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Division algebras and MRD codes from skew polynomials

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Abstract

Let *D* be a division algebra, finite-dimensional over its center, and $R = D[t; \sigma, \delta]$ a skew polynomial ring. Using skew polynomials $f \in R$, we construct division algebras and maximum rank distance codes consisting of matrices with entries in a noncommutative division algebra or field. These include Jha Johnson semifields, and the classes of classical and twisted Gabidulin codes constructed by Sheekey.

1. Introduction

Rank distance codes are important both in coding theory and cryptography. One of the best-known maximum rank distance (MRD) codes is probably the Gabidulin code [10] which was mentioned already by Delsarte [8]. In coding theory, MRD codes are well suited to correct errors [6, 31]. In cryptography, they are used to design public-key cryptosystems, see for instance [9, 12].

MRD codes over general (non-finite) fields, in particular number fields, were already studied in [2] and later touched on in [34]. Rank metric codes over both cyclic and more general Galois extensions were considered in [3, 31, 32]. Although rank metric codes have been also constructed over finite principal ideal rings [19] and discretely valued rings [21], to our knowledge they have not yet been studied over noncommutative rings. In this paper, we also consider MRD codes in $M_k(B)$, where B is a noncommutative division algebra.

We construct these MRD codes using skew polynomials. Skew polynomials have been successfully used in constructions of both division algebras (mostly semifields) and linear codes [2, 4, 5, 13, 26–28], in particular building space-time block codes (STBCs) [29] and MRD codes [33, 34].

Our codes can be seen as generalizations of both the classical and twisted Gabidulin codes in [10], resp., [33]. We put Sheekey's construction [34] in a broader context which helps to understand it better and potentially allows other ways to generalize MRD coding using skew polynomials. The drawback is that rather early on we have to rigorously restrict the choice of the polynomials f we can employ and that the construction remains rather theoretical.

Sheekey [34] only considers skew polynomials $f \in K[t; \sigma]$ with coefficients in cyclic Galois field extensions for his construction and limits himself to the case that the minimal central left multiple of fhas maximal degree. He misses out on codes (with matrix entries both in a noncommutative division algebra, and with entries in fields) and algebras that can be obtained by employing skew polynomials with coefficients in a noncommutative division algebra. He also misses out on constructions using $f \in D[t; \delta]$. We construct both new division algebras and MRD codes with entries in a noncommutative division algebra, and with entries in fields.

The first five Sections of the paper contain the preliminaries (Section 1) and theoretical background needed to obtain the main results (Sections 2-5). Let *D* be a division algebra of degree *d* over its center,

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and $f \in R = D[t; \sigma, \delta]$ a monic irreducible skew polynomial with a bound that lies in the center C(R) of R.

While developing the theory, we point out how the choice of *D* and the polynomial *f* has to be restricted in order to construct both division algebras and MRD codes out of *f*, a scalar $v \in D$ and a suitable $\rho \in Aut(D)$.

Apart from Section 9, we fix the following general assumptions unless specified otherwise: $R = D[t; \sigma]$, where σ is an automorphism of D of finite order n modulo inner automorphisms, i.e. $\sigma^n = i_u$ for some inner automorphism $i_u(z) = uzu^{-1}$, and $F = C \cap \text{Fix}(\sigma)$. Choose $\rho \in \text{Aut}(D)$, such that F/F' with $F' = \text{Fix}(\rho) \cap F$ is finite-dimensional. Let $\nu \in D^{\times}$.

Let $f \in R$ be monic and irreducible of degree m > 1, and h the minimal central left multiple of f, so that $R/Rh \cong M_k(B)$ for some division algebra B (Theorem 3). Let l < k be a positive integer. Define $S_{n,m,l}(v, \rho, f) = \{a + Rh \mid a \in P\} \subset R/Rh$ with the set $P = \{d_0 + d_1t + \cdots + d_{lm-1}t^{lm-1} + v\rho(d_0)t^{lm} \mid d_i \in D\}$. Let $L_a : R/Rf \to R/Rf$ be the left multiplication map $L_a(b + Rf) = ab + Rf$. We have well-defined maps $S_{n,m,l}(v, \rho, f) \longrightarrow \text{End}_B(R/Rf) \longrightarrow M_k(B), a \mapsto L_a \mapsto M_a$, where M_a is the matrix representing L_a with respect to a right B-basis of R/Rf. The image $C_{n,m,l} = \{M_a \mid a \in S_{n,m,l}(v, \rho, f)\}$ of $S_{n,m,l}(v, \rho, h)$ in $M_k(B)$ is an F'-linear rank metric code. If $C_{n,m,l}$ has distance $d_C = k - l + 1$, then $C_{n,m,l}$ is called a maximum rank distance code in $M_k(B)$. We will usually deal with the case that deg(h) = dmn, so that B is a field.

The most general results are contained in Section 6: If *P* does not contain a polynomial of degree *lm*, whose irreducible factors are all similar to *f*, then $C_{n,m,l}$ is an *F*'-linear MRD code in $M_k(B)$ with minimum distance k - l + 1 (Theorem 19).

Furthermore, let $D = (E/C, \gamma, a)$ be a cyclic division algebra such that $\sigma|_E \in \operatorname{Aut}(E)$ and $\gamma \circ \sigma|_E = \sigma|_E \circ \gamma$, and $\sigma^n(z) = u^{-1}zu$ for some $u \in E$. Let $f(t) = \sum_{i=0}^m a_i t^i \in E[t; \sigma]$ be a monic irreducible polynomial of degree *m*, such that deg(*h*) = *dmn*, and such that all monic f_i similar to *f* lie in $E[t; \sigma]$. Then, the algebra $S_{n,m,1}(v, \rho, f)$ is a division algebra, if one of the following holds: (i) $v \notin E$ and $\rho|_E \in \operatorname{Aut}(E)$; (ii) $v \in E^{\times}$ and $\rho|_E \in \operatorname{Aut}(E)$, such that $N_{E/F'}(a_0)N_{E/F'}(v) \neq 1$ (Theorem 16). MRD codes are canonically obtained from the matrices representing the left multiplication of these division algebras.

In Section 7, the nuclei of the algebras and codes are investigated. We give some examples of algebras obtained from our construction employing $f(t) = t^n - \theta \in K[t; \sigma]$ in Section 8.

We conclude with a brief look at the constructions using a differential polynomial $f \in D[t; \delta]$, where the center of *D* is a field of characteristic *p*, in Section 9.

The fact that we are using $f \in D[t; \sigma]$, respectively $f \in D[t; \gamma]$, means we have a larger choice of skew polynomials to build codes that Sheekey does, who only considers f with coefficients in a cyclic field extension.

This work is part of the second author's PhD thesis [35].

2. Preliminaries

2.1. Nonassociative algebras

Let *F* be a field. We call *A* an *algebra* over *F* if there exists an *F*-bilinear map $A \times A \rightarrow A$, $(x, y) \mapsto x \cdot y$, denoted simply by juxtaposition *xy*, the *multiplication* of *A*. An algebra *A* is called *unital* if there is an element in *A*, denoted by 1, such that 1x = x1 = x for all $x \in A$. We will only consider unital algebras. A nonassociative algebra $A \neq 0$ is called a *division algebra* if for any $a \in A$, $a \neq 0$, the left multiplication with *a*, $L_a(x) = ax$, and the right multiplication with *a*, $R_a(x) = xa$, are bijective. If *A* is finite-dimensional as an *F*-vector space, then *A* is a division algebra if and only if *A* has no zero divisors. The *left nucleus* of *A* is defined as Nuc_{*l*}(*A*) = { $x \in A \mid [x, A, A] = 0$ }, the *middle nucleus* of *A* is Nuc_{*m*}(*A*) = { $x \in A \mid [A, x, A] = 0$ }, where [x, y, z] = (xy)z - x(yz) is the *associator*. Nuc_{*l*}(*A*), Nuc_{*m*}(*A*), and Nuc_{*r*}(*A*) are associative subalgebras of *A*. Their intersection Nuc(*A*) = { $x \in A \mid [x, A, A] = [A, x, A] = [A, A, x] = 0$ } is the *nucleus* of *A*. Nuc(*A*) is an associative subalgebra of *A*, and x(yz) = (xy)z whenever one of the elements x, y, z is in Nuc(*A*). The *center* of *A* is C(*A*) = { $x \in$ Nuc(*A*) | xy = yx for all $y \in A$ }.

Let *A* be a finite-dimensional central simple associative algebra over *F* of degree *d* and let \overline{F} denote the algebraic closure of *F*. Then, $A \otimes_F \overline{F} \cong M_d(\overline{F})$, so that we can fix an embedding $A \longrightarrow M_d(\overline{F})$ and view every $a \in A$ as a matrix in $M_d(\overline{F})$. The characteristic polynomial

$$m_a(X) = X^d - s_1(a)X^{d-1} + s_2(a)X^{d-2} - \dots + (-1)^d s_d(a),$$

of $a \in A$ has coefficients in F and is independent of the choice of the embedding. The coefficient $N_A(a) = s_d(a)$ is called the *reduced norm of a* [20]. Let K/F be a cyclic Galois extension of degree d with Galois group $Gal(K/F) = \langle \gamma \rangle$ and norm $N_{K/F}$. Let $c \in F^{\times}$. An *associative cyclic algebra* $(K/F, \gamma, c)$ of degree d over F is a d-dimensional K-vector space

$$(K/F, \gamma, c) = K \oplus eK \oplus e^2K \oplus \cdots \oplus e^{d-1}K,$$

with multiplication given by the relations $e^d = c$, $le = e\sigma(l)$, for all $l \in K$. $(K/F, \gamma, c)$ is a division algebra for all $c \in F^{\times}$, such that $c^s \notin N_{K/F}(K^{\times})$ for all *s* which are prime divisors of *d*, $1 \le s \le d - 1$.

2.2. MRD codes

Let *K* be a field. A *code* is a set of matrices $C \subset M_{n,m}(K)$. Let $L \subset K$ be a subfield, then C is *L*-linear if C is a vector space over *L*. A *rank metric code* is a code $C \subset M_{n,m}(K)$ equipped with the rank distance function $d(X, Y) = \operatorname{rank}(X - Y)$. Define the *minimum distance* of a rank metric code C as

 $d_{\mathcal{C}} = \min\{d(X, Y) \mid X, Y \in \mathcal{C}, X \neq Y\}.$

An L-linear rank metric code C satisfies the Singleton-like bound

$$\dim_L(\mathcal{C}) \le n(m - d_{\mathcal{C}} + 1)[K:L],$$

where dim_{*L*}(C) is the dimension of the *L*-vector space C [2, Proposition 6].

An *L*-linear rank metric code attaining the Singleton-like bound is called a *maximum rank distance code* or *MRD code* (for MRD codes over cyclic field extensions see [2]).

If now *B* is a not necessarily commutative division algebra then more generally, we again define a *code* as a set of matrices $C \subset M_{n,m}(B)$. Let $B' \subset B$ be a subalgebra, then C is B'-linear (or simply linear), if C is a right B'-module.

A *rank metric code* $C \subset M_{n,m}(B)$ is a code together with the distance function

$$d(X, Y) = \operatorname{colrank}(X - Y),$$

for all $X, Y \in M_{n,m}(B)$, where colrank is the column rank of A (the rank of the right B-module generated by the columns of A). A matrix in $M_{n,m}(B)$ has column rank at most m; any matrix which attains this bound is said to have attained *full column rank*. The *minimum distance* of a rank metric code $C \subset M_{n,m}(B)$ is defined as

$$d_{\mathcal{C}} = \min\{d(X, Y) \mid X, Y \in \mathcal{C}, X \neq Y\}.$$

To our knowledge, such codes $\mathcal{C} \subset M_{n,m}(B)$ have not previously been considered in the literature.

2.3. Skew polynomial rings

In the following, let *D* be a central simple division algebra of degree *d* over its center *C*, σ a ring endomorphism of *D* and $\delta : D \to D$ a *left* σ -*derivation*, i.e. an additive map such that $\delta(ab) = \sigma(a)\delta(b) + \delta(a)b$ for all $a, b \in D$. The *skew polynomial ring* $D[t; \sigma, \delta]$ is the set of skew polynomials $g(t) = a_0 + a_1t + \cdots + a_nt^n$ with $a_i \in D$, with term-wise addition and multiplication defined via $ta = \sigma(a)t + \delta(a)$ for all $a \in D$ [22]. Define Fix(σ) = { $a \in D | \sigma(a) = a$ } and Const(δ) = { $a \in D | \delta(a) = 0$ }. If $\delta = 0$, define $D[t; \sigma] = D[t; \sigma, 0]$. If $\sigma = id$, define $D[t; \delta] = D[t; id, \delta]$.

For $f(t) = a_0 + a_1 t + \dots + a_n t^n \in R = D[t; \sigma, \delta]$ with $a_n \neq 0$, we define the degree of f as deg(f) = n and deg $(0) = -\infty$. A skew polynomial $f \in R$ is *irreducible* if it is not a unit and it has no proper factors,

i.e if there do not exist $g, h \in R$ with $1 \le \deg(g)$, deg $(h) < \deg(f)$ such that f = gh [18, p. 2 ff.]. We call $f \in R$ right-invariant if Rf is a left and a right ideal in R, and a two-sided maximal element, if f is right-invariant and Rf is a nonzero maximal ideal in R (equivalently, if $f \ne 0$ and R/Rf is a simple ring) [18, p. 13]. Two nonzero skew polynomials $f_1, f_2 \in R$ are similar, written $f_1 \sim f_2$, if $R/Rf_1 \cong R/Rf_2$ [18, p. 11].

A skew polynomial $f \in R$ is *bounded* if there exists a nonzero polynomial $f^* \in R$ such that Rf^* is the largest two-sided ideal of R contained in Rf. The polynomial f^* is uniquely determined by f up to scalar multiplication by elements of D^{\times} and is called a *bound* of f.

If $f \in R$ has degree *m*, then for all $g \in R$ of degree $l \ge m$, there exist uniquely determined $r, q \in R$ with deg $(r) < \deg(f)$, such that g = qf + r. Let $\operatorname{mod}_{x} f$ denote the remainder of right division by *f*. The skew polynomials $R_m = \{g \in R \mid \deg(g) < m\}$ of degree less that *m* canonically represent the elements of the left *R*-modules R/Rf. Furthermore, R_m together with the multiplication $g \circ h = gh \operatorname{mod}_{x} f$ is a unital nonassociative algebra $S_f = (R_m, \circ)$ over $F_0 = \{a \in D \mid ah = ha \text{ for all } h \in S_f\} = \operatorname{Comm}(S_f) \cap D$, called a *Petit algebra*. When the context is clear, we simply use juxtaposition for multiplication in S_f . Note that $C(D) \cap \operatorname{Fix}(\sigma) \cap \operatorname{Const}(\delta) \subset F_0$. For all $a \in D^{\times}$, we have $S_f = S_{af}$; thus, without loss of generality we can assume *f* is monic when working with Petit algebras S_f . If *f* has degree 1 then $S_f \cong D$.

Lemma 1. Let R be a ring with no zero divisors. For all $g \in C(R)$, every right divisor of g in R also divides g on the left.

Proof. Suppose γ is a right divisor of g. Then, $g = \delta \gamma$ for some $\delta \in R$. As g lies in the center of R, we have $\delta g = g\delta = \delta \gamma \delta$. This rearranges to $0 = \delta g - \delta \gamma \delta = \delta(g - \gamma \delta)$. As R contains no zero divisors and $\delta \neq 0$, it follows that $g = \gamma \delta$.

2.4. The minimal central left multiple of $f \in D[t; \sigma]$

From now on let, σ be an automorphism of D of finite order n modulo inner automorphisms, i.e. $\sigma^n = i_u$ for some inner automorphism $i_u(z) = uzu^{-1}$. Then, the order of $\sigma|_C$ is n. W.l.o.g., we choose $u \in Fix(\sigma)$. Let $R = D[t; \sigma]$ and define $F = C \cap Fix(\sigma)$. R has center

$$C(R) = F[u^{-1}t^n] = \left\{ \sum_{i=0}^k a_i (u^{-1}t^n)^i \mid a_i \in F \right\} \cong F[x]$$

with $x = u^{-1}t^n$ [18, Theorem 1.1.22]. All polynomials $f \in R$ are bounded.

For any $f \in R = D[t; \sigma]$ with a bound in C(R), we define the *minimal central left multiple mclm*(f) of f in R to be the unique polynomial of minimal degree $h \in C(R) = F[u^{-1}t^n]$ such that h = gf for some $g \in R$, and such that $h(t) = \hat{h}(u^{-1}t^n)$ for some monic $\hat{h}(x) \in F[x]$. Define $E_{\hat{h}} = F[x]/(\hat{h}(x))$. If f has nonzero constant term, then $f^* \in C(R)$ [11, Lemma 2.11]). From now on, we assume that f has nonzero constant term and denote by $h \in C(R)$, $h(t) = \hat{h}(u^{-1}t^n)$, the minimal central left multiple of f. Then, h equals the bound of f up to a scalar multiple from D. If f is irreducible in R, then $\hat{h}(x)$ is irreducible in F[x]. If $\hat{h} \in F[x]$ is irreducible, then $f = f_1 \cdots f_r$ for irreducible $f_i \in R$ such that $f_i \sim f_i$ for all i, j ([23], cf. [36]).

Lemma 2. Let $f \in R$.

- (i) If $f \in R$ is irreducible, then every $g \in R$ similar to f has h as its minimal central left multiple.
- (ii) Suppose that $h \in F[x]$ is irreducible. Then, $f = f_1 \cdots f_r$ for irreducible $f_i \in R$ such that $f_i \sim f_j$ for all i, j.

This follows easily from [7, p. 9, Corollary 2] and [18, Theorem 1.2.9].

The quotient algebra R/Rh has center $C(R/Rh) \cong F[x]/(\hat{h}(x))$, cf. [11, Lemma 4.2]. Define $E_{\hat{h}} = F[x]/(\hat{h}(x))$. Suppose that $\hat{h}(x) \neq x$ and that \hat{h} is irreducible in F[x]. Then, h generates a maximal two-sided ideal Rh in R [18, p. 16] and R/Rh is simple over its center $E_{\hat{h}}$.

Theorem 3 [23]. Let $f \in R = D[t; \sigma]$ be monic and irreducible of degree m > 1 with minimal central left multiple $h(t) = \hat{h}(u^{-1}t^n)$. Then,, Nuc_r(S_f) is a central division algebra over $E_{\hat{h}}$ of degree s = dn/k, where k is the number of irreducible factors of h in R, and

$$R/Rh \cong M_k(\operatorname{Nuc}_r(S_f)).$$

In particular, this means $\deg(\hat{h}) = \frac{dm}{s}$, $\deg(h) = km = \frac{dnm}{s}$, and

$$[\operatorname{Nuc}_r(S_f):F] = s^2 \cdot \frac{dm}{s} = dms.$$

Moreover, s divides gcd(dm, dn). If f is not right-invariant, then k > 1 *and s* \neq *dn.*

We know that $[S_f : F] = [S_f : C][C : F] = d^2m \cdot n$. Since $Nuc_r(S_f)$ is a subalgebra of S_f , comparing dimensions we obtain that

$$d^2mn = [S_f : F] = [S_f : \operatorname{Nuc}_r(S_f)] \cdot [\operatorname{Nuc}_r(S_f) : F] = k \cdot dms,$$

that is $[S_f : \operatorname{Nuc}_r(S_f)] = k$.

If *f* is not right-invariant which is equivalent to S_f being not associative, which in turn is equivalent to k > 1, then $s \neq dn$ looking at the degree of *h*. Note that deg(*h*) = *dnm* is the largest possible degree of *h*.

All of the above applies in particular to the special case that *D* is a finite field extension *K* of *C* of degree *n*, and $\sigma \in \operatorname{Aut}(K)$ has order *n*. Then, $R = K[t; \sigma]$ has center $C(R) = F[t^n] = \{\sum_{i=0}^k a_i(t^n)^i \mid a_i \in F\} = F[x]$ where $F = \operatorname{Fix}(\sigma)$ [18, Theorem 1.1.22].

3. Constructing sets of matrices employing irreducible $f \in D[t; \sigma]$

Let $R = D[t; \sigma]$ be as in Section 2.3 and $f \in R$ be an irreducible monic polynomial of degree m > 1 with nonzero constant term and minimal central left multiple $h(t) = \hat{h}(u^{-1}t^n)$. Let

$$E_f = \{z(t) + Rf \mid z(t) = \hat{z}(u^{-1}t^n) \in F[u^{-1}t^n]\} \subset R/Rf.$$

Together with the multiplication $(x + Rf) \circ (y + Rf) = (xy) + Rf$ for all $x, y \in F[u^{-1}t^n]$, E_f becomes an *F*-algebra.

Lemma 4. (i) For each $z(t) = \hat{z}(u^{-1}t^n) \in F[u^{-1}t^n]$ with $\hat{z} \in F[x]$, we have $z \in Rf$ if and only if $z \in Rh$. (ii) (E_f, \circ) is a field isomorphic to E_h .

Proof. (i) As h = gf for some $g \in R$, each $z \in Rh$ also lies in Rf.

Conversely, let $z(t) = \hat{z}(u^{-1}t^n) \in F[u^{-1}t^n]$ with $\hat{z} \in F[x]$ be such that $z \in Rf$. By the Euclidean division algorithm in F[x], there exist unique $\hat{q}(x), \hat{r}(x) \in F[x]$ such that $\hat{z} = \hat{q}\hat{h} + \hat{r}$, where $\deg(\hat{r}) < \deg(\hat{h}) = s$ or $\hat{r} = 0$. If $\hat{r} \neq 0$, then $\hat{r} = \hat{z} - \hat{q}\hat{h}$, i.e. we found $q(t) = \hat{q}(u^{-1}t^n), r(t) = \hat{r}(u^{-1}t^n) \in F[u^{-1}t^n]$, such that $r(t) = z(t) - q(t)h(t) \in Rf$. Let $\hat{r}'(x) = r_0^{-1}\hat{r}(x) \in F[x]$, where $r_0 \in F^{\times}$ is the leading coefficient of $\hat{r}(x)$, then $r'(t) = \hat{r}'(u^{-1}t^n)$ is monic by definition.

As $r'(t) = \hat{r'}(u^{-1}t^n) \in Rf$, too, there exists $a(t) \in R$ such that r'(t) = a(t)f(t). Thus, $r'(t) \in F[u^{-1}t^n]$ is a monic polynomial of degree less than *s* which is right divisible by *f*. This contradicts the definition of *h* as the minimal central left multiple of *f*. Thus, we conclude that r = 0 and $z = qh \in Rh$, as required. (ii) E_f is a commutative associative ring with identity 1 + Rf. Define the map $G: E_f \to E_{\hat{h}}$, G(z+Rf) = z + Rh for all $z \in F[u^{-1}t^n]$. *G* is well-defined and surjective. For all $x, y \in F[u^{-1}t^n]$, we have G(x+Rf) + G(y+Rf) = (x+Rh) + (y+Rh) = (x+y) + Rh = G(x+y+Rf), G(1+Rf) = 1 + P(x) *Rh*, and G(x + Rf)G(y + Rf) = (x + Rh)(y + Rh) = xy + Rh = G(xy + Rf), yielding that *G* is an isomorphism. To check injectivity, we note that G(x + Rf) = 0 + Rh if and only if $x \in Rh$. By Lemma 4 (i), this implies $x \in Rf$ and so x + Rf = 0 + Rf.

Let $B = \text{Nuc}_r(S_f)$ and k be the number of irreducible factors of h(t) in R.

Lemma 5. The left R-module R/Rf is a right B-module of rank k via the scalar multiplication $R/Rf \times B \longrightarrow R/Rf$, (a + Rf)(z + Rf) = az + Rf for all $z \in F[u^{-1}t^n]$ and $a \in R$. We can identify R/Rf with B^k via a canonical basis.

Proof. Since the Petit algebra $S_f = R/Rf$ with its multiplication $ab = ab \mod_r f$ is a nonassociative unital algebra with right nucleus B, R/Rf is a right B-module via the given scalar multiplication. As R/Rf is a vector space of dimension d^2mn over F, R/Rf is free of rank k over B.

Let $\nu \in D^{\times}$ and $\rho \in Aut(D)$, and define $F' = Fix(\rho) \cap F$. We assume in the following that F/F' is finite-dimensional. Let *s* be the degree of *B* over $E_{\hat{h}}$. We assume *f* is not right-invariant, i.e. k > 1.

Let l < k = dn/s be a positive integer. Define the set $S_{n,m,l}(v, \rho, f) = \{a + Rh \mid a \in P\} \subset R/Rh$, where

$$P = \{d_0 + d_1t + \dots + d_{lm-1}t^{lm-1} + \nu\rho(d_0)t^{lm} \mid d_i \in D\} \subset D[t;\sigma].$$

 $S_{n,m,l}(\nu, \rho, f)$ is a vector space over F' of dimension $d^2nml[F:F']$. R/Rf is a right *B*-module of rank *k*, as shown above. Let $L_a: R/Rf \to R/Rf$ be the left multiplication map $L_a(b+Rf) = ab + Rf$. Then, L_a is *B*-linear, as we have $a(x\alpha) = (ax)\alpha$ for all $\alpha \in B$, $a, x \in R/Rf$, and therefore, $L_a(x\alpha) = L_a(x)\alpha$ for all $\alpha \in B$. Thus, $L_a \in \text{End}_B(R/Rf)$ and

$$R/Rh \cong M_k(B) \cong \operatorname{End}_B(B^k) = \operatorname{End}_B(R/Rf)$$

by Theorem 3. Hence, we have well-defined maps

$$L: S_{n,m,l}(\nu, \rho, f) \to \operatorname{End}_{\mathcal{B}}(R/Rf), a \mapsto L_a,$$

$$\lambda: S_{n,m,l}(\nu, \rho, f) \to M_k(B), a \mapsto L_a \mapsto M_a,$$

where M_a is the matrix representing L_a with respect to a *B*-basis of *R*/*Rf*. We denote the image of $S_{n,m,l}(\nu, \rho, h)$ in $M_k(B)$ by

$$\mathcal{C}_{n,m,l} = \{ M_a \mid a \in S_{n,m,l}(\nu, \rho, f) \}.$$

The code $C = C_{n,m,l}$ is F'-linear by construction, and a generalized rank metric code. If C has minimum distance d_C , the Singleton-like bound canonically generalizes to the bound

$$\dim_{F'}(\mathcal{C}) \le k(k - d_{\mathcal{C}} + 1)[B:F'],$$

with [B:F'] = s[F:F']. If $d_{\mathcal{C}} = k - l + 1$, then $\dim_{F'}(S_{n,m,l}(\nu, \rho, f)) = d^2nml/dms[B:F'] = d^2mnl$ [F:F'] = lk[B:F'] = lkdms[F:F']. Thus, if $d_{\mathcal{C}} = k - l + 1$, then \mathcal{C} attains this bound and \mathcal{C} is a maximum rank distance code in $M_k(B)$.

We will usually deal with the case that deg (h) = dmn, so that $B = E_{\hat{h}}$ is a field, s = 1, and $C_{n,m,l} \subset M_{dn}(E_f)$. Note that if l = 1 and $d_c = k$, this generalized Singleton-like bound is achieved trivially: we obtain examples of MRD codes in $M_k(B)$. This arises when we look at division algebras $S_{n,m,1}(\nu, \rho, f)$ and the matrices representing their left multiplication, cf. Remark 17 and Corollary 18.

4. The rank of the matrix that corresponds to the element a + Rh

Let $R = D[t; \sigma]$ be as in Section 3, and $f \in R$ be an irreducible monic polynomial of degree m > 1with minimal central left multiple h. Let $B = \text{Nuc}_r(S_f)$. We have $\text{deg}(\hat{h}) = km$ and $R/Rh \cong M_k(B)$ as $E_{\hat{h}}$ algebras by Theorem 3. Let $\Psi : R/Rh \to M_k(B)$, $\Psi(a + Rh) = M_a$, be this isomorphism. For $M_a \in M_k(B)$, consider the right *B*-linear map $L_{M_a}: M_k(B) \to M_k(B), L_{M_a}: X \mapsto M_a X$. Then, we obtain the following generalization of [34], Proposition 7 (which was only proved for *f* with coefficients in a finite field, i.e. for the special case that deg(*h*) = *nm* is maximal):

Theorem 6. Let deg(h) = km. Then $M_a \in M_k(B)$ and

$$\dim_{B}(\operatorname{im}(L_{M_{a}})) = k^{2} - \frac{k}{m} \operatorname{deg}(\operatorname{gcrd}(a, h)), \quad \operatorname{colrank}(M_{a}) = k - \frac{1}{m} \operatorname{deg}(\operatorname{gcrd}(a, h))$$

for all $a + Rh \in R/Rh$. In particular, if $\deg(h) = dmn$, then $M_a \in M_n(E_{\hat{h}})$, and

$$\operatorname{rank}(M_a) = dn - \frac{1}{m} \operatorname{deg}(\operatorname{gcrd}(a(t), h(t))).$$

Proof. For each $M_a \in M_k(B)$, define $\operatorname{Ann}_r(M_a) = \{N \in M_k(B) \mid M_a N = 0\}$. Then, $\operatorname{Ann}_r(M_a)$ is the kernel of the endomorphism $L_{M_a} : M_k(B) \to M_k(B)$. By the Rank-Nullity Theorem for free right *B*-modules of finite rank [16, ch. IV, Cor. 2.14], it follows that

$$k^2 = \dim_B(\operatorname{im}(L_{M_a})) + \dim_B(\operatorname{Ann}_r(M_a)).$$

We conclude that $\dim_B(im(L_{M_a})) = k^2 - \dim_B(\operatorname{Ann}_r(M_a))$. Now for each b + Rh, $M_aM_b = 0$ if and only if $\Psi(a + Rh)\Psi(b + Rh) = 0$. As Ψ is multiplicative, this is true if and only if $\Psi((a + Rh)(b + Rh)) = 0$. This means (a + Rh)(b + Rh) = 0. Hence, it is clear that $\operatorname{Ann}_r(M_a) \cong \operatorname{Ann}_r(a)$, where

$$Ann_{r}(a) = \{b + Rh \in R/Rh \mid (a + Rh)(b + Rh) = 0 + Rh\},\$$

so dim(Ann_r(M_a)) = dim(Ann_r(a)). Let $\gamma = gcrd(a, h)$ so $h = \delta\gamma$ for some $\delta \in R$. As $h \in C(R)$ and R is a domain, we also have $h = \gamma\delta$ by Lemma 1. Let $b \in R$ be the unique element such that $a = b\gamma$. Then, $gcrd(b, \delta) = 1$, else γ is not the greatest common right divisor of a and h.Let $v \in R$. By the left Euclidean division algorithm, there exist unique $u, w \in R$ such that $v = \delta u + w$ where $deg(w) < deg(\delta)$ and $gcld(w, \delta) = 1$. It follows that $av = a\delta u + aw = b\gamma\delta u + b\gamma w = bhu + b\gamma w$, and therefore, $av + Rh = b\gamma w + Rh$. Suppose $b\gamma w \equiv 0 \mod_r h$. As $gcrd(b, \delta) = 1$, there exist $c, d \in R$ such that $cb + d\delta = 1$, so $cb\gamma + d\delta\gamma = \gamma$. As $\delta\gamma = h$, this implies $cb\gamma \equiv \gamma : \mod_r h$. Hence, $\gamma w \equiv cb\gamma w \equiv 0 : \mod_r h$. However, $deg(w) < deg(\delta)$ so $deg(\gamma w) < deg(\gamma \delta) = deg(h)$; due to this, $\gamma w \equiv 0 : \mod_r h$ implies that $\gamma w = 0$. As $\gamma \neq 0$ and R is a domain, we conclude that w = 0.Hence, (a + Rh)(v + Rh) = 0 + Rh if and only if $v = \delta u$ where $deg(u) < deg(\gamma)$. As δ is uniquely defined by a and h, every element of Ann_r(a) is determined by $u \in R$ such that $deg(u) < deg(\gamma)$. Thus,

$$\operatorname{Ann}_{r}(a) = \{ v + Rh \in R/Rh \mid (a + Rh)(v + Rh) = 0 + Rh \}$$
$$= \{ \delta u \mid u \in R, \deg(u) < \deg(\gamma) \} \cong R_{\deg(\gamma)} = \{ g \in R \mid \deg(g) < \deg(\gamma) \}.$$

As $\{1, t, \dots, t^{\deg(\gamma)-1}\}$ is a *D*-basis for the free left *D*-module R_{γ} , it follows that $\dim_D(\operatorname{Ann}_r(a)) = \deg(\gamma)$, so $\dim_F(\operatorname{Ann}_r(a)) = \deg(\gamma)d^2n$. Since $\dim_{E_{\hat{h}}}(B) = s^2 = d^2n^2/k^2$ and $[E_{\hat{h}}:F] = km/n$, we obtain $\dim_F(B) = d^2mn/k$. Hence, we get

$$\dim_{B}(\operatorname{Ann}_{r}(a)) = \frac{\operatorname{deg}(\gamma)d^{2}nk}{d^{2}mn} = \frac{\operatorname{deg}(\gamma)k}{m},$$

and so

$$\dim_B(\operatorname{im}(L_A)) = k^2 - \dim_B(\operatorname{Ann}_r(M_a)) = k^2 - \frac{k}{m}\operatorname{deg}(\gamma).$$

Let $\underline{c_i}$, respectively $\underline{r_i}$, denote the columns, and rows of M_a and $\underline{x_i}$ denote the columns of X. Computing the matrix using dot product notation, we have

$$M_a X = \begin{pmatrix} \underline{r_1} \cdot \underline{x_1} & \dots & \underline{r_1} \cdot \underline{x_k} \\ \vdots & \ddots & \vdots \\ \underline{r_k} \cdot \underline{x_1} & \dots & \underline{r_k} \cdot \underline{x_k} \end{pmatrix}$$

The i^{th} column of $M_a X$ is equal to

$$\begin{pmatrix} \underline{r_1} \cdot \underline{x_i} \\ \vdots \\ \underline{r_k} \cdot \underline{x_i} \end{pmatrix} = \underline{c_1}\lambda_1 + \dots + \underline{c_k}\lambda_k$$

for some $\lambda_j \in B$. Hence, the dimension of the right *B*-module generated by the i^{th} column of $M_a X$ is exactly the column rank of M_a . As there are *k* columns of $M_a X$, it follows that $\dim_B(\operatorname{im}(L_{M_a})) = k \operatorname{colrank}(M_a)$.

All of the above applies in particular to the special case that K/F is a field extension and $\sigma \in \operatorname{Aut}_F(K)$ of finite order $n, R = K[t; \sigma]$ and $C(R) = F[t^n] \cong F[x]$. Let $f \in R$ be a monic irreducible polynomial of degree $m > 1, B = \operatorname{Nuc}_r(S_f)$, and $h(t) = \hat{h}(t^n)$ its minimal central left multiple, $\operatorname{deg}(\hat{h}) = km$. Then, $\Psi : R/Rh \to M_k(B), \Psi(a + Rh) = M_a$ is an E_f -algebra isomorphism. For each $M_a \in M_k(B)$, we have the endomorphism $L_{M_a} : M_k(B) \to M_k(B)$ by $L_{M_a} : X \mapsto M_a X$. Analogously to Theorem 6, we can prove:

Theorem 7 (for fnite fields and thus deg(h) = nm maximal, cf. [34], Proposition 7). Suppose that deg(h) = km, then for all $a + Rh \in R/Rh$ we have

$$\dim_B(\operatorname{im}(L_{M_a})) = k^2 - \frac{k}{m} \operatorname{deg}(\operatorname{gcrd}(a, h)), \quad \operatorname{colrank}(M_a) = k - \frac{1}{m} \operatorname{deg}(\operatorname{gcrd}(a, h)).$$

In particular, if deg(h) = mn then $M_a \in M_n(E_{\hat{h}})$ and rank $(M_a) = n - \frac{1}{m} \text{deg}(\text{gcrd}(a, h))$.

This generalizes [34, Remark 6].

5. Using the norm of $D(t; \sigma)$ to investigate f

5.1. The algebra $(D(x), \tilde{\sigma}, ux)$

Let C/F be a finite cyclic field extension of degree *n* with $\operatorname{Gal}(C/F) = \langle \sigma \rangle$. Let *D* be a finite-dimensional division algebra of degree *d* with center *C* and suppose that σ extends to a *C*-algebra automorphism of *D* that we call σ , too. Let $R = D[t; \sigma]$ as in Section 3. Then, there exists $u \in D^{\times}$ such that $\sigma^n = i_u$ and $\sigma(u) = u$. These two relations determine *u* up to multiplication with elements from F^{\times} [25, Lemma 19.7].

The quotient algebra $(D, \sigma, a) = D[t; \sigma]/(t^n - a)D[t; \sigma]$, where $f(t) = t^n - a \in D[t; \sigma]$ with $d \in F^{\times}$, is called a *generalized cyclic algebra*. The special case where D = C yields the cyclic algebra $(C/F, \gamma, a)$ [18, p. 19].

Let $D(t; \sigma) = \{f/g \mid f \in D[t; \sigma], g \in C(D[t; \sigma])\}$ be the ring of central quotients of $D[t; \sigma]$. Let $\tilde{\sigma}$ denote the extension of σ to D(x) that fixes x [14, Lemma 2.1]. Then, $C(D(t; \sigma)) = \text{Quot}(C(D[t; \sigma])) = F(x), x = u^{-1}t^n$, is the center of $D(t; \sigma)$, where Quot(U) denotes the quotient field of an integral domain U. More precisely, $D(t; \sigma) \cong (D(x), \tilde{\sigma}, ux)$ is a generalized cyclic algebra of degree dn over its center F(x) and a division algebra [14, Theorems 2.2, 2.3].

Let *N* be the reduced norm of $(D(x), \tilde{\sigma}, ux)$.

Lemma 8. Let $f \in R$. If N(f) is irreducible in F[x], then f is irreducible in R.

Proof. If f = gp for $g, p \in R$ then N(f) = N(g)N(p) is reducible in F[x], since both N(g) and N(p) lie in F[x], which immediately yields the assertion.

From now on, we assume that

 $D = (E/C, \gamma, a)$ is a cyclic division algebra over C of degree d,

 $\sigma|_E \in \operatorname{Aut}(E)$ such that $\gamma \circ \sigma = \sigma \circ \gamma$ and $u \in E$.

Then, $\sigma|_E$ has order *n*. Write m = kn + r for some $0 \le r < n$. Let $f = \sum_{i=0}^{m} a_i t^i \in R$ be a polynomial such that $a_0 \ne 0$ and $h \in R$ be the minimal central left multiple of *f* in *R*.

Theorem 9 [36]. For $f \in E[t; \sigma] \subset D[t; \sigma]$, we have

 $N(f(t)) = N_{E/F}(a_0) + \dots + (-1)^{dr(n-1)} N_{E/F}(a_m) N_{E/C}(u)^m x^{dm}.$

Theorem 10. Suppose that deg(h) = dmn.

- (i) [36, Theorem 14 (i)] If \hat{h} is irreducible in F[x], then f is irreducible in R.
- (ii) [36, Theorem 14 (ii)] If f is irreducible, then N(f) is irreducible in F[x].
- (*iii*) If $f \in E[t; \sigma]$, then $N(f) = (-1)^{dr(n-1)} N_{E/F}(a_m) N_{E/C}(u)^m \hat{h}$ and

$$N_{E/F}(a_0) = (-1)^{dr(n-1)} N_{E/F}(a_m) N_{E/C}(u)^m h_{0}$$

if h_0 denotes the constant term of \hat{h} .

Proof. (iii) By Theorem 9, we have $\deg(N(f)) = dmn$ in R. N(f) is a two-sided multiple of f in R; therefore, the bound f^* of f divides N(f) in R. Since $(f, t)_r = 1$, $f^* \in C(R)$ and therefore f^* equals h up to some factor in F^{\times} . Thus, $h(t) = \hat{h}(u^{-1}t^n)$ must divide N(f) in R. Write N(f) = g(t)h(t) for some $g \in R$. Comparing degrees in R, we obtain $\deg N(f) = \deg(g(t)) + dmn = dmn$, which implies $\deg(g) = 0$, i.e. $g(t) = a \in A^{\times}$. This implies that $N(f) = ah(t) = a\hat{h}(u^{-1}t^n)$. Comparing highest coefficients of N(f) and $a\hat{h}$ yields that $a = (-1)^{dr(n-1)}N_{E/F}(a_m)N_{E/C}(u)^m$ by Theorem 9, so that comparing constant terms we get that $N_{E/F}(a_0) = (-1)^{dr(n-1)}N_{E/F}(a_m)N_{E/C}(u)^m h_0$, if h_0 is the constant term of $\hat{h}(x)$.

Theorem 11. Let $f \in E[t; \sigma] \subset R$ be monic and irreducible of degree m. Let $\deg(\hat{h}) = dm$ and suppose that all the monic polynomials similar to f lie in $E[t; \sigma]$. If g is a monic divisor of h in R of degree lm, then

$$N_{E/F}(g_0) = N_{E/F}(a_0)^l$$
.

Proof. We know that $h(t) = \hat{h}(u^{-1}t^n)$, with $\hat{h}(x)$ irreducible in F[x], since f is irreducible. Thus, h is a t.s.m. element in Jacobson's terminology [18] and the irreducible factors $f_1(t), \ldots, f_k(t)$ of any decomposition of h(t) are all similar and are all similar to f, as f must be one of them by the definition of h. Now, g(t) is a monic divisor of h. Thus, we can decompose g(t) into a product of irreducible factors and up to similarity the irreducible factors of g will be the same as suitably chosen irreducible factors of h by [18, Theorem 1.2.9.]. Hence, w.l.o.g. $g = f_1 f_2 \cdots f_l$, where the f_i are irreducible in R and f_i is similar to f for all $i = 1, 2, \ldots, l$ [18, Theorem 1.2.19]. Thus by Lemma 2, the minimal central left multiple of each f_i is equal to h. Since f is monic, we may assume w.l.o.g. that all f_i are monic. By Theorem 10 and since all $f_i \in E[t; \sigma]$ by our assumption, this implies that $N_{E/F}(f_i(0)) = (-1)^{dm(n-1)} N_{E/C}(u)^m h_0 = N_{E/F}(a_0)$. As the constant term of g is equal to $\prod_{i=1}^{l} f_i(0)$, we see that

$$N_{E/F}(g_0) = \prod_{i=1}^{l} N_{E/F}(f_i(0)) = [(-1)^{dm(n-1)} N_{E/C}(u)^m h_0]^l$$
$$= (-1)^{ldm(n-1)} N_{E/C}(u)^{lm} h_0^l = N_{E/F}(a_0)^l.$$

We are not able to say if the assumptions on the f_i 's in the above result are empty or trivial.

5.2. The algebra $(K(x)/F(x), \tilde{\sigma}, x)$

Let K/F be a cyclic field extension of degree *n* with $Gal(K/F) = \langle \sigma \rangle$, $R = K[t; \sigma]$ and $x = t^n$. We now look at the cyclic algebra $(K(x)/F(x), \tilde{\sigma}, x)$ (this case corresponds to D = C in the previous Section).

Let *N* be the reduced norm of $(K(x)/F(x), \tilde{\sigma}, x)$ over F(x) (cf. also [18, Proposition 1.4.6]). We have $\tilde{\sigma}|_{K} = \sigma$, and *N* is a nondegenerate form of degree *n*. Let $f = \sum_{i=0}^{m} a_{i}t^{i} \in R$ be a polynomial of degree *m* such that $a_{0} \neq 0$ and $h \in R$ be the minimal central left multiple of *f* in *R*. Then $N(f(t)) = N_{K/F}(a_{0}) + \cdots + (-1)^{m(n-1)}N_{K/F}(a_{m})x^{m}$ [36, Theorem 3].

Theorem 12. Suppose that deg(h) = mn.

- (i) [36, Theorem 6 (i)] If \hat{h} is irreducible in F[x], then f is irreducible in R.
- (ii) [36, Theorem 6 (ii)] If f is irreducible, then N(f) is irreducible in F[x].
- (iii) $N_{K/F}(a_0) = (-1)^{m(n-1)}h_0$, if h_0 denotes the constant term of \hat{h} .

Theorem 12 (iii) is proved analogously as Theorem 10 (iii).

Theorem 13 (cf. [34, Theorem 5] for finite fields, the proof is the same). Suppose that f is not rightinvariant. If deg(h) = mn and g is a monic divisor of h(t) in R of degree ml, then

$$N_{K/F}(g_0) = N_{K/F}(a_0)^l.$$

6. Division algebras and MRD codes employing $f \in R$

6.1. The case that $f \in D[t; \sigma]$

Let $f \in R = D[t; \sigma]$ be a monic polynomial of degree *m*. Let $\rho \in Aut(D)$, $\nu \in D$ and $F' = Fix(\rho) \cap F$ where $F = C \cap Fix(\sigma)$. Let $b(t), c(t) \in R_m = \{g \in R \mid \deg(g) < m\}$ and b_0 be the constant term of b(t). Then, the multiplication defined via

$$b(t) \circ c(t) = (b(t) + \nu \rho(b_0)t^m)c(t) \operatorname{mod}_r f,$$

makes R_m into a non-unital nonassociative ring (R_m, \circ) . When the context is clear, we will drop the \circ notation and simply use juxtaposition. (R_m, \circ) is an algebra over F'.

Example 14. If
$$f(t) = t - c \in D[t; \sigma]$$
 for some $c \in D$, $v \neq 0$, then (R_m, \circ) has the multiplication
 $a \circ b = (a + v\rho(a)t)b) \mod_{z} f$
 $= ab + v\rho(a)\sigma(b)t \mod_{z} f$
 $= ab + v\rho(a)\sigma(b)c$,

for all $a, b \in D$. This generalizes the algebras studied in [30]. If $R = K[t; \sigma]$ for some finite field extension K/F; this is the multiplication of Albert's twisted semifields [1]. If F/F' is finite and (R_m, \circ) is not a division algebra, $a \circ b = 0$ for some nonzero $a, b \in D$, if and only if $ab = -v\rho(a)\sigma(b)c$. Taking norms of both sides and canceling $N_{D/F'}(ab)$ on both sides, we obtain that $N_{D/F'}(-vc) = (-1)^{d^2n[F:F']}N_{D/F'}(vc) = 1$. Thus, if F/F' is finite and $N_{D/F'}(vc) \neq (-1)^{d^2n[F:F']}$ then (R_m, \circ) is a division algebra.

From now on for the rest of the paper, we again assume that f is an irreducible monic polynomial of degree m > 1, $(f, t)_r = 1$, and that h is the minimal central left multiple of f. Let F/F' be finite-dimensional, and

$$P = \{d_0 + d_1t + \dots + d_{lm-1}t^{lm-1} + \nu\rho(d_0)t^{lm} \mid d_i \in D\} \subset D[t;\sigma].$$

Theorem 15. Let l = 1. Then:

(i) Let $b(t) \in R_m$ with constant coefficient b_0 . If $b(t) + v\rho(b_0)t^m \in P$ is reducible in R, then b(t) is not a left zero divisor in (R_m, \circ) .

- (ii) If v = 0, then (R_m, \circ) is a division algebra over F', which for $m \ge 2$ is a Petit algebra.
- (iii) If P does not contain any polynomial similar to f, then (R_m, \circ) is a division algebra over F'.

Note that f may be right-invariant.

Proof. Suppose that there are $b(t) = b_0 + b_1 t + \cdots + b_{m-1} t^{m-1}$, $c(t) \in R_m$, such that

 $b(t) \circ c(t) = (b(t) + \nu \rho(b_0)t^m)c(t) \operatorname{mod}_r f = 0.$

Then, there exists $g \in R_m$ such that $(b(t) + v\rho(b_0)t^m)c(t) = g(t)f(t)$. Since f is irreducible and of degree m, while deg(c) < m, f must be similar to an irreducible factor of $b(t) + v\rho(b_0)t^m$, because of the uniqueness of an irreducible decomposition in R up to similarity. But $b(t) + v\rho(b_0)t^m$ has degree at most m, so f is similar to $b(t) + v\rho(b_0)t^m$. Thus, $b(t) + v\rho(b_0)t^m$ must have degree m and be irreducible as well. Hence if $b(t) + v\rho(b_0)t^m$ is not similar to f then $b(t) + v\rho(b_0)t^m$ is not a left zero divisor in (R_m, \circ) . This happens for instance, if v = 0 or if $b(t) + v\rho(b_0)t^m$ is reducible. Moreover, (R_m, \circ) is a division algebra if P does not contain any polynomial similar to f.

We are again not able to say if the assumptions on the f_i 's in the following result are empty or trivial.

Theorem 16. Let $D = (E/C, \gamma, a)$ be a cyclic division algebra over C of degree d such that $\sigma|_E \in Aut(E)$ and $\gamma \circ \sigma|_E = \sigma|_E \circ \gamma$. Suppose that $\sigma^n(z) = u^{-1}zu$ with $u \in E$.

Let $f(t) = \sum_{i=0}^{m} a_i t^i \in E[t; \sigma] \subset D[t; \sigma]$ be monic and irreducible, and let $\deg(h) = dmn$. Suppose that all monic f_i similar to f lie in $E[t; \sigma]$. Then, (R_m, \circ) is a division algebra over F', if one of the following holds: (i) $v \notin E$ and $\rho|_E \in \operatorname{Aut}(E)$. (ii) $v \in E^{\times}$ and $\rho|_E \in \operatorname{Aut}(E)$, such that

$$N_{E/F'}(a_0)N_{E/F'}(\nu) \neq 1.$$

Note that our global assumption that $\sigma^n(z) = u^{-1}zu$ for all $z \in D$, so that $\sigma^n(e) = u^{-1}eu = e$ for all $e \in E$, forces $(\sigma|_E)^n = id$.

Proof. By Theorem 15, (R_m, \circ) is a division algebra, if the set *P* with l = 1 does not contain any polynomial similar to *f*. All polynomials similar to *f* are irreducible factors of h(t), so (R_m, \circ) is a division algebra, if *P* does not contain any irreducible factor of h(t). Suppose that *P* contains an irreducible factor *g* of *h* with constant term g_0 . Then, *g* has degree *m* as it is similar to *f*. Let $g_m t^m$ be its highest coefficient, so that $g_m^{-1}g$ is a monic divisor of *h*.

By Theorem 10 and since $g \in E[t; \sigma]$ by assumption, this implies

$$N_{E/F}(g_0g_m^{-1}) = (-1)^{m(n-1)}h_0 = N_{E/F}(a_0),$$

and in particular, that g_0 and g_m are both nonzero. Since $g \in P$, we also have $g_m = \nu \rho(g_0)$. Suppose $\nu \notin E$ and $\rho(E) \subset E$. Since the coefficients of the f_i all lie in E, we have $g_m \neq \nu \rho(g_0)$ which yields a contradiction. Hence, there is no divisor g of h in P and S is a division algebra. Suppose that $\nu \in E^{\times}$ and $\rho(E) \subset E$. Substituting $g_m = \nu \rho(g_0)$ into the above equation yields

$$N_{E/F}(g_0) = N_{E/F}(a_0)N_{E/F}(\nu\rho(g_0)).$$

Applying $N_{F/F'}$ to both sides implies that

$$N_{E/F'}(g_0) = N_{F/F'}(N_{E/F}(a_0))N_{E/F'}(\nu\rho(g_0)).$$

Now $N_{E/F'}(\rho(g_0)) = N_{E/F'}(g_0)$, so we can cancel the nonzero term $N_{E/F'}(g_0)$ to obtain $1 = N_{E/F'}(a_0)$ $N_{E/F'}(\nu)$.

Remark 17. Let $S = S_{n,m,1}(v, \rho, f) = \{a + Rh \mid a \in P\}$. We can use $C(S) \subset M_k(B)$ to define a multiplication on B^m . As $\dim_F(D) = d^2n$ and $\dim_F(B) = d^2mn/k$, there exists an F-vector space isomorphism between D^m and B^k . Similarly, there exists an isomorphism $G : V_f \to B^k$, $G(a + Rf) = \underline{a}$. Define $* : B^k \times B^k \to B^k$ by

$$\underline{a} * \underline{b} = M_a \cdot \underline{b}$$

for all $\underline{a}, \underline{b} \in B^k$, where $M_a \in C(S)$ is the representation of the map $L_{a(t)+\nu\rho(a_0)t^m} \in \operatorname{End}_B(R/Rf)$ induced by G. (Each $a \in R_m$ corresponds to a map $L_{a(t)+\nu\rho(a_0)t^m}$. As $\operatorname{End}_B(R/Rf) \cong M_k(B)$ and $\dim(R_m) = \dim(\mathcal{C}(S))$, there is a canonical bijection between $L_{a(t)+\nu\rho(a_0)t^m}$ and M_a .) As M_a represents $L_a \in \operatorname{End}_B(R/Rf)$, $(B^k, *)$ is isomorphic to R/Rf equipped with the multiplication $(a + Rf)(b + Rf) = L_{a(t)+\nu\rho(a_0)t^m}(b + Rf)$. Thus, (R_m, \circ) and $(B^k, *)$ are isomorphic algebras and $S_{n,m,1}(\nu, \rho, f)$ is the same algebra as (R_m, \circ) .

If l = 1, we write $S(\nu, \rho, f) = S_{n,m,1}(\nu, \rho, f)$ for (R_m, \circ) . $S(\nu, \rho, f)$ is a division algebra if and only if every matrix in $C_{n,m,1}$ has full column rank. It then canonically defines an *F*'-linear *MRD* code in $M_k(B)$, $B = \text{Nuc}_r(S_f)$. Therefore, we obtain from all of the above results:

Corollary 18. Let $D = (E/C, \gamma, a)$ be a cyclic division algebra over C of degree d such that $\sigma|_E \in Aut(E)$ and $\gamma \circ \sigma = \sigma \circ \gamma$. Suppose that $\sigma^n(z) = u^{-1}zu$ with $u \in E$.

Let $f = \sum_{i=0}^{m} a_i t^i \in R$ be monic and irreducible of degree m. Then, B is a division algebra over $E_{\hat{h}}$ and $S(v, \rho, f)$ defines an F'-linear MRD code in $M_k(B)$ with minimum distance k, if one of the following holds:

- (*i*) v = 0. Then, $S(v, \rho, f)$ is a (unital) Petit algebra.
- (ii) P does not contain any polynomial similar to f.
- (iii) Suppose $\rho|_E \in \operatorname{Aut}(E)$, $f = \sum_{i=0}^m a_i t^i \in E[t; \sigma] \subset R$, $\deg(h) = dmn$, all the monic polynomials similar to f lie in $E[t; \sigma]$, and one of the following holds:
 - (a) $v \notin E$,
 - (b) $N_{E/F'}(v)N_{E/F'}(a_0) \neq 1$. Then, we get an F'-linear MRD code in $M_{dn}(E_{\hat{h}})$ with minimum distance dn.

The case v = 0 produces the MRD codes which are associated with the unital Petit algebras. They can be viewed as generalized Gabidulin codes.

More generally, we can also construct MRD codes for l > 1. Let $f \in R$ not be right-invariant, and let l < k be a positive integer.

Theorem 19. Suppose that P does not contain any polynomial of degree lm, whose irreducible factors are all similar to f. Then, the set $S_{n,m,l}(v, \rho, f)$ defines an F'-linear MRD code in $M_k(B)$ with minimum distance k - l + 1. In particular, if deg(h) = dmn, then this code is an F'-linear MRD code in $M_{dn}(E_{\hat{h}})$ with minimum distance dn - l + 1.

We are not able to say if the assumption on P can be satisfied in this general setup. It is satisfied in the case considered in [34, Theorem 7].

Proof. We have to show that the minimum column rank of the matrix corresponding to a nonzero element in $S_{n,m,l}(v, \rho, f)$ is k - l + 1. By Theorem 6, this is equivalent to finding an element $g \in A$ such that the greatest common right divisor of g and h has degree at most (l - 1)m. Suppose towards a contradiction that there exists $g \in A$ such that deg(gcrd(g, h)) = lm; since deg $(g) \leq lm$, it follows that g must be a divisor of h. As any divisor of h is a product of irreducible polynomials similar to f, g must be a product of polynomials similar to f. This contradicts our assumption, so any matrix has rank at least k - l + 1.

Theorem 20 (for $f \in K[t; \sigma]$, K a finite field, this is [34, Theorem 7]). Let $f = \sum_{i=0}^{m} a_i t^i \in E[t; \sigma] \subset R = D[t; \sigma]$ be monic irreducible, and let deg(h) = dmn. Suppose that all monic f_i similar to f lie in $E[t; \sigma]$. Then, $S_{n,m,l}(v, \rho, f)$ defines an F'-linear MRD code in $M_{dn}(E_{\hat{h}})$ with minimum distance dn - l + 1, if one of the following holds:

(i) v = 0(ii) $v \notin E$ and $\rho|_E \in \operatorname{Aut}(E)$. (iii) $v \in E$, $\rho|_E \in \operatorname{Aut}(E)$ and $N_{E/F'}(v)N_{E/F'}(a_0)^l \neq 1$.

The proof is straightforward.

6.2. The case $R = K[t; \sigma]$

Let $f = \sum_{i=0}^{m} a_i t^i \in R = K[t; \sigma]$ be an irreducible monic polynomial of degree *m* with minimal central left multiple *h*. Suppose throughout this section that F/F' is a finite field extension, $v \in K$, and $\rho \in \operatorname{Aut}(K)$. Let 1 < l < k and $S_{n,m,l}(v, \rho, f) = \{a + Rh \mid a \in P\} \subset R/Rh$, where $P = \{d_0 + d_1t + \cdots + d_{lm-1}t^{lm-1} + v\rho(d_0)t^{lm} \mid d_i \in K\}$. Then, we obtain the following results:

Theorem 21. *Let* l = 1.

- (i) Let $b(t) \in R_m$ with constant coefficient b_0 . If $b(t) + v\rho(b_0)t^m \in P$ is reducible in R, then b(t) is not a left zero divisor in $S(v, \rho, f)$.
- (ii) If v = 0, then $S(v, \rho, f)$ is a division algebra over F', a unital Petit algebra.
- (iii) If P does not contain any polynomial similar to f, then $S(v, \rho, f)$ is a division algebra over F'.

The proof is analogous to the one of Theorem 15. Note that f may be right-invariant here. Using Theorems 10 and 13, we obtain (for finite fields, cf. [34], the proof is analogous):

Theorem 22. Suppose that deg(h) = mn. Then, $S(v, \rho, f)$ is a division algebra over F' if

$$N_{K/F'}(a_0)N_{K/F'}(\nu) \neq 1.$$

Corollary 23. $B = \text{Nuc}_r(S_f)$ is a division algebra and the left multiplication of the algebra $S(v, \rho, f)$ defines an *F'*-linear MRD code in $M_k(B)$ with minimum distance *k*, if one of the following holds:

- (*i*) v = 0.
- (ii) P does not contain any polynomial similar to f.
- (iii) deg(h) = mn and $v \in K$ such that $N_{K/F'}(v) \neq 1/N_{K/F'}(a_0)$. In this case, the algebra $S(v, \rho, f)$ defines an F'-linear MRD code in $M_n(E_{\hat{h}})$ with minimum distance n.

Note that the condition on *f* in (iii) is satisfied for all *f* if gcd(m, n) = 1 or if *n* is prime. We now look at the case that 1 < l < k and also assume that *f* is not right-invariant.

Theorem 24 (for finite fields, cf. [34, Theorem 7]). If deg(h) = mn, then the set $S_{n,m,l}(v, \rho, f)$ defines an F'-linear MRD code in $M_n(E_{\hat{h}})$ with minimum distance n - l + 1 for any $v \in K$ such that

$$N_{K/F'}(\nu) \neq 1/N_{K/F'}(a_0)^l$$
.

Note that k = n here since deg(h) = mn.

Corollary 25. The set $S_{n,m,l}(\nu, \rho, h)$ defines an F'-linear MRD code in $M_n(E_{\hat{h}})$ with minimum distance n - l + 1, if one of the following holds:

- (*i*) $\deg(h) = mn \text{ and } v = 0$,
- (*ii*) *n* is prime or gcd(m, n) = 1, and $1 \neq N_{K/F'}(v)N_{K/F'}(a_0)^l$
- (iii) $\deg(h) = mn \text{ and } N_{K/F'}(v) \notin (F'^{\times})^l$.

The codes $S_{n,m,l}(0, \rho, h)$ generalize the Gabidulin codes constructed in [10] that go back to [8]. Note that $N_{K/F'}(\nu) \notin (F'^{\times})^l$ implies $N_{K/F'}(\nu) \neq N_{K/F'}(a_0)^l$ for any f. Thus, (ii) implies (iii) above.

Theorem 26. Suppose that P does not contain any polynomial of degree lm, whose irreducible factors are all similar to f. Then the set $S_{n,m,l}(v, \rho, f)$ defines an F'-linear MRD code in $M_k(B)$ with minimum distance k - l + 1. In particular, if deg(h) = mn then $S_{n,m,l}(v, \rho, f)$ defines an F'-linear MRD code in $M_n(E_{\hat{h}})$ with minimum distance n - l + 1.

Proof. We have to show that the minimum column rank of the matrix corresponding to a nonzero element in $S_{n,m,l}(\nu, \rho, f)$ is k - l + 1. By Theorem 7, this is equivalent to finding an element $g \in P$ such that the greatest common right divisor of g and h has degree at most (l - 1)m. Suppose towards a contradiction that deg(gcrd(g, h)) = lm; since deg(g) $\leq lm$, it follows that g must be a divisor of h.

As any divisor of h is a product of irreducible polynomials similar to f, g must be a product of polynomials similar to f. This contradicts our assumption, so any matrix has rank at least k - l + 1.

7. Nuclei

Let $\mathcal{M} = \mathcal{M}(A) = \{L_a \mid a \in A\} \subseteq \operatorname{End}_F(A)$ be the spread set of an *F*-algebra *A*, where L_a is the left multiplication map in *A*. We define the *left* and *right idealizers* of \mathcal{M} as

$$I_{l}(\mathcal{M}) = \{ \Phi \in \operatorname{End}_{F}(A) \mid \Phi \mathcal{M} \subseteq \mathcal{M} \}, \text{ respectively, } I_{r}(\mathcal{M}) = \{ \Phi \in \operatorname{End}_{F}(A) \mid \mathcal{M} \Phi \subseteq \mathcal{M} \}.$$

The *centralizer* of \mathcal{M} is defined as $\text{Cent}(\mathcal{M}) = \{\Phi \in \text{End}_F(A) \mid \Phi M = M\Phi \quad \forall M \in \mathcal{M}\}$. We call $Z(\mathcal{M}) = I_l(\mathcal{M}) \cap \text{Cent}(\mathcal{M})$ the *center* of \mathcal{M} .

Theorem 27 (cf. [34, Proposition 5] for finite fields). Let A be a unital division algebra and \mathcal{M} be the spread set of A. Let \mathcal{M}^* be the spread set associated with the opposite algebra A^{op} . Then

 $\operatorname{Nuc}_{l}(A) \cong I_{l}(\mathcal{M}), \quad \operatorname{Nuc}_{m}(A) \cong I_{r}(\mathcal{M}), \quad \operatorname{Nuc}_{r}(A) \cong \operatorname{Cent}(\mathcal{M}^{*}), \quad C(A) \cong Z(\mathcal{M}).$

The proof from [34] holds verbatim in our general setting.

The above results can now be applied to determine the nuclei and center of the non-unital algebras $S = S_{n,m,l}(v, \rho, f)$.

In the following, let $R = D[t; \sigma]$. We use the assumptions on *D*, respectively *K*, and σ from Section 6.

Let $f \in R$ be an irreducible monic polynomial of degree *m*, and let *h* be the minimal central left multiple of *f*. We assume throughout that *f* is not right-invariant, so that k > 1.

Remark 28. The algebras $S_{n,m,l}(0, \rho, f)$ are unital Petit algebras and hence have left nucleus $Nuc_m(S) = D$, and their right nucleus $\{g \in R_m | fg \in Rf\}$ is the eigenspace of f. If $S_{n,m,l}(0, \rho, f)$ is not associative, then $\{d \in D | dg = gd \text{ for all } g \in S\}$ is their center [26].

Theorem 29. Let $R = D[t; \sigma]$ and $\deg(h) = dmn$. Suppose $l \le dn/2$, n > 1 and lm > 2. Let $S = S_{n,m,l}(v, \rho, f)$ and \mathcal{M} be the image of S in $\operatorname{End}_{E_f}(R/Rf)$, that means the corresponding rank metric code lies in $M_n(E_{\hat{h}})$. If $v \ne 0$, we have

(i) $I_l(\mathcal{M}) \cong \{g_0 \in D \mid g_0 v = v \rho(g_0)\} \subset D$ (in particular, $I_l(\mathcal{M}) \cong \operatorname{Fix}(\rho)$ if $v \in C$),

(*ii*) $I_r(\mathcal{M}) \cong \operatorname{Fix}(\rho^{-1} \circ \sigma^{lm}) \subset D$,

(*iii*) Cent(\mathcal{M}) $\cong E_{\hat{h}}$, $Z(\mathcal{M}) \cong F'$.

If v = 0, we have

- (*iv*) $I_l(\mathcal{M}) \cong D, I_r(\mathcal{M}) \cong D$,
- (v) $\operatorname{Cent}(\mathcal{M}) \cong E_{\hat{h}}, Z(\mathcal{M}) \cong F.$

Much of the proof works identically to the proof of [34, Theorem 9]. We sketch the proof to highlight the main differences in this more general case. The lm = 2 case has to be considered separately, and we have only been able to solve that for $F = \mathbb{R}$.

Proof. Let $\mathcal{M} = \{L_a \in \operatorname{End}_{E_f}(R/Rf) \mid a \in P\}$ be the image of *S* in $\operatorname{End}_{E_f}(R/Rf) \subset \operatorname{End}_F(R/Rf)$. In the following, we identify each element in \mathcal{M} with the element $g \in S$ that induces it. Analogously to the proof of [34, Theorem 9], $\{g \in I_l(\mathcal{M}) \mid \deg(g) \leq lm\} = \{g_0 \in D \mid g_0 v = v\rho(g_0)\}$. If v = 0, then $1 \in \mathcal{M}$ so $I_l(\mathcal{M}) \subset \mathcal{M}$ so all $g \in I_l(\mathcal{M})$ have degree at most *lm*. Consider $v \neq 0$. To check there are no elements $g \in I_l(\mathcal{M})$ of degree higher than *lm*, we follow the approach of [34, Theorem 9] and consider $gt \mod \hat{h}(u^{-1}t^n)$. Recalling $\deg(h) = dm$, we have $h(t) = (u^{-1}t^n)^{dm} + \cdots = u^{-dm}(t^n + h'_{dm-1}t^{(dm-1)n} + \cdots + h'_0)$ so

$$gt \bmod h(t) = \left(\sum_{i=0}^{dmn-1} g_{i-1}t^{i}\right) - g_{dmn-1}u^{dm}\left(\sum_{j=0}^{dm-1} h'_{j}t^{n_{j}}\right).$$

As $g \in I_l(\mathcal{M})$, this implies $gt \mod h \in \mathcal{M}$, so for all $i \in \{lm + 1, \dots, dmn - 1\}$, we have

$$g_{i-1} = \begin{cases} 0 & \text{for } i \neq 0 \mod n \\ g_{dmn-1} u^{dm} h'_{i/n} & \text{for } i \equiv 0 \mod n \end{cases}$$
(1)

where $h'_{i/n} = 0$ if i/n is not an integer. We will show that $g_{dmn-1} = 0$ and thus $\deg(g) \le lm - 1$. As lm > 2, this follows verbatim from [34, Theorem 9].

The same holds for $I_r(\mathcal{M})$ following Sheekey's proof with the appropriate amendments made for $D[t; \sigma]$. The results for Cent(\mathcal{M}) and $Z(\mathcal{M})$ hold verbatim from [34, Theorem 9].

Corollary 30. Let $R = D[t; \sigma]$ and $\deg(h) = dmn$. Suppose n > 1, m > 2 and $S = S_{n,m,1}(\nu, \rho, f)$ with $\nu \neq 0$ be a division algebra. Then,

- (*i*) $\operatorname{Nuc}_{l}(S) \cong \{g_{0} \in D \mid g_{0}\nu = \nu\rho(g_{0})\} \subset D$, so in particular $\operatorname{Nuc}_{l}(S) = \operatorname{Fix}(\rho) \subset D$, if $\nu \in C$.
- (*ii*) Nuc_m(S) \cong Fix($\rho^{-1} \circ \sigma^m$) $\subset D$.
- (*iii*) $C(S) = \operatorname{Fix}(\rho) \cap F = F'$.
- (*iv*) $\dim_{F'} \operatorname{Nuc}_{r}(S) = \dim_{F'}(E_{\hat{h}}) = \operatorname{deg}(\hat{h})[F:F'] = [F:F']dm.$

Theorem 31. Let $R = K[t; \sigma]$ and deg(h) = mn. Suppose $l \le n/2$, n > 1 and lm > 2. Let $S = S_{n,m,l}(\nu, \rho, f)$ with $\nu \ne 0$ and \mathcal{M} be the image of S in End_{Ef}(R/Rf), so that the corresponding rank metric code lies in $M_n(E_{\hat{h}})$. Then,

- (*i*) $I_l(\mathcal{M}) \cong \operatorname{Fix}(\rho) \subset K, I_r(\mathcal{M}) \cong \operatorname{Fix}(\rho^{-1} \circ \sigma^{lm}) \subset K$,
- (*ii*) Cent(\mathcal{M}) \cong $E_{\hat{h}}$, $Z(\mathcal{M}) \cong F'$. If v = 0, we have
- (*iii*) $I_l(\mathcal{M}) \cong K, I_r(\mathcal{M}) \cong K$,
- (*iv*) Cent(\mathcal{M}) $\cong E_{\hat{h}}$, $Z(\mathcal{M}) \cong F$.

Again, the proof is analogous to the one of [34], Theorem 9 (it does not use the fact that for finite fields the right nucleus of S_f is $E_{\hat{h}}$, it only uses that R/Rh has center $E_{\hat{h}}$).

Corollary 32. Let $R = K[t; \sigma]$ and deg(h) = mn. Suppose that n > 1, m > 2 and that $S = (S(v, \rho, f), \circ)$ is a division algebra with $v \neq 0$. Then,

- (*i*) $\operatorname{Nuc}_{l}(S) = \operatorname{Fix}(\rho) \subset K$,
- (*ii*) Nuc_m(S) = Fix($\rho^{-1} \circ \sigma^m$) $\subset K$,

(*iii*) $C(S) = \operatorname{Fix}(\rho) \cap F = F'$.

(*iv*) $\dim_{F'} \operatorname{Nuc}_{r}(S) = \dim_{F'}(E_{\hat{h}}) = \deg(\hat{h})[F:F'] = [F:F']m.$

Theorems 29, 31 and Corollaries 30, 32 generalize [34, Theorem 9, Corollary 1] which were proved for semifields.

8. Examples of division algebras and an MRD code when $f(t) = t^n - \theta \in K[t; \sigma]$

8.1. $K = F(\theta)$

Let $K = F(\theta)$ be an extension of prime degree *n*. Let $f(t) = t^n - \theta \in K[t; \sigma]$. We now compute the rank metric code associated with the *F'*-algebra $S_{n,n,1}(v, \rho, f)$. Note that $f(t) = t^3 - \theta \in K[t; \sigma]$ is irreducible if and only if $\theta \neq \sigma^2(z)\sigma(z)z$ for all $z \in K$. If *F* contains a primitive *n*th root of unity, then f(t) is irreducible if and only if $\theta \neq \sigma^{n-1}(z) \cdots \sigma(z)z$ for all $z \in K$.

We assume that *f* is irreducible. Define $h(t) = (t^n - \theta)(t^n - \sigma(\theta)) \cdots (t^n - \sigma^{n-1}(\theta)) = (t^n)^n + \cdots + (-1)^n N_{K/F}(\theta)$, then h(t) = mclm(f): as $t^n - \sigma^i(\theta) \in K[t^n]$, the factors of h(t) all commute and $h(t) \in K[t^n]$. Since $\sigma(h(t)) = (t^n - \sigma(\theta)) \cdots (t^n - \sigma^{n-1}(\theta))(t^n - \theta) = h(t)$, we know that $h(t) \in \text{Fix}(\sigma)[t] = F[t]$ so $h(t) \in F[t] \cap K[t^n] = F[t^n] = C(R)$. Hence $h(t) = \hat{h}(t^n)$ with $\hat{h}(x) = x^n + (-1)^n N_{K/F}(\theta) \in F[x]$. Thus *f* divides *h* both from the left and the right by Lemma 1.

As *n* is prime, the minimal central left multiple of *f* must have degree *n* in *F*[*x*] by Theorem 3; thus, $h(t) = \operatorname{mclm}(f)$, and hence, $\hat{h}(x) = x^n + (-1)^n N_{K/F}(\theta)$ also is an irreducible polynomial in *F*[*x*]. As a field, $E_f = \{z + Rf \mid z \in F[t^n]\}$ is generated by $\{1 + Rf, t^n + Rf, t^{2n} + Rf, \dots, t^{n(n-1)} + Rf\} = \{1 + Rf, \theta + Rf, \theta^2 + Rf, \dots, \theta^{n-1} + Rf\}$ over *F*. As *K* is generated by $\{1, \theta, \dots, \theta^{n-1}\}$, there is a canonical isomorphism $E_f \longrightarrow K, x + Rf \mapsto x$.

It is clear that $\{1 + Rf, t + Rf, \dots, t^{n-1} + Rf\}$ is an E_f -basis for R/Rf. Let $a = a_0 + a_1t + \dots + a_{n-1}t^{n-1} + \nu\rho(a_0)t^n \in S(\nu, \rho, h)$. In order to determine M_a , we consider how $L_{a_it^i}$ acts on the basis elements of R/Rf. As left multiplication is distributive, i.e. $L_{a+b}(x) = L_a(x) + L_b(x)$, it follows that $L_a = \sum_{i=0}^n L_{a_it^i}$, where $a_n = \nu\rho(a_0)$. For each *i*, we have:

$$L_{a_{i}t^{i}}(1+Rf) = a_{i}t^{i} + Rf = (t^{i} + Rf)(\sigma^{n-i}(a_{i}) + Rf)$$

$$L_{a_{i}t^{i}}(t+Rf) = a_{i}t^{i+1} + Rf = (t^{i+1} + Rf)(\sigma^{n-i-1}(a_{i}) + Rf)$$

$$\vdots \qquad \vdots$$

$$L_{a_{i}t^{i}}(t^{n-i} + Rf) = a_{i}t^{n} + Rf = a_{i}\theta + Rf = (1 + Rf)(a_{i}\theta + Rf)$$

$$L_{a_{i}t^{i}}(t^{n-i+1} + Rf) = a_{i}t^{n+1} + Rf = a_{i}x\theta + Rf = (t + Rf)(\sigma(a_{i})\theta + Rf)$$

Thus, the matrix representing $L_{a;t^i}$ is given by

$$M_{a_it^i} = \begin{pmatrix} 0 & 0 & \cdots & 0 & \sigma^{n-i}(a_i) & 0 & \cdots & 0 \\ 0 & 0 & \cdots & 0 & 0 & \sigma^{n-(i+1)}(a_i) & \cdots & 0 \\ \vdots & & \ddots & & & \ddots & \vdots \\ 0 & & & \ddots & & & & \sigma(a_m) \\ a_i\theta & & & \ddots & & & \sigma(a_m) \\ a_i\theta & & & \ddots & & & \sigma(a_m) \\ \vdots & & \ddots & & & & \ddots & \vdots \\ 0 & 0 & \cdots & \sigma^{n-(i-1)}(a_i)\theta & 0 & 0 & \cdots & 0 \end{pmatrix}.$$

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As $M_a = \sum_{i=0}^n M_{a_i i^i}$, we obtain $M_a = (m_{i,j})_{i,j}$ where

$$m_{i,j} = \begin{cases} \sigma^{n+1-i}(a_0) + \sigma^{n+1-i}(\nu \rho(a_0))\theta & \text{for } i = j, \\ \sigma^{n+1-i}(a_{i-j}) & \text{for } i > j, \\ \sigma^{n+1-i}(a_{n+i-j})\theta & \text{for } i < j. \end{cases}$$

This yields $C_{n,n,1} = \{M_a \mid a_k \in K \text{ for } k = 0, 1, \dots, n-1\} \subset M_n(K)$ as the matrix spread set of the $n^2[F:F']$ -dimensional F'-algebra $S_{n,n,1}(\nu, \rho, f)$. The algebra associated with this spread set is a division algebra if $N_{K/F'}(\theta)N_{K/F'}(\nu) \neq 1$ (Theorem 22). In that case, the spread set will be an MRD code. In particular, for $\nu = 0$ this condition is satisfied for any irreducible $f(t) = t^n - \theta$. This is the well known result that for irreducible f, the Petit algebra S_f is a division algebra and so are all its isotopes. For n > 2 and $\nu \neq 0$, Corollary 32 yields

 $\operatorname{Nuc}_{l}(S) = \operatorname{Nuc}_{m}(S) = \operatorname{Fix}(\rho) \subset K, \quad C(S) = F', \quad \dim_{F'}\operatorname{Nuc}_{r}(S) = [F:F']m.$

8.2. Real division algebras of dimension 4

Over a finite field *F*, all division algebras of dimension 4 over *F* which have *F* as their center and a nucleus of dimension 2 over *F* can be constructed as algebras $S_{n,m,1}(\nu, \rho, f)$ for suitable parameters [34]. Let us now look at some real division algebras we obtain with our construction. If $\nu = 0$, then any choice of an irreducible $f \in \mathbb{C}[t; -]$ will yield an algebra isotopic to a real Petit division algebra. If $\nu \neq 0$, any choice of irreducible $f \in \mathbb{C}[t; -]$ where $N_{\mathbb{C}/\mathbb{R}}(a_0) \neq 1/N_{\mathbb{C}/\mathbb{R}}(\nu)$ also yields a division algebra (Theorem 22).

Let $b \in \mathbb{R}$ and $f(t) = t^2 - bi \in \mathbb{C}[t; \neg]$. Then, $h(t) = \hat{h}(t^2)$, $\hat{h}(x) = x^2 + b^2 \in \mathbb{R}[x]$, is the minimal central left multiple of f, as $h(t) = t^4 + b^2 = (t^2 + bi)(t^2 - bi)$. For all b > 0, $f(t) = t^2 - bi$ is irreducible in $\mathbb{C}[t; \neg]$.

For every irreducible $f(t) = t^2 - bi$, and $v \in \mathbb{C}$ such that $N_{\mathbb{C}/\mathbb{R}}(v) \neq \frac{1}{b^2}$, we obtain a four-dimensional real division algebra $S_{2,2,1}(v, \rho, f)$ and an MRD code given by its matrix spread set

$$\mathcal{C}_{2,2,1} = \left\{ \begin{pmatrix} z_0 + \nu \rho(z_0)bi & z_1 bi \\ \overline{z_1} & \overline{z_0} + \overline{\nu \rho(z_0)}bi \end{pmatrix} \mid z_0, z_1 \in \mathbb{C} \right\},\$$

where ρ is either the identity or the complex conjugation.

As mentioned in Theorem 7, [34, Theorem 9] uses results to deal with the case when lm = 2 that are valid over finite fields, but can be extended to $R = \mathbb{C}[t; -]$, for instance for $f(t) = t^2 - i$:

Theorem 33. Let $f(t) = t^2 - i \in \mathbb{C}[t; \overline{}]$. Suppose $S = S_{2,2,1}(v, \rho, f)$ is a division algebra for some $v \neq 0$ and $\rho \in \operatorname{Aut}_{\mathbb{R}}(\mathbb{C})$. Then,

- (*i*) $\operatorname{Nuc}_{l}(S) = \operatorname{Nuc}_{m}(S) = \operatorname{Fix}(\rho)$,
- (*ii*) $C(S) = \mathbb{R}$,
- (*iii*) dim_{\mathbb{R}}(Nuc_{*r*}(*S*)) = dim_{\mathbb{R}}($\mathbb{R}[t^2]$) = 2.

Proof. We have $h(t) = t^4 + 1 \in \mathbb{R}[t^2]$. Suppose $g + Rh \in I_l(\mathcal{M})$ for some $g(t) = g_0 + g_1t + g_2t^2 + g_3t^3 \in \mathbb{R}$. Then, $ga \in S(v, \rho, h)$ for all $a \in S$. Direct and laborious computation yields $g_2 = 0$, $g_3 = -g_1\overline{v}$, and $v\rho(g_0a_0 + g_1\overline{va_1}) = g_0v\rho(a_0) + g_1\overline{a_1}$. This is satisfied for all $a_0, a_1 \in \mathbb{C}$ if and only if $v\rho(g_0) = g_0v$ and $g_1 = v\rho(g_1\overline{v})$.

If $\rho = id$, it follows that either $g_1 = 0$ or $N_{\mathbb{C}/\mathbb{R}}(\nu) = 1$; as *S* is a division algebra, Theorem 13 (or [34, Theorem 4]) forces $g_1 = 0$ and so $g = g_0$ for some $g_0 \in \mathbb{C}$. Thus $I_l(\mathcal{M}) = \mathbb{C}$. If $\rho = \overline{}$, then $g_1 = \nu^2 g_1$ so either $g_1 = 0$ or $\nu = \pm 1$. As $N_{\mathbb{C}/\mathbb{R}}(\nu) \neq 1$ [34, Theorem 4], this forces $g_1 = 0$ so $g = g_0$ for some $g_0 \in \mathbb{R}$. In this case, $I_l(\mathcal{M}) = \mathbb{R}$.

The computations for $I_r(\mathcal{M})$ follow analogously and $Cent(\mathcal{M})$ and $Z(\mathcal{M})$ follow from the proof of [34, Theorem 4]. We obtain the final result on the nuclei using Theorem 27 to relate the idealizers and centralizer of \mathcal{M} to the nuclei of the algebra *S*.

Example 34. If $f(t) = t^2 - i$, we obtain division algebras S for all $v \in \mathbb{C}$ such that $N_{\mathbb{C}/\mathbb{R}}(v) \neq 1$. If $v \neq 0$ and S is a division algebra, then $C(S) = \mathbb{R}$ and $\dim_{\mathbb{R}} \operatorname{Nuc}_r(S) = 2$. Since therefore $\operatorname{Nuc}_r(S)$ is a two-dimensional division algebra over \mathbb{R} , $\operatorname{Nuc}_r(S)$ is an Albert isotope of \mathbb{C} and can be found in the classification in [15, Theorem 1]: it must be \mathbb{C} , $\mathbb{C}^{(-,-)}$, $\mathbb{C}^{(1+L(v)-,-)}$, or $\mathbb{C}^{(1+L(v)-,1+L(w)-)}$, with $u, v \in \mathbb{C}$ suitably chosen.

If additionally $\rho = id$, then $\operatorname{Nuc}_{l}(S) = \operatorname{Nuc}_{m}(S) = \mathbb{C}$, and if $\rho = -$ then $\operatorname{Nuc}_{l}(S) = \operatorname{Nuc}_{m}(S) = \mathbb{R}$. Note that the four-dimensional algebras in the first class are all isotopes of nonassociative quaternion algebras.

9. Constructing algebras and codes using irreducible $f \in R = D[t; \delta]$

We now consider the same construction using differential polynomial rings. Let *C* a field of characteristic *p* and *D* be a finite-dimensional division algebra with center *C*. Let $R = D[t; \delta]$, where δ is a derivation of *D*, such that $\delta|_C$ is algebraic with minimum polynomial $g(t) = t^{p^e} + c_1 t^{p^{e-1}} + \cdots + c_e t \in F[t]$ of degree p^e , with $F = C \cap \text{Const}(\delta)$. (This includes the special case where d = 1, that is $R = K[t; \delta]$, and δ is an algebraic derivation with minimum polynomial *g*.) Then, $g(\delta) = id_{d_0}$ is an inner derivation of *D*. W.l.o.g. we choose $d_0 \in F$, so that $\delta(d_0) = 0$ [18, Lemma 1.5.3]. Then, $C(D[t; \delta]) = F[x] = \{\sum_{i=0}^k a_i(g(t) - d_0)^i | a_i \in F\}$ with $x = g(t) - d_0$. The two-sided $f \in R$ are of the form f(t) = uc(t) with $u \in D$ and $c(t) \in C(R)$ [18, Theorem 1.1.32]. All polynomials $f \in R$ are bounded.

For every $f \in R$, the minimal central left multiple of f in R is the unique polynomial of minimal degree $h \in C(R) = F[x]$ such that h = gf for some $g \in R$, and such that $h(t) = \hat{h}(g(t) - d_0)$ for some monic $\hat{h}(x) \in F[x]$. The bound f^* of f is the unique minimal central left multiple of f up to some scalar.

From now on, let $f \in R = D[t; \delta]$ be a monic irreducible polynomial of degree m > 1 and let $h(t) = \hat{h}(g(t) - d_0)$ be its minimal central left multiple. Then, $\hat{h}(x)$ is irreducible in F[x] and h generates a maximal two-sided ideal Rh in R [18, p. 16]. We have

$$C(R/Rh) \cong F[x]/F[x]\hat{h}(x)$$

[17, Proposition 4], and $\deg(h) = p^e deg(\hat{h})$. Define $E_{\hat{h}} = F[x]/F[x]\hat{h}(x)$ and let k be the number of irreducible factors of h in R.

Theorem 35 [23]. Nuc_r(S_f) is a central division algebra over $E_{\hat{h}}$ of degree $s = dp^e/k$, and

$$R/Rh \cong M_k(\operatorname{Nuc}_r(S_f)).$$

In particular, this means that $\deg(\hat{h}) = \frac{dm}{s}$, $\deg(h) = km = \frac{dp^e m}{s}$, and

$$[\operatorname{Nuc}_r(S_f):F] = s^2 \cdot \frac{dm}{s} = dms.$$

Moreover, s divides $gcd(dm, dp^e)$ *. If f is not right-invariant, then* k > 1 *and* $s \neq dp^e$ *.*

The proof is analogous to the one of Theorem 3. In particular, $[S_f : F] = [S_f : C]p^e = d^2m \cdot p^e$. Comparing dimensions, we obtain again that $[S_f : Nuc_r(S_f)] = k$, and if *f* is not right-invariant, k > 1.

For each $z(t) = \hat{z}(g(t) - d_0) \in F[g(t) - d_0]$ with $\hat{z} \in F[x]$, we have $z \in Rf$ if and only if $z \in Rh$. Let

$$E_f = \{z(t) + Rf \mid z(t) = \hat{z}((g(t) - d_0)) \in F[(g(t) - d_0)]\} \subset R/Rf.$$

Together with the multiplication $(x + Rf) \circ (y + Rf) = (xy) + Rf$ for all $x, y \in F[(g(t) - d_0)]$, E_f is a field extension of F of degree deg (\hat{h}) isomorphic to $E_{\hat{h}}$. Let $B = \text{Nuc}_r(S_f)$, then B has degree s over $E_{\hat{h}}$, and

R/Rf is a free right *B*-module of dimension k via $R/Rf \times B \longrightarrow R/Rf$, (a + Rf)(z + Rf) = az + Rf. We assume f is not right-invariant which yields k > 1.

For $\rho \in \operatorname{Aut}(D)$ define $F' = \operatorname{Fix}(\rho) \cap F$. We assume that F/F' is finite-dimensional. Let $\nu \in D$ and $1 \le l < k = dp^e/s$. Define the set $S_{p^e,m,l}(\nu, \rho, f) = \{a + Rh \mid a \in P\} \subset R/Rh$, where

$$P = \{a_0 + a_1t + \dots + a_{lm-1}t^{lm-1} + \nu\rho(a_0)t^{lm} \mid a_i \in D\} \subset D[t;\delta].$$

 $S_{p^e,m,l}(\nu, \rho, f)$ is a vector space over F' of dimension $d^2p^em[F:F']$. We identify each element of $S_{p^e,m,l}(\nu, \rho, f)$ with a map in $\operatorname{End}_B(R/Rf)$ as follows: For each $a \in S_{p^e,m,l}(\nu, \rho, f)$, let $L_a: R/Rf \to R/Rf$ be the left multiplication map $L_a(b+Rf) = ab+Rf$. Let M_a be the matrix in $M_k(B)$ representing L_a with respect to a *B*-basis of R/Rf and denote the image of $S = S_{p^e,m,l}(\nu, \rho, f)$ in $M_k(B)$ by

$$\mathcal{C}_{p^e,m,l} = \{M_a \mid a \in S_{p^e,m,l}(\nu,\rho,f)\}.$$

For l = 1, this construction again yields algebras over F': define a multiplication on the F'-vector space $R_m = \{g \in R \mid deg(g) < m\}$ via

$$a(t) \circ b(t) = (a(t) + \nu \rho(a_0)t^m)b(t) \operatorname{mod}_r(f).$$

For m > 1, (R_m, \circ) is isomorphic to $S(\nu, \rho, f) = S_{\rho^e, m, 1}(\nu, \rho, f)$. Therefore, we also denote (R_m, \circ) by $S(\nu, \rho, f) = S_{\rho^e, m, 1}(\nu, \rho, f)$.

Example 36. Let $R = D[t; \delta]$ and f(t) = t + c for some $c \in D$. For $v \in D^{\times}$ and $\rho \in Aut(D)$, $S_{p^e, 1, 1}(v, \rho, f) = (D, \circ)$ has the multiplication

$$x \circ y = (x + \nu\rho(x)t)y) : \operatorname{mod}_{x}f = xy + \nu\rho(x)yt + \nu\rho(x)\delta(y) : \operatorname{mod}_{r}f$$
$$= xy + \nu\rho(x)(\delta(y) - yc)$$

for all $x, y \in D$.

Theorem 37. Let $f \in D[t; \delta]$ be irreducible and $\deg(h) = km$. For all $a + Rh \in R/Rh$,

$$\dim_{B}(\operatorname{im}(L_{M_{a}})) = k^{2} - \frac{k}{m} \operatorname{deg}(\operatorname{gcrd}(a, \hat{h}(g(t) - d_{0}))),$$
$$\operatorname{colrank}(M_{a}) = k - \frac{1}{m} \operatorname{deg}(\operatorname{gcrd}(a, \hat{h}(g(t) - d_{0})).$$

In particular, if $\deg(h) = dmp^e$ then $M_a \in M_{p^e}(E_{\hat{h}})$ and

$$\operatorname{rank}(M_a) = dp^e - \frac{1}{m} \operatorname{deg}(\operatorname{gcrd}(a, \hat{h}(g(t) - d_0))).$$

Thus, $S_{p^e,m,1}(\nu, \rho, f)$ is a division algebra if and only if there are no divisors of *h* in $S_{p^e,m,1}(\nu, \rho, f)$. More generally for l > 1, the above result means:

Theorem 38. Suppose that P does not contain any polynomial of degree lm, whose irreducible factors are all similar to f. Then, the set $S_{p^e,m,l}(v, \rho, f)$ defines an F'-linear MRD code in $M_k(B)$ with minimum distance k - l + 1. In particular, if $\deg(h) = dmp^e$, then this code is an F'-linear MRD code in $M_{dp^e}(E_{\hat{h}})$ with minimum distance $dp^e - l + 1$.

Corollary 39. Suppose that l = 1.

- (i) If $a(t) + v\rho(a_0)t^m \in P$ is reducible, then a(t) is not a left zero divisor of (R_m, \circ) .
- (ii) If v = 0 then (R_m, \circ) is a division algebra over F', which for $m \ge 2$ is a Petit algebra.
- (iii) If P does not contain any polynomial similar to f, then (R_m, \circ) is a division algebra over F'.

The proofs are all identical to their analogues where $f \in D[t; \sigma]$.

Remark 40. One can also use the reduced norm N of the central simple algebra $D(t; \delta)$ in this setting: given a central simple algebra D with a maximal subfield E and $R = D[t; \delta]$, take the ring of central quotients $D(t; \delta) = \{f/g \mid f \in R, g \in C(R)\}$ of R. It has center $C(D(t; \delta)) = \text{Quot}(C(R)) = F(x)$, where x = $g(t) = d_0$. Let $\tilde{\delta}$ be the extension of δ to D(x) such that $\tilde{\delta} = id_{t|D(x)}$. Then, $D(t; \delta)$ is a central simple F(x)algebra, more precisely $D(t; \delta) \cong (D(x), \tilde{\delta}, d_0 + x)$, i.e. $D(t; \delta)$ is a generalized differential algebra.

Let N be the reduced norm of $D(t; \delta)$. For all $f \in R$, $N(f) \in F[x]$ and f divides N(f). Let $\omega : D \to M_d(E)$ be the left regular representation of D. For any $f \in R$ of degree m, $N(f) = \pm \det(\omega(a_m))^{p^e} x^{dm} + \dots$ In particular, N(f) has degree dm [36]. As the bound of f has degree dm in F[x], it follows that N(f) is equal to the bound of f. Thus if $\deg(\hat{h}) = dm$, we conclude that $\hat{h}(x) = \pm N(f)$.

There is more work to be done, for example, to determine the constant term of N(f(t)). This may lead to criteria on how to obtain division algebras using our construction. Additionally, the nuclei of both the algebras and the codes need to be calculated. For instance, consider the special case where d = 1, i.e. $R = K[t, \delta]$ for some field extension K/F. If $f(t) = a_0 + a_1t + \cdots + a_mt^m \in R = K[t; \delta]$ has degree m, then $N(f(t)) = (-1)^{m(p^e-1)}a_m^{p^e}x^m + \ldots$ [36, Thm 18(ii)] Thus,

$$N(f(t)) = a_m^{p^e} x^m + \dots$$

To find the constant term of N(f) is difficult. It is possible to compute special cases though, for example, for $f(t) = g(t) + a \in K[t; \delta]$, $N(f(t)) = (x + a)^{p^e}$ [36].

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