_0 .

DEFINITION 1.4 (Compactness and around). For a metric space (M, d) , we say that:

- (M, d) is *countably-compact* if for any $(a_n)_{n \in \mathbb{N}}$ in M and sequence of rationals $(r_n)_{n \in \mathbb{N}}$ such that we have $M \subset \cup_{n \in \mathbb{N}} B_d^M(a_n, r_n)$, there is $m \in \mathbb{N}$ such that $M \subset \cup_{n \leq m} B_d^M(a_n, r_n)$.
- (M, d) is *strongly countably-compact* if for any sequence $(O_n)_{n \in \mathbb{N}}$ of open sets in M such that $M \subset \cup_{n \in \mathbb{N}} O_n$, there is $m \in \mathbb{N}$ such that $M \subset \cup_{n \leq m} O_n$.
- (M, d) is *compact* in case for any $\Psi : M \rightarrow \mathbb{R}^+$, there are $x_0, \dots, x_k \in M$ such that $\cup_{i \leq k} B_d^M(x_i, \Psi(x_i))$ covers M .
- (M, d) is *sequentially compact* if any sequence has a convergent subsequence.
- (M, d) is *limit point compact* if any infinite set in M has a limit point.
- (M, d) is *complete* in case every *Cauchy*³ sequence converges.
- (M, d) is *weakly complete* if every *effectively*³ Cauchy sequence converges.
- (M, d) is *totally bounded* if for all $k \in \mathbb{N}$, there are $w_0, \dots, w_m \in M$ such that $\cup_{i \leq m} B_d^M(w_i, \frac{1}{2^k})$ covers M .
- (M, d) is *effectively totally bounded* if there is a sequence of finite sequences $(w_n)_{n \in \mathbb{N}}$ in M such that for all $k \in \mathbb{N}$ and $x \in M$, there is $i < |w_k|$ such that $x \in B_d^M(w(i), \frac{1}{2^k})$.
- a set $C \subset M$ is *sequentially closed* if for any sequence $(w_n)_{n \in \mathbb{N}}$ in C converging to $w \in M$, we have $w \in C$.
- (M, d) has the *Cantor intersection property* if any sequence of nonempty closed sets with $M \supseteq C_0 \supseteq \dots \supseteq C_n \supseteq C_{n+1}$, has a nonempty intersection,

³A sequence $(w_n)_{n \in \mathbb{N}}$ in (M, d) is *Cauchy* if $(\forall k \in \mathbb{N})(\exists N \in \mathbb{N})(\forall m, n \geq N)(d(w_m, w_n) < \frac{1}{2^k})$. A sequence is *effectively Cauchy* if there is $g \in \mathbb{N}^{\mathbb{N}}$ such that $g(k) = N$ in the previous formula.

- (M, d) has the *sequential Cantor intersection property* if the sets in the previous item are sequentially closed.
- (M, d) is *separable* if there is a sequence $(x_n)_{n \in \mathbb{N}}$ in M such that $(\forall x \in M, k \in \mathbb{N})(\exists n \in \mathbb{N})(d(x, x_n) < \frac{1}{2^k})$.

Thirdly, full second-order arithmetic Z_2 is the ‘upper limit’ of second-order RM. The systems Z_2^ω and Z_2^Ω are conservative extensions of Z_2 by [15, Corollary 2.6]. The system Z_2^Ω is RCA_0^ω plus Kleene’s quantifier (\exists^3) (see, e.g., [15, 32]), while Z_2^ω is RCA_0^ω plus (S_k^2) for every $k \geq 1$; the latter axiom states the existence of a functional S_k^2 deciding Π_k^1 -formulas in Kleene normal form. The system $\Pi_1^1\text{-CA}_0^\omega \equiv RCA_0^\omega + (S_1^2)$ is a Π_3^1 -conservative extension of $\Pi_1^1\text{-CA}_0$ [35], where S_1^2 is also called the Suslin functional. We also write ACA_0^ω for $RCA_0^\omega + (\exists^2)$ where the latter is as follows:

$$(\exists E : \mathbb{N}^{\mathbb{N}} \rightarrow \{0, 1\})(\forall f \in \mathbb{N}^{\mathbb{N}})[(\exists n \in \mathbb{N})(f(n) = 0) \leftrightarrow E(f) = 0]. \quad (\exists^2)$$

Over RCA_0^ω , (\exists^2) is equivalent to the existence of Feferman’s μ (see [19, Proposition 3.9]), defined as follows for all $f \in \mathbb{N}^{\mathbb{N}}$:

$$\mu(f) := \begin{cases} n, & \text{if } n \text{ is the least natural such that } f(n) = 0, \\ 0, & \text{if } f(n) > 0 \text{ for all } n \in \mathbb{N}. \end{cases}$$

Fourth, the uncountability of the reals, formulated as follows, is studied in [31].

- $\text{NIN}_{[0,1]}$: there is no $Y : [0, 1] \rightarrow \mathbb{N}$ that is injective⁴.
- $\text{NBI}_{[0,1]}$: there is no $Y : [0, 1] \rightarrow \mathbb{N}$ that is both injective and surjective⁵.

It is shown in [30, 31] that Z_2^ω cannot prove $\text{NBI}_{[0,1]}$ and that $Z_2^\omega + \text{QF-AC}^{0.1}$ cannot prove $\text{NIN}_{[0,1]}$, where the latter is countable choice⁶ for quantifier-free formulas. Moreover, many third-order theorems imply $\text{NIN}_{[0,1]}$, as also established in [31]. By contrast, that \mathbb{R} cannot be enumerated is formalised by Theorem 1.5.

THEOREM 1.5. *For any sequence of distinct real numbers $(x_n)_{n \in \mathbb{N}}$ and any interval $[a, b]$, there is $y \in [a, b]$ such that y is different from x_n for all $n \in \mathbb{N}$.*

The previous theorem is rather tame, especially compared to $\text{NIN}_{[0,1]}$. Indeed, [13] includes an efficient computer program that computes the number y from Theorem 1.5 in terms of the other data; a proof of Theorem 1.5 in RCA_0 can be found in [42, II.4.9], while a proof in Bishop’s *Constructive Analysis* is found in [2, page 25].

⁴A function $f : X \rightarrow Y$ is injective if different $x, x' \in X$ yield different $f(x), f(x') \in Y$.

⁵A function $f : X \rightarrow Y$ is surjective if for every $y \in Y$, there is $x \in X$ with $f(x) = y$.

⁶To be absolutely clear, $\text{QF-AC}^{0.1}$ states that for every $Y^2, (\forall n \in \mathbb{N})(\exists f \in \mathbb{N}^{\mathbb{N}})(Y(f, n) = 0)$ implies $(\exists \Phi^{0 \rightarrow 1})(\forall n \in \mathbb{N})(Y(\Phi(n), n) = 0)$.

Finally, the following remark discusses an interesting aspect of (\exists^2) and $\text{NIN}_{[0,1]}$.

REMARK 1.6 (On excluded middle). Despite the grand stories told in mathematics and logic about Hilbert and the law of excluded middle, the ‘full’ use of the latter law in RM is almost somewhat of a novelty. To be more precise, the law of excluded middle as in $(\exists^2) \vee \neg(\exists^2)$ is extremely useful, namely as follows: suppose we are proving $T \rightarrow \text{NIN}_{[0,1]}$ over RCA_0^ω . Now, in case $\neg(\exists^2)$, all functions on \mathbb{R} (and $\mathbb{N}^{\mathbb{N}}$) are continuous by [19, Proposition 3.12]. Clearly, any continuous $Y : [0, 1] \rightarrow \mathbb{N}$ is not injective, i.e., $\text{NIN}_{[0,1]}$ follows in the case that $\neg(\exists^2)$. Hence, what remains is to establish $T \rightarrow \text{NIN}_{[0,1]}$ in case we have (\exists^2) . However, the latter axiom, e.g., implies ACA_0 (and sequential compactness) and can uniformly convert reals to their binary representations, which can simplify the remainder of the proof.

Here, $\text{NIN}_{[0,1]}$ is just one example and there are many more, all pointing to a more general phenomenon: while invoking $(\exists^2) \vee \neg(\exists^2)$ may be non-constructive, it does lead to a short proof via case distinction: in case (\exists^2) , one has access to a stronger system while in case $\neg(\exists^2)$, the theorem at hand is a triviality (like for $\text{NIN}_{[0,1]}$ in the previous paragraph), or at least has a well-known second-order proof. We can also work over $\text{RCA}_0^\omega + \text{WKL}_0$, noting that the latter establishes that continuous functions on $[0, 1]$ or $2^{\mathbb{N}}$ have RM-codes (see [32, Section 2] and [18, Section 4]).

§2. Analysis on unrepresented metric spaces. We show that some (very specific) properties of continuous functions on compact metric spaces are classified in the (second-order) Big Five systems of Reverse Mathematics (Section 2.2), while most variations/generalisations are not provable from the latter, and much stronger systems (Section 2.1). The negative results are (mostly) established by deriving $\text{NBI}_{[0,1]}$ (Theorem 2.2), which is not provable in Z_2^ω . We also show that $\text{NIN}_{[0,1]}$ does not follow in most cases (Theorem 2.4).

2.1. Obtaining the uncountability of the reals. In this section, we show that basic properties of continuous functions on compact metric spaces, like Heine’s theorem in item (b), imply the uncountability of the reals as in $\text{NBI}_{[0,1]}$. These basic properties are therefore not provable in Z_2^ω .

First of all, fragments of the induction axiom are sometimes used in an essential way in second-order RM (see, e.g., [22]). The equivalence between induction and bounded comprehension is also well-known in second-order RM [42, X.4.4]. We seem to need a little bit of the induction axiom as follows.

PRINCIPLE 2.1 IND_1 . *Let Y^2 satisfy $(\forall n \in \mathbb{N})(\exists! f \in 2^{\mathbb{N}})[Y(n, f) = 0]$. Then $(\forall n \in \mathbb{N})(\exists w^{1^*})[|w| = n \wedge (\forall i < n)(Y(i, w(i)) = 0)]$.*

Note that IND_1 is a special case of the axiom of finite choice, and is valid in all models considered in [23–29, 31], i.e., $Z_2^\omega + IND_1$ cannot prove $NBI_{[0,1]}$. We have (first) used IND_1 in the RM of the Jordan decomposition theorem in [30].

Secondly, the items in Theorem 2.2 are essentially those in [5, Theorem 4.1] or [42, IV.2.2], but without codes. Equivalences of certain (coded) definitions of compactness are studied in second-order RM in, e.g., [3, 4].

THEOREM 2.2 ($RCA_0^\omega + IND_1$). *The principle $NBI_{[0,1]}$ follows from any of the items (a)–(s) where (M, d) is a metric space with $M \subset \mathbb{R}$.*

- (a) *For countably-compact (M, d) and sequentially continuous $F : M \rightarrow \mathbb{R}$, F is bounded on M .*
- (b) *Item (a) with ‘bounded’ replaced by ‘uniformly continuous’.*
- (c) *Item (a) with ‘bounded’ replaced by ‘has a supremum’.*
- (d) *Item (a) with ‘bounded’ replaced by ‘attains a maximum’.*
- (e) *A countably-compact (M, d) has the sequential Cantor intersection property.*
- (f) *A countably-compact metric space (M, d) is separable.*

The previous items still imply $NBI_{[0,1]}$ if we replace ‘countably-compact’ by ‘compact’ or ‘(weakly) complete and totally bounded’ or ‘strongly countably-compact’.

- (h) *For sequentially compact (M, d) , any continuous $F : M \rightarrow \mathbb{R}$ is bounded.*
- (i) *Item (h) with ‘bounded’ replaced by ‘uniformly continuous’.*
- (j) *Item (h) with ‘bounded’ replaced by ‘has a supremum’.*
- (k) *Item (h) with ‘bounded’ replaced by ‘attains a maximum’.*
- (l) *Items (h)–(k) assuming a modulus of continuity.*
- (m) *Dini’s theorem ([1, 8, 9]). Let (M, d) be sequentially compact and let $F_n : (M \times \mathbb{N}) \rightarrow \mathbb{R}$ be a monotone sequence of continuous functions converging to continuous $F : M \rightarrow \mathbb{R}$. Then the convergence is uniform.*
- (n) *On a sequentially compact metric space (M, d) , equicontinuity implies uniform equicontinuity.*
- (o) *(Pincherle [34, page 67]). For sequentially compact (M, d) and continuous $F : M \rightarrow \mathbb{R}^+$, we have $(\exists k \in \mathbb{N})(\forall w \in M)(F(w) > \frac{1}{2^k})$.*
- (p) *(Ascoli-Arzelà, [42, III.2]). For sequentially compact (M, d) , a uniformly bounded and equicontinuous sequence of functions on M has a uniformly convergent sub-sequence.*
- (q) *Any sequentially compact (M, d) is strongly countably-compact.*
- (r) *Any sequentially compact (M, d) is separable.*
- (s) *Any sequentially compact (M, d) has the seq. Cantor intersection property.*
- (t) *A sequentially compact metric space (M, d) is limit point compact.*

Items (h)–(l) are provable in Z_2^Ω (via the textbook proof).

PROOF. First of all, by Remark 1.6, we may assume (\exists^2) as $\text{NBI}_{[0,1]}$ is trivial in case $\neg(\exists^2)$. Now suppose $Y : [0, 1] \rightarrow \mathbb{N}$ is a bijection, i.e., injective and surjective. Define M as the union of the new symbol $\{0_M\}$ and the set $N := \{w^{1^*} : (\forall i < |w|)(Y(w(i)) = i)\}$. We define ‘ $=_M$ ’ as $0_M =_M 0_M$, $u \neq_M 0_M$ for $u \in N$, and $w =_M v$ if $w =_{1^*} v$ and $w, v \in N$. The metric $d : M^2 \rightarrow \mathbb{R}$ is defined as $d(0_M, 0_M) =_{\mathbb{R}} 0$, $d(0_M, u) = d(u, 0_M) = \frac{1}{2^{|u|}}$ for $u \in N$ and $d(w, v) = |\frac{1}{2^{|v|}} - \frac{1}{2^{|w|}}|$ for $w, v \in N$. Since Y is an injection, we have $d(v, w) =_{\mathbb{R}} 0 \leftrightarrow v =_M w$. The other properties of a metric space from Definition 1.3 follow by definition (and the triangle equality of the absolute value on the reals).

Secondly, to show that (M, d) is countably-compact, fix a sequence $(a_n)_{n \in \mathbb{N}}$ in M and a sequence of rationals $(r_n)_{n \in \mathbb{N}}$ such that we have $M \subset \cup_{n \in \mathbb{N}} B_d^M(a_n, r_n)$. Suppose $0_M \in B_d^M(a_{n_0}, r_{n_0})$ for $a_{n_0} \neq_M 0_M$, i.e., $\frac{1}{2^{|a_{n_0}|}} = d(0_M, a_{n_0}) < r_{n_0}$. Then $|\frac{1}{2^{|y|}} - \frac{1}{2^{|a_{n_0}|}}| = d(y, a_{n_0}) < r_{n_0}$ holds for all $y \in N$ such that $|y| > |a_{n_0}|$. Now use IND_1 to enumerate the (finitely many) reals $z \in M$ with $|z| < |a_{n_0}|$. In this way, there exists a finite sub-covering of $\cup_{n \in \mathbb{N}} B_d^M(a_n, r_n)$ of at most $|a_{n_0}| + 1$ elements. The proof is analogous (and easier) in case $a_{n_0} =_M 0_M$. Thus, (M, d) is a countably-compact metric space.

Thirdly, define the function $F : M \rightarrow \mathbb{R}$ as follows: $F(0_M) := 0$ and $F(w) := |w|$ for any $w \in N$. Clearly, if the sequence $(w_n)_{n \in \mathbb{N}}$ in M converges to 0_M , either it is eventually constant 0_M or lists all reals in $[0, 1]$. The latter case is impossible by Theorem 1.5. Hence, F is sequentially continuous at 0_M , but not continuous at 0_M . To show that F is (sequentially) continuous at $w \neq 0_M$, consider the formula $|\frac{1}{2^{|w|}} - \frac{1}{2^{|v|}}| = d(v, w) < \frac{1}{2^N}$; the latter is false for $N \geq |w| + 2$ and any $v \neq_M 0_M$. Thus, the following formula is (vacuously) true:

$$(\forall k \in \mathbb{N})(\exists N \in \mathbb{N})(\forall v \in B_d^M(w, \frac{1}{2^N}))(|F(w) - F(v)| < \frac{1}{2^k}).$$

i.e., F is continuous at $w \neq_M 0_M$, with a (kind of) modulus of continuity given. Applying item (a) (or item (c) and (d)), we obtain a contradiction as F is clearly unbounded on M . This contradiction yields $\text{NBI}_{[0,1]}$ and the same for item (b) as F is not (uniformly) continuous.

Fourth, to obtain $\text{NBI}_{[0,1]}$ from item (e), suppose again the former is false and $Y : [0, 1] \rightarrow \mathbb{R}$ and (M, d) are as above. Define $C_n := \{x \in N : |x| > n + 1\}$ and note that this set is non-empty (as Y is a surjection) but satisfies $\cap_n C_n = \emptyset$. Item (e) now yields a contradiction if we can show that C_n is sequentially closed. To the latter end, let $(w_k)_{k \in \mathbb{N}}$ be a sequence in C_n with limit $w \in M$. In case $w =_M 0_M$, we make the same observation as in the third paragraph: either the sequence $(w_k)_{k \in \mathbb{N}}$ is eventually constant 0_M or enumerates the reals in $[0, 1]$. Both are impossible, i.e., this case does not

occur. In case $w \neq 0_M$, we have

$$(\forall k \in \mathbb{N})(\exists N \in \mathbb{N})(\forall n \geq N)(|\frac{1}{2^{|w|}} - \frac{1}{2^{|w_n|}}| = d(w, w_n) < \frac{1}{2^k}),$$

which is only possible if $(w_n)_{n \in \mathbb{N}}$ is eventually constant w . In this case of course, $w \in C_n$, i.e., C_n is sequentially closed, and $(e) \rightarrow \text{NBI}_{[0,1]}$ follows. Regarding item (f), suppose (M, d) is separable, i.e., there is a sequence $(w_n)_{n \in \mathbb{N}}$ such that

$$(\forall w \in M, k \in \mathbb{N})(\exists n \in \mathbb{N})(|\frac{1}{2^{|w|}} - \frac{1}{2^{|w_n|}}| = d(w, w_n) < \frac{1}{2^k}). \tag{2.1}$$

As in the above, for $w \neq_M 0_M$ and $k_0 = |w| + 2$, the formula $d(w, w_n) < \frac{1}{2^{k_0}}$ is false for any $n \in \mathbb{N}$, i.e., we also obtain a contradiction in this case, yielding $\text{NBI}_{[0,1]}$.

Fifth, for the sentences between items (f) and (h), (M, d) is also complete and (strongly countably) compact, which is proved in (exactly) the same way as in the second paragraph: any ball around 0_M covers ‘most’ of M ; to show that (M, d) is complete, let $(w_n)_{n \in \mathbb{N}}$ be a Cauchy sequence, i.e., we have

$$(\forall k \in \mathbb{N})(\exists N \in \mathbb{N})(\forall n, m \geq N)(d(w_n, w_m) < \frac{1}{2^k}).$$

Then $(w_n)_{n \in \mathbb{N}}$ is either eventually constant or enumerates all reals in $[0, 1]$. The latter is impossible by Theorem 1.5, i.e., $(w_n)_{n \in \mathbb{N}}$ converges to some $w \in M$. Note that a continuous function is trivially sequentially continuous.

Sixth, to obtain $\text{NBI}_{[0,1]}$ from item (h) and higher, recall the set $N := \{w^{1^*} : (\forall i < |w|)(Y(w(i)) = i)\}$ and consider (N, d) , which is a metric space in the same way as for (M, d) . To show that (N, d) is sequentially compact, let $(w_n)_{n \in \mathbb{N}}$ be a sequence in N . In case $(\forall n \in \mathbb{N})(|w_n| < m)$ for some $m \in \mathbb{N}$, then $(w_n)_{n \in \mathbb{N}}$ contains at most m different elements, as Y is an injection. The pigeon hole principle now implies that (at least) one w_{n_0} occurs infinitely often in $(w_n)_{n \in \mathbb{N}}$, yielding an obviously convergent sub-sequence. In case $(\forall m \in \mathbb{N})(\exists n \in \mathbb{N})(|w_n| \geq m)$, the sequence $(w_n)_{n \in \mathbb{N}}$ enumerates the reals in $[0, 1]$ (as Y is a bijection), which is impossible by Theorem 1.5. Thus, (N, d) is a sequentially compact space; the function $G : N \rightarrow \mathbb{R}$ defined as $G(u) = |u|$ is continuous (in the same way as for F above) but not bounded. This contradiction establishes that item (h) implies $\text{NBI}_{[0,1]}$, and the same for items (g)–(k). For item (l), the function $H(x, k) := \frac{1}{2^{|x|+k+2}}$ is a modulus of continuity for G .

Seventh, for item (m), assume again $\neg \text{NBI}_{[0,1]}$ and define $G_n(w)$ as $|w|$ in case $|w| \leq n$, and 0 otherwise. As for G above, G_n is continuous and $\lim_{n \rightarrow \infty} G_n(w) = G(w)$ for $x \in N$. Since $G_n \leq G_{n+1}$ on N , item (m) implies that the convergence is uniform, i.e., we have

$$(\forall k \in \mathbb{N})(\exists m \in \mathbb{N})(\forall w \in N)(\forall n \geq m)(|G_n(w) - G(w)| < \frac{1}{2^k}), \tag{2.2}$$

which is clearly false. Indeed, take $k = 1$ and let $m_1 \in \mathbb{N}$ be as in (2.2). Since Y is surjective, IND_1 provides $w_1 \in N$ of length $m_1 + 1$, yielding $|G(w_1) - G_{m_1}(w_1)| = |(m_1 + 1) - 0| > \frac{1}{2}$, contradicting (2.2) and thus $\text{NBI}_{[0,1]}$ follows from item (m). For item (n), $(G_n)_{n \in \mathbb{N}}$ is equicontinuous by the previous, but not uniformly equicontinuous, just like for item (m) using a variation of (2.2). For item (o), the function $J(w) := \frac{1}{2^{|w|}}$ is continuous on N in the same way as for F, G . However, assuming $\neg \text{NBI}_{[0,1]}$, J becomes arbitrarily small on N , contradicting item (o). For item (p), define $J_n(w)$ as $J(w)$ if $|w| \leq n$, and 1 otherwise. Similar to the previous, J_n converges to J , but not uniformly, i.e., item (p) also implies $\text{NBI}_{[0,1]}$.

For item (q), note that $O_n := \{w \in N : |w| = n\}$ is open as $B_d^M(v, \frac{1}{2^{n+2}}) \subset O_n$ in case $v \in O_n$. Then $\cup_{n \in \mathbb{N}} O_n$ covers N , assuming N (and $\neg \text{NBI}_{[0,1]}$) as above. However, there clearly is no finite sub-covering.

Finally, for items (r) and (s), the above proof for items (e) and (f) goes through without modification. For item (t), note that N is an infinite set in (N, d) without limit point. The final sentence speaks for itself: one uses (\exists^3) and (μ^2) to obtain a modulus of continuity. For $\varepsilon = 1$, the latter yields an uncountable covering, which has finite sub-covering assuming (\exists^3) by [28, Theorem 4.1]. This immediately yields an upper bound while the supremum and maximum are obtained using the usual interval-halving technique using (\exists^3) . \dashv

We could restrict item (q) to *R2-open* sets [27, 33], where the latter are open sets such that $x \in U$ implies $B(x, h_U(x)) \subset U$ with the notation of Definition 1.2.

2.2. Variations on a theme. Lest the reader believe that third-order metric spaces are somehow irredeemable, we show that certain (very specific) variations of the items in Theorem 2.2 are provable in rather weak systems, sometimes assuming countable choice as in $\text{QF-AC}^{0,1}$ (Theorems 2.3 and 2.4). We also show that certain items in Theorem 2.2 are just very hard to prove by deriving some of the new ‘Big’ systems from [30, 31, 38, 40], namely the Jordan decomposition theorem and the uncountability of \mathbb{R} as in $\text{NIN}_{[0,1]}$ (Theorem 2.6).

First, we establish the following theorem, which suggests a strong need for open sets as in Definition 1.2 if we wish to prove basic properties of metric spaces in the base theory, potentially extended with the Big Five. The fourth item should be contrasted with item (e) in Theorem 2.2. Many variations of the below results are of course possible based on the associated second-order results.

THEOREM 2.3 (RCA_0^w).

- (a) *For strongly countably open (M, d) , a continuous $F : M \rightarrow \mathbb{R}$ is bounded.*

- (b) *Dini's theorem for strongly countably-compact (M, d) .*
- (c) *Pincherle's theorem for strongly countably-compact (M, d) .*
- (d) *A metric space (M, d) with the Cantor intersection property, is strongly countably-compact.*
- (e) *The following are equivalent:*
 - (e.1) *weak König's lemma WKL_0 ,*
 - (e.2) *for any weakly complete and effectively totally bounded metric space (M, d) with $M \subset [0, 1]$, a continuous $F : M \rightarrow \mathbb{R}$ is bounded above,*
 - (e.3) *the previous item for sequentially continuous functions.*
- (f) *The following are equivalent.*
 - (f.1) *arithmetical comprehension ACA_0 ,*
 - (f.2) *any weakly complete and effectively totally bounded metric space (M, d) with $M \subset [0, 1]$, is sequentially compact.*

PROOF. For the first item, since F is continuous, the set $E_n := \{x \in M : |F(x)| > n\}$ is open and exists in the sense of Definition 1.2. Since $\cup_{n \in \mathbb{N}} E_n$ covers (M, d) , there is a finite sub-covering $\cup_{n \leq n_0} E_n$ for some $n_0 \in \mathbb{N}$, implying $|F(x)| \leq n_0 + 1$ for all $x \in M$, i.e., F is bounded as required.

For the second item, let F, F_n be as in Dini's theorem and define $G_n(w) := F(w) - F_n(w)$. Now fix $k \in \mathbb{N}$ and define $E_n := \{w \in M : G_n(w) < \frac{1}{2^k}\}$. The latter yields a countable open covering and one obtains uniform convergence from any finite sub-covering. For the third item, fix $F : M \rightarrow \mathbb{R}^+$ and define $E_n := \{w \in M : F(w) > \frac{1}{2^n}\}$. The proof proceeds as for the previous items.

For the fourth item, this amounts to a manipulation of definitions. For the fifth item, that (e.2) and (e.3) imply WKL_0 is immediate by [32, Theorem 2.8] for $M = [0, 1]$ and [19, Proposition 3.6]. For the downward implication, fix $F : M \rightarrow \mathbb{R}$ for $M \subset [0, 1]$ as in item (e.2). In case $\neg(\exists^2)$, all functions on \mathbb{R} are continuous by [19, Proposition 3.12]. By [32, Theorem 2.8], all (continuous) $[0, 1] \rightarrow \mathbb{R}$ -functions are bounded. Since we may (also) view F as a (continuous) function from reals to reals, F is bounded on $[0, 1]$ and hence M , i.e., this case is finished.

In case (\exists^2) , we follow the well-known proof to show that (M, d) is sequentially compact. Indeed, for a sequence $(x_n)_{n \in \mathbb{N}}$ in M , define a sub-sequence as follows: M can be covered by a finite number of balls with radius $1/2^k$ with $k = 1$. Find a ball with infinitely many elements of $(x_n)_{n \in \mathbb{N}}$ inside (which can be done explicitly using (\exists^2)) and choose x_{n_0} in this ball to define $y_0 := x_{n_0}$. Now repeat the previous steps for $k > 1$ and note that the resulting sequence is effectively Cauchy and hence convergent (by the assumptions on M). Hence, (M, d) is sequentially compact and suppose $F : M \rightarrow \mathbb{R}$ is unbounded, i.e., $(\forall n \in \mathbb{N})(\exists x \in M)(F(x) > n)$. It is now important to note that the underlined quantifier can be replaced by a quantifier over \mathbb{N} using the sequence $(w_n)_{n \in \mathbb{N}}$ provided by M being effectively totally bounded.

Applying $\text{QF-AC}^{0,0}$, included in RCA_0^ω , there is a sequence $(x_n)_{n \in \mathbb{N}}$ such that $|F(x_n)| > n$. This sequence has a convergent sub-sequence, say with limit y , and F is not continuous at y , a contradiction. Thus, F is bounded for both disjuncts of $(\exists^2) \vee \neg(\exists^2)$. The equivalence involving ACA_0 has a similar proof. \dashv

The final part of the proof seems to crucially depend on *effective* totally boundedness. Indeed, by the first part of Theorem 2.2, item (e.3) of Theorem 2.3 with ‘effectively’ omitted, implies $\text{NBI}_{[0,1]}$. In other words, the equivalences in Theorem 2.3 do not seem robust.

Secondly, we show that certain items from Theorem 2.2 fit nicely with RM, assuming an extended base theory. Other items turn out to be connected to the ‘new’ Big systems studied in [30, 38, 39].

We now show that certain items from Theorem 2.2 are provable assuming countable choice as in $\text{QF-AC}^{0,1}$. Thus, these items do not imply $\text{NIN}_{[0,1]}$ as the latter is not provable in $\text{Z}_2^\omega + \text{QF-AC}^{0,1}$. The third item should be contrasted with [42, III.2]. Many results in RM do not go through in the absence of $\text{QF-AC}^{0,1}$, as studied at length in [27, 28].

THEOREM 2.4 ($\text{RCA}_0^\omega + \text{QF-AC}^{0,1}$). *The following are provable for (M, d) any metric space with $M \subset \mathbb{R}$.*

- *Items (h), (g), (m), (n), (o), (q), (s), and (t) from Theorem 2.2.*
- *The following are equivalent:*
 - *weak König’s lemma WKL_0 ,*
 - *the unit interval is strongly countably-compact.*
- *The following are equivalent:*
 - *arithmetical comprehension ACA_0 ,*
 - *a weakly complete and effectively totally bounded (M, d) with $M \subset [0, 1]$ is limit point compact.*

PROOF. First of all, we prove item (h) from Theorem 2.2 in $\text{RCA}_0^\omega + \text{QF-AC}^{0,1}$. To this end, suppose the continuous function $F : M \rightarrow \mathbb{R}$ is unbounded, i.e., $(\forall n \in \mathbb{N})(\exists w \in M)(|F(w)| > n)$. Applying $\text{QF-AC}^{0,1}$, there is a sequence $(x_n)_{n \in \mathbb{N}}$ such that $|F(x_n)| > n$. Since (M, d) is assumed to be sequentially complete, let $(y_n)_{n \in \mathbb{N}}$ be a convergent sub-sequence with limit $y \in M$. Clearly, F cannot be continuous at $y \in M$, a contradiction, which yields item (h). Item (g) is proved in the same way: suppose F is not uniformly continuous and apply $\text{QF-AC}^{0,1}$ to the latter statement to obtain a sequence. Then F is not continuous at the limit of the convergent sub-sequence. Items (m)–(o) are proved in the same way. To prove item (q), let $(O_n)_{n \in \mathbb{N}}$ be a countable open covering of M with $(\forall n \in \mathbb{N})(\exists x \in M)(x \notin \cup_{m \leq n} O_m)$. Apply $\text{QF-AC}^{0,1}$ to obtain a sequence $(x_n)_{n \in \mathbb{N}}$, which has a convergent sub-sequence $(y_n)_{n \in \mathbb{N}}$ by assumption, say with limit $y \in M$. Then $y \in O_{n_0}$ for some $n_0 \in \mathbb{N}$, which implies that y_n is also eventually in O_{n_0} ,

a contradiction. To prove item (s), let $(C_n)_{n \in \mathbb{N}}$ be as in the sequential Cantor intersection property and apply QF-AC^{0,1} to $(\forall n \in \mathbb{N})(\exists x \in M)(x \in C_n)$. The convergent sub-sequence has a limit $y \in \bigcap_{n \in \mathbb{N}} C_n$. To prove item (t), let X be an infinite set, i.e., $(\forall N \in \mathbb{N})(\exists w^{1^*})(\forall i < |w|)(|w| = N \wedge w(i) \in X)$. Now apply QF-AC^{0,1} to obtain a sequence $(w_n)_{n \in \mathbb{N}}$ in X . Since (M, d) is sequentially closed, the latter sequence has a convergent sub-sequence, the limit of which is a limit point of X .

Secondly, the equivalence in the second item is proved in [27, Theorem 4.1]. For the third item, the upwards implication is immediate for $M = [0, 1]$. For the downwards implication, assume (M, d) as in the final sub-item. Theorem 2.3 implies that (M, d) is sequentially compact. As in the previous paragraph, an infinite set in M now has a limit point. −

A similar proof should go through for many of the other items in Theorem 2.2 and for QF-AC^{0,1} replaced by NCC from [29]; the latter is provable in Z_2^Ω while the former is not provable in ZF.

Secondly, the Jordan decomposition theorem is studied in [30, 40] where various versions are shown to be equivalent to the enumeration principle for countable sets. Many equivalences exist for the following principle, elevating it to a new ‘Big’ system, as shown in [30].

PRINCIPLE 2.5 (cocode₀). *Let $A \subset [0, 1]$ and $Y : [0, 1] \rightarrow \mathbb{N}$ be such that Y is injective on A . Then there is a sequence of reals $(x_n)_{n \in \mathbb{N}}$ that includes A .*

This principle is ‘explosive’ in that $ACA_0^\omega + \text{cocode}_0$ proves ATR_0 and $\Pi_1^1\text{-}CA_0^\omega + \text{cocode}_0$ proves $\Pi_2^1\text{-}CA_0$ (see [30, Section 4]). As it turns out, the separability of metric spaces is similarly explosive.

THEOREM 2.6 (ACA₀^ω).

- Item (f) or (r) from Theorem 2.2 implies cocode_0 .
- Item (f) or (r) for $M = [0, 1]$ from Theorem 2.2 implies $\text{NIN}_{[0,1]}$.

PROOF. For the first item, let $Y : [0, 1] \rightarrow \mathbb{N}$ be injective on $A \subset [0, 1]$; without loss of generality, we may assume $0 \in A$. Now define $d(x, y) := |\frac{1}{2^{Y(x)}} - \frac{1}{2^{Y(y)}}|$, $d(x, 0) = d(0, x) := \frac{1}{2^{Y(x)}}$ for $x, y \neq 0$ and $d(0, 0) := 0$. The metric space (A, d) is countably-compact as $0 \in B_d^A(x, r)$ implies $y \in B_d^A(x, r)$ for $y \in A$ with only finitely many exceptions (as Y is injective on A). Similarly, (A, d) is sequentially compact: in case a sequence $(z_n)_{n \in \mathbb{N}}$ in A has at most finitely many distinct elements, there is an obvious convergent/constant sub-sequence. Otherwise, $(z_n)_{n \in \mathbb{N}}$ has a sub-sequence $(y_n)_{n \in \mathbb{N}}$ such that $Y(y_n)$ becomes arbitrary large with n increasing; this sub-sequence is readily seen to converge to 0.

Now let $(x_n)_{n \in \mathbb{N}}$ be the sequence provided by item (f) or (r) of Theorem 2.2, implying $(\forall x \in A)(\exists n \in \mathbb{N})(d(x, x_n) < \frac{1}{2^{Y(x)+1}})$ by taking

$k = Y(x) + 1$. The latter formula implies

$$(\forall x \in A)(\exists n \in \mathbb{N})(x \neq_{\mathbb{R}} 0 \rightarrow |\frac{1}{2^{Y(x)}} - \frac{1}{2^{Y(x_n)}}| <_{\mathbb{R}} \frac{1}{2^{Y(x)+1}}) \quad (2.3)$$

by definition. Note that x_n from (2.3) cannot be 0 by the definition of the metric d . Clearly, $|\frac{1}{2^{Y(x)}} - \frac{1}{2^{Y(x_n)}}| < \frac{1}{2^{Y(x)+1}}$ is only possible if $Y(x) = Y(x_n)$, implying $x =_{\mathbb{R}} x_n$. Hence, we have shown that $(x_n)_{n \in \mathbb{N}}$ lists all reals in $A \setminus \{0\}$. The same proof now yields the second item for $A = [0, 1]$ as Theorem 1.5 implies the reals cannot be enumerated. \dashv

In conclusion, the coding of metric spaces does distort the logical properties of basic properties of continuous functions on metric spaces by Theorem 2.2. This is established by deriving $\text{NBI}_{[0,1]}$ while noting that $\text{NIN}_{[0,1]}$ generally does not follow by Theorem 2.4. The latter also shows that in an enriched base theory, one can obtain ‘rather vanilla’ RM. By contrast, other properties of metric spaces imply new ‘Big’ systems, as is clear from Theorem 2.6.

§3. Foundational musings.

3.1. Thoughts on coding. The results in this paper have implications for the coding of higher-order objects in second-order RM, as discussed in this section.

First of all, our results shed new light on the following problem from [11, page 135].

PROBLEM. [...] Show that Simpson’s neighborhood condition coding of partial continuous functions between complete separable metric spaces is “optimal”.

A coding is called *optimal* in [11] in case RCA_0 can prove ‘as much as possible’, i.e., as many as possible of the basic properties of the coding can be established in RCA_0 . Theorem 2.2 show that without separability, basic properties of continuous functions on compact metric spaces are no longer provable from second-order (comprehension) axioms. Thus, separability is an essential ingredient *if* one wishes to study these matters using second-order arithmetic/axioms.

Secondly, second-order (comprehension) axioms can establish many (third-order) theorems about continuous *and* discontinuous functions on the reals (see [32, 40]), assuming RCA_0^ω . Hence, large parts of (third-order) real analysis can be developed using second-order comprehension axioms in a weak third-order background theory, namely RCA_0^ω , using little-to-no-coding. The same does not hold for continuous functions on compact metric spaces by the above results. In particular, Theorem 2.3 suggests we have to choose a *very specific* representation, namely ‘weakly complete and effectively totally bounded’ to obtain third-order statements that are

classified in the Big Five. Indeed, Theorem 2.2 implies that many (most?) other variations are not provable from second-order (comprehension) axioms.

In conclusion, our results show that separability is an essential ingredient *if* one wishes to study these matters using second-order arithmetic/axioms. However, our results also show that this is a *very specific* choice that is ‘non-standard’ in the sense that many variations cannot be established using second-order arithmetic/axioms.

3.2. Set theory and ordinary mathematics. In this section, we explore a theme introduced in [39]. Intuitively speaking, we collect evidence for a parallel between our results and some central results in set theory. Formulated slightly differently, one could say that interesting phenomena in set theory have ‘miniature versions’ to be found in third-order arithmetic, or that the seeds for interesting phenomena in set theory can already be found in third-order arithmetic.

First of all, the cardinality of \mathbb{R} is mercurial in nature: the famous work of Gödel [12] and Cohen [6, 7] shows that the *Continuum Hypothesis* cannot be proved or disproved in ZFC, i.e., Zermelo-Fraenkel set theory with AC, the usual foundations of mathematics. In particular, the exact cardinality of \mathbb{R} cannot be established in ZFC. A parallel observation in higher-order RM is that $Z_2^\omega + \text{QF-AC}^{0,1}$ cannot prove that \mathbb{R} is uncountable in the sense of there being no injection from \mathbb{R} to \mathbb{N} (see [31] for details). In a conclusion, the cardinality of \mathbb{R} has a particularly mercurial nature, in both set theory and higher-order arithmetic.

Secondly, many standard results in mainstream mathematics are not provable in ZF, i.e., ZFC with AC removed, as explored in great detail [14]. The absence of AC is even said to lead to *disasters* in topology and analysis (see [17]). A parallel phenomenon was observed in [27, 28], namely that certain rather basic equivalences go through over $\text{RCA}_0^\omega + \text{QF-AC}^{0,1}$, but not over Z_2^ω .

Examples include the equivalence between compactness results and local-global principles, which are intimately related according to Tao [45]. In this light, it is fair to say that disasters happen in both set theory and higher-order arithmetic in the absence of AC. It should be noted that $\text{QF-AC}^{0,1}$ (not provable in ZF) can be replaced by NCC from [29] (provable in Z_2^Ω) in the aforementioned.

Thirdly, we discuss the essential role of AC in measure and integration theory, which leads to rather concrete parallel observations in higher-order arithmetic. Indeed, the full pigeonhole principle for measure spaces is not provable in ZF, which immediately follows from, e.g., [14, Diagram 3.4]. A parallel phenomenon in higher-order arithmetic (see [39]) is that even the restriction to closed sets, namely $\text{PHP}_{[0,1]}$ cannot be proved in $Z_2^\omega + \text{QF-AC}^{0,1}$ (but Z_2^Ω suffices).

A more ‘down to earth’ observation pertains to the intuitions underlying the Riemann and Lebesgue integral. Intuitively, the integral of a non-negative function represents the area under the graph; thus, if the integral is zero, then this function must be zero for ‘most’ reals. Now, AC is needed to establish this intuition for the Lebesgue integral [16]. Similarly, [39, Theorem 3.8] establishes the parallel observation that this intuition for the *Riemann* integral cannot be proved in $Z_2^\omega + \text{QF-AC}^{0,1}$ (but Z_2^Ω suffices as usual).

Fourth, the *pointwise* equivalence between sequential and ‘epsilon-delta’ continuity cannot be proved in ZF while $\text{RCA}_0^\omega + \text{QF-AC}^{0,1}$ suffices for functions on Baire space (see [19]). A parallel observation is provided by (the proof of) Theorem 2.2, namely that the following statement is not provable in Z_2^ω :

for countably-compact (M, d) and sequentially continuous $F : M \rightarrow \mathbb{R}$, F is continuous on M .

Thus, the *global* equivalence between sequential and ‘epsilon-delta’ continuity on metric spaces cannot be proved in Z_2^ω . In other words, the exact relation between sequential and ‘epsilon-delta’ continuity is hard to pin down, both in set theory and third-order arithmetic.

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