

## ON THE DIVISIBILITY OF HOMOGENEOUS DIRECTED GRAPHS

M. EL-ZAHAR AND N. W. SAUER

**ABSTRACT** Let  $\mathcal{T}$  be a finite set of finite tournaments. We will give a necessary and sufficient condition for the  $\mathcal{T}$ -free homogeneous directed graph  $H_{\mathcal{T}}$  to be divisible. That is, that there is a partition of  $H_{\mathcal{T}}$  into two classes such that neither of them contains an isomorphic copy of  $H_{\mathcal{T}}$ .

**Introduction.** Let  $T$  be a finite tournament. A *decomposition* of a tournament  $T$  is a partition of the vertex set into three classes  $L, M, N$  such that  $L \neq \emptyset$  and there are directed edges from  $x$  to  $y$  and from  $z$  to  $x$  whenever  $x \in L, y \in M$  and  $z \in N$ ; we write  $T = (L, M, N)$  to indicate such a decomposition. We assume that  $1 \leq |L| \leq |T|$ . If a finite tournament  $S$  together with a partition of  $S$  into two classes  $A$  and  $B$  is given we will indicate this by saying that  $[A, B]$  is a *partitioned tournament*. Two partitioned tournaments  $[A_0, B_0]$  and  $[A_1, B_1]$  are isomorphic if there exists a tournament isomorphism from  $[A_0, B_0]$  onto  $[A_1, B_1]$  which also preserves the partition. If  $(L, M, N)$  is a decomposition of a finite tournament, then  $p(L, M, N)$  is the partitioned tournament  $[M, N]$ .

Let  $\mathcal{T}$  be a finite set of finite tournaments, such that no element of  $\mathcal{T}$  is a subtournament of an other tournament of  $\mathcal{T}$ . If  $[A, B]$  is a partitioned tournament then:

$\mathcal{T}_{[A,B]} = \{L: \text{there is a decomposition } (L, M, N) \text{ of some element of } \mathcal{T} \text{ such that } p(L, M, N) \text{ is isomorphic to } [A, B]\}$ .

The set of minimal elements of  $\mathcal{T}_{[A,B]}$  is called a *derived set* of  $\mathcal{T}$ . Clearly, there are only finitely many derived sets of  $\mathcal{T}$  and every derived set of  $\mathcal{T}$  is finite. If  $\mathcal{R}$  and  $\mathcal{S}$  are two sets of tournaments, we write  $\mathcal{R} \prec \mathcal{S}$  iff for every  $R \in \mathcal{R}$  there is an  $S \in \mathcal{S}$  such that  $S$  is a subtournament of  $R$ . Observe that the relation  $\prec$  is transitive.

$H_{\mathcal{T}}$  is the countable homogeneous directed graph which embeds every finite  $\mathcal{T}$ -free directed graph [4]. We say  $H_{\mathcal{T}}$  is *indivisible*, if for every partition of  $H_{\mathcal{T}}$  into two classes, one of the classes contains an isomorphic copy of  $H_{\mathcal{T}}$ . The main result of this paper is the following theorem:

**THEOREM.** *If  $\mathcal{T}$  is a finite set of finite tournaments, then  $H_{\mathcal{T}}$  is indivisible if and only if the set of derived sets of  $\mathcal{T}$  is totally ordered under the relation  $\prec$ .*

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1. **Preliminaries.** A directed graph, or digraph,  $D$  is a set of vertices  $V(D)$  together with a set of directed edges or arcs  $\vec{E}(D) \subset V(D) \times V(D)$ . The arc  $(a, b)$  is directed from  $a$  to  $b$  and is denoted by  $\vec{ab}$ .  $D$  contains no loops and at most one of  $\vec{ab}, \vec{ba}$  is an arc of  $D$ . A *tournament* is a digraph in which every pair of vertices is linked by an arc. A digraph  $G$  is a *subdigraph* of  $D$  if  $V(G) \subset V(D)$  and  $\vec{E}(G) = (G \times G) \cap \vec{E}(D)$ . We say that  $D$  *embeds*  $G$  if it contains an induced subdigraph isomorphic to  $G$ . This will be denoted by  $G \rightarrow D$ .

If  $G \not\rightarrow D$  then we say that  $D$  *omits*  $G$  or is  $G$ -free. If  $\mathcal{A}$  is a set of digraphs we say  $D$  is  $\mathcal{A}$ -free if  $D$  is  $A$ -free for every  $A \in \mathcal{A}$ . A countable digraph  $D$  is *homogeneous* if every isomorphism (local automorphism)  $\alpha: A \rightarrow B$  between finite subdigraphs  $A, B$  of  $D$  extends to an automorphism of  $D$ ; (see [4], page 313). Lachlan [7] has classified the countable homogeneous tournaments, and Schmerl has classified the countable homogeneous partially ordered sets [11], and independently, Henson [5] and Peretyiatkin [9] have shown the existence of  $2^{\aleph_0}$  nonisomorphic homogeneous digraphs. Recently, Cherlin classified all homogeneous directed graphs [1].

A digraph  $D$  is *indivisible* if for every partition  $V(D) = R \cup B$  there is an isomorphism  $f: D \rightarrow D$  such that either  $f(D) \subset R$  or  $f(D) \subset B$ . (Excellent references for this and related concepts are [4] and [10]). The problem of classifying all countable divisible homogeneous undirected graphs has been completely solved [2] and [6]. In [8] all countable homogeneous undirected graphs have been classified.

Let  $D$  be a countable homogeneous digraph. The *age* of  $D$ , denoted by  $\mathcal{A}(D)$ , is the set of all finite digraphs (up to isomorphism) that can be embedded in  $D$ . It is well known that  $\mathcal{A}(D)$  has the following *amalgamation property*:

(AP) if  $A_0, A_1, A_2 \in \mathcal{A}(D)$  and  $f_i: A_0 \rightarrow A_i$  ( $i = 1, 2$ ) are embeddings then there is a digraph  $A \in \mathcal{A}(D)$  and embeddings  $g_i: A_i \rightarrow A$  ( $i = 1, 2$ ) such that  $g_1 \circ f_1 = g_2 \circ f_2$ .

In fact a countable class  $\Sigma$  of finite digraphs is the age of some countable homogeneous digraph if and only if  $\Sigma$  satisfies AP and is closed under taking subdigraphs. Furthermore, two homogeneous digraphs with the same age are isomorphic. Another characterizing property of homogeneous digraphs is the following *embedding property*:

(EP) if  $A \in \mathcal{A}(D)$  and  $x \in A$  then every embedding  $f: A - x \rightarrow D$  extends to an embedding  $g: A \rightarrow D$ .

$\mathcal{A}(D)$  is called *indivisible* if for every partition of  $D$  into two classes,  $R$  and  $B$ , one of the classes embeds every member of  $\mathcal{A}(D)$ . It is a standard compactness argument to show that indivisibility of  $\mathcal{A}(D)$  is equivalent to the following Ramsey property [10]. For every  $A_1, \dots, A_n \in \mathcal{A}(D)$  there exists  $B \in \mathcal{A}(D)$  such that for every partition  $B = B_1 \cup \dots \cup B_n$  there is an  $i$  such that  $A_i \rightarrow B_i$  [10]. Folkman, [3], has proven that the set of  $K_n$ -free undirected graphs is indivisible.

We shall consider the following strengthening of AP called the *free amalgamation property*:

(FAP) if  $A_0, A_1, A_2 \in \mathcal{A}(D)$  and  $f_i: A_0 \rightarrow A_i$  ( $i = 1, 2$ ) are embeddings then there is  $A \in \mathcal{A}(D)$  and embeddings  $g_i: A_i \rightarrow A$  ( $i = 1, 2$ ) such that the following are satisfied:

- (a)  $g_1 \circ f_1 = g_2 \circ f_2$
- (b) if  $x_1 \in g_1(A_1 - f_1(A_0))$  and  $x_2 \in g_2(A_2 - f_2(A_0))$  then  $x_1 \neq x_2$  and none of  $\overrightarrow{x_1x_2}, \overrightarrow{x_2x_1}$  is an arc of  $A$ .

REMARK. We shall say that  $g_1(A_1) \cup g_2(A_2)$  is obtained from  $A_1$  and  $A_2$  by free amalgamation over  $f_1(A_0) \approx f_2(A_0)$ .

A countable homogeneous digraph whose age has the FAP property will be called a *freely amalgamated homogeneous* (FAP) digraph. Necessarily, every FAP digraph is infinite.

Let  $\vec{\mathcal{T}}(D) = \{T : T \text{ is a tournament and } T \in \mathcal{A}(D)\}$ . Obviously  $\vec{\mathcal{T}}(D)$  is hereditary, *i.e.* is closed under forming subtournaments.

LEMMA 1. *Let  $\Sigma$  be a class of tournaments closed under forming subtournaments. Then there is a unique FAP digraph  $D$  such that  $\vec{\mathcal{T}}(D) = \Sigma$ .*

PROOF. Let  $\Sigma^*$  be the class of finite digraphs  $A$  satisfying: Every tournament embeddable into  $A$  belongs to  $\Sigma$ . Clearly  $\Sigma^*$  is closed under taking subgraphs and has FAP since the free amalgamation creates no new tournaments. Hence  $\Sigma^*$  is the age of a unique FAP digraph  $D$  which also satisfies  $\vec{\mathcal{T}}(D) = \Sigma$ . The uniqueness of  $D$  follows from the fact that every FAP digraph  $D$  for which  $\Sigma \subset \mathcal{A}(D)$  must satisfy  $\Sigma^* \subset \mathcal{A}(D)$ . ■

The previous lemma asserts that a FAP digraph is characterised by the set of tournaments which it embeds. This can be re-stated in terms of a set of finite forbidden tournaments as follows.

Let  $\mathcal{T}(D) = \{T : T \text{ is a tournament minimal w.r.t. being non-embeddable in } D\}$ . Then no member of  $\mathcal{T}(D)$  is embeddable into any other member and for every class of tournaments  $\Sigma$  with this property there exists a unique FAP digraph  $D$  such that  $\Sigma = \mathcal{T}(D)$ . Or, in other words, if  $\mathcal{T}$  is a set of finite tournaments then there exists exactly one FAP digraph  $H_{\mathcal{T}}$  which is  $\mathcal{T}$ -free. The age of  $H_{\mathcal{T}}$  is the set of all finite  $\mathcal{T}$ -free digraphs. Due to the uniqueness property of homogeneous structures [4],  $H_{\mathcal{T}}$  can also be characterised as the unique countable homogeneous digraph whose age consists of all finite  $\mathcal{T}$ -free digraphs.

FAP digraphs have many useful properties which are not true in general for homogeneous digraphs. One such property which the reader can easily verify is that the extension  $g$  in the aforementioned embedding property EP can be chosen in infinitely many ways. Other properties are stated below.

LEMMA 2. *Let  $D$  be a homogeneous digraph with the FAP and  $X$  a finite subset of  $D$ . Then the subdigraph induced by*

$$N(X) = \{y \in D - X : \forall x \in X (\overrightarrow{xy} \notin \vec{E}(D) \wedge \overrightarrow{yx} \notin E(D))\} \text{ is isomorphic to } D.$$

PROOF. Let  $A \in \mathcal{A}(D)$  and  $z \in A$ . We show that every embedding  $f: A - z \rightarrow N(X)$  extends to an embedding of  $A$  into  $N(X)$ . Let  $G$  be the digraph consisting of the disjoint

union of  $X$  and  $A$  with no further arcs.  $G$  is obtained by freely amalgamating  $A$  and  $X$  over the empty set. Therefore  $G \in \mathcal{A}(D)$ . Let  $i: X \rightarrow D$  be the inclusion map, then  $f \cup i$  embeds  $G - z$  into  $D$  and therefore can be extended to an embedding  $g: G \rightarrow D$ . The restriction  $g|_A$  is the required extension of  $f$ . ■

DEFINITION. A homogeneous digraph  $D$  is called *weakly indivisible* if it satisfies the following:

(WI) for every  $A \in \mathcal{A}(D)$  and every  $X \subset D$ , if  $A \not\rightarrow X$  then  $D \rightarrow D - X$ .

LEMMA 3. Let  $D$  be a FAP digraph. Then  $D$  is weakly indivisible.

PROOF. Let  $A \in \mathcal{A}(D)$  and  $X \subset D$  be such that  $A \not\rightarrow X$ . We use induction on the cardinality of  $A$  assuming that  $|A| \geq 2$ . Let  $B \in \mathcal{A}(D)$ ,  $z \in B$  and  $f: B - z \rightarrow D - X$  be an embedding. Let  $x \in A$ . By the induction hypothesis, the statement (WI) holds for  $A - x$ . Therefore, by Lemma 2, we can assume that there exists an embedding  $g: A - x \rightarrow X \cap N(f(B - z))$ . Consider the digraph  $G$  obtained from  $A$  and  $B$  by free amalgamation over  $\{x\} \approx \{z\}$ . Let  $x_0$  be the image of both  $x$  and  $z$  in  $G$ . the map  $f \cup g$  is an embedding of  $G - x_0$  into  $D$  and therefore it extends to an embedding  $h: G \rightarrow D$ . Clearly,  $h(x_0) \in D - X$ . This proves that  $f$  can be extended to an embedding of  $f': B \rightarrow D - X$  by letting  $f'(z) = h(x_0)$ . ■

Clearly the weak indivisibility of a homogeneous digraph  $D$  implies that  $\mathcal{A}(D)$  is indivisible. So we have the following theorem which extends Folkman’s theorem [3] to the case of digraphs.

THEOREM 4. Let  $\Sigma$  be a class of tournaments. Then for every set of finite  $\Sigma$ -free digraphs  $G_1, \dots, G_n$  there exists a finite  $\Sigma$ -free digraph  $H$  such that for every partition of  $H$  into  $H_1 \cup \dots \cup H_n$ , there exists an  $i \leq n$  and an embedding  $f_i: G_i \rightarrow H_i$ .

2. **The indivisibility of FAP digraphs.** Let  $\mathcal{T}$  be a set of finite tournaments and put  $D = H_{\mathcal{T}}$ . We shall assume that  $D$  is defined on  $N$  the set of positive integers. For each vertex  $x \in D$  let

$$\begin{aligned} \Gamma'(x) &= \{y : y < x \text{ and } \overrightarrow{xy} \in \vec{E}(D)\}, \\ \Gamma''(x) &= \{y : y < x \text{ and } \overleftarrow{yx} \in \vec{E}(D)\}, \\ \Gamma(x) &= (\Gamma'(x), \Gamma''(x)). \end{aligned}$$

For  $A, B \subset N$  we write  $A < B$  if  $\max A < \min B$ . For  $m \in N$ , we denote by  $[m] := \{k \in \omega : k \leq m\}$ . Let  $\mathcal{F}$  denote the class

$$\mathcal{F} = \{(A, B) : A, B \text{ are finite subsets of } \omega \text{ and } A \cap B = \emptyset\}.$$

For  $(A, B), (A', B') \in \mathcal{F}$  we write  $(A, B) \subset (A', B')$  if  $A \subset A'$  and  $B \subset B'$ . For a set  $A$  of positive integers, we also write  $A$  to denote the digraph induced by  $D$  on  $A$ . An embedding  $f: (A, B) \rightarrow (A', B')$  will always mean an embedding  $f: A \cup B \rightarrow A' \cup B'$  which sends  $A$  into  $A'$  and  $B$  into  $B'$ .  $(A, B), (A', B')$  are isomorphic,  $(A, B) \approx (A', B')$ , if

there is an embedding  $f : (A, B) \rightarrow (A', B')$  such that the restrictions of  $f$  to  $A$  and  $B$  are graph-isomorphisms

For  $(A, B) \in \mathcal{F}$  we define

$$C(A, B) = \{x \in \omega \mid \forall a \in A \forall b \in B (\overrightarrow{ax} \in \vec{E}(D) \wedge \overrightarrow{bx} \in \vec{E}(D))\}$$

Note that  $C(A, B) \cap (A \cup B) = \emptyset$

**LEMMA 5** For all  $(A, B) \in \mathcal{F}$ ,  $C(A, B)$  is an FAP digraph

**PROOF** If  $(A, B) = (\emptyset, \emptyset)$  then  $C(A, B) = D$  by definition. Let  $(A, B) \neq (\emptyset, \emptyset)$ . Let  $M, N \subset C(A, B)$  and  $f : M \rightarrow N$  be an isomorphism. Let  $\iota$  be the identity map on  $A \cup B$ . Then  $f \cup \iota : A \cup B \cup M \rightarrow A \cup B \cup N$  is an isomorphism. Then there is an automorphism  $\sigma : D \rightarrow D$  which extends  $f \cup \iota$ . Since  $\sigma$  fixes both  $A$  and  $B$ , it must map  $C(A, B)$  onto itself. The restriction  $\sigma \upharpoonright C(A, B)$  extends  $f$  to an automorphism of  $C(A, B)$ . This means that  $C(A, B)$  is a homogeneous digraph. To show that  $C(A, B)$  is an FAP digraph, let  $M, N, N', L$  be finite subgraphs of  $C(A, B)$  where  $N, N'$  are isomorphic and  $N$  is a subgraph of  $M$  and  $N'$  is a subgraph of  $L$ . To amalgamate  $M$  and  $L$  freely over  $N \approx N'$  we simply amalgamate  $M \cup A \cup B$  and  $L \cup A \cup B$  freely over  $A \cup B \cup N \approx A \cup B \cup N'$  and then discard the elements of  $A \cup B$ . ■

We define a preorder on  $\mathcal{F}$  by letting  $(A, B) \prec (A', B')$  if there exists an embedding  $C(A', B') \rightarrow C(A, B)$ . Obviously  $(A, B) \subset (A', B')$  implies that  $(A, B) \succ (A', B')$ . We now state

**THEOREM 6** If  $D$  is indivisible then  $\prec$  is a total preorder on  $\mathcal{F}$

**PROOF** Assume for a contradiction that  $(A, B), (A', B') \in \mathcal{F}$  are not  $\prec$  comparable. Then  $(A, B), (A', B') \in \mathcal{F}$  and  $C(A, B), C(A', B')$  are both non-empty FAP digraphs by Lemma 5. For any pair of pairs  $(E, F), (E', F') \in \mathcal{F}$  we write  $(E, F) <_1 (E', F')$  if  $\max(E \cup F) \Delta (E' \cup F') \in E' \cup F'$ . Here  $\Delta$  is the symmetric difference operator. Observe that  $<_1$  imposes a total order on the elements of  $\mathcal{F}$ . We shall define a partition  $D = D_1 \cup D_2 \cup D_3$  as follows. Let  $x \in D$ , then

- (a)  $x \in D_1$  if there exists  $(E, F) \subset \Gamma(x)$  such that  $(E, F) \approx (A, B)$  and  $(E, F) <_1 (E', F')$  for every  $(E', F') \subset \Gamma(x)$  satisfying  $(E', F') \approx (A', B')$
- (b)  $x \in D_2$  if there exists  $(E', F') \subset \Gamma(x)$  such that  $(E', F') \approx (A', B')$  and  $(E', F') <_1 (E, F)$  for every  $(E, F) \subset \Gamma(x)$  with  $(E, F) \approx (A, B)$
- (c)  $x \in D_3$  otherwise

We shall show that none of  $D_1, D_2$  and  $D_3$  embeds  $D$ . First we observe that every embedding  $\sigma : D \rightarrow D_i$  can be assumed to be orderpreserving. The reason for this is that every isomorphism  $f : C - x \rightarrow D$ , where  $x \in C \in \mathcal{A}(D)$ , can be extended to an isomorphism  $g : C \rightarrow D$  in infinitely many ways. This implies that for every embedding  $\sigma : D \rightarrow D_i$ , we can define another order-preserving embedding  $\sigma_1 : D \rightarrow \sigma(D)$  which we might as well consider instead of  $\sigma$ .

(a)  $D_1$  contains no isomorphic copy of  $D$ . Assume that  $\sigma : D \rightarrow D_1$  is an embedding. Let  $y \in C(A', B')$  be such that  $y \succ \max(A' \cup B')$  and put  $z = \sigma(y)$ . Then  $z \in C(M, N)$

where  $(M, N) = \sigma(A', B') = (\sigma(A'), \sigma(B'))$ . Since  $z \in D_1$  there must exist  $(E, F) \subset \Gamma(z)$  such that  $(E, F) \approx (A, B)$  and  $(E, F) <_1 (M, N)$ . There are only finitely many such  $(E, F)$  which we enumerate by  $(E_1, F_1), \dots, (E_k, F_k)$ . This defines a partition  $C(M, N) = C_1 \cup \dots \cup C_k$  where  $y \in C_j$  if  $\min\{i : y \in C(E_i, F_i)\} = j$ .

Since the age of  $C(M, N)$  is indivisible, Lemma 3 and Lemma 5, there is a class, say  $C_j$  which embeds every element of  $\mathcal{A}(C(M', N')) = \mathcal{A}(C(A, B))$ . Therefore  $\mathcal{A}(C(A', B')) \subset \mathcal{A}(C(A, B))$ . This implies that  $C(A, B) \rightarrow C(A', B')$  by the homogeneity of  $C(A', B')$ . Hence we arrived at a contradiction to our assumptions. The proof that  $D_2$  contains no isomorphic copy of  $D$  is similar.

(b)  $D_3$  contains no isomorphic copy of  $D$ . This is easy to see since  $D_3$  does not contain a vertex  $y$  such that  $(A, B) \rightarrow \Gamma(y) \cap D_3$ . ■

Before we proceed to discuss the sufficiency of the condition in Theorem 6, we investigate this condition further in terms of the set of tournaments forbidden in  $D$ .

We wish to describe  $C(A, B)$  in terms of  $(A, B)$ . From Lemma 5,  $C(A, B)$  is a FAP and therefore is characterized by its set of forbidden tournaments  $\mathcal{T}(C(A, B))$ . Let  $T \in \mathcal{T}(D)$  and assume that  $(K, M, N)$  is a decomposition of  $T$ . If  $(M, N) \rightarrow (A, B)$  then, clearly  $K \not\rightarrow C(A, B)$ . It is also true that every tournament  $P$  for which  $C(A, B)$  is  $P$ -free must arise in this way. Let  $L(A, B) = \{K : \exists T \in \mathcal{T}(D)$  with decomposition  $(K, M, N)$  such that  $(M, N) \rightarrow (A, B)$  is an embedding $\}$ . Then  $\mathcal{T}(C(A, B))$  is exactly the set of minimal (w.r.t. embedding) tournaments in  $L(A, B)$ .

LEMMA 7.  $(A, B) \prec (A', B')$  if and only if for each  $L \in L(A, B)$  there exists  $L' \in L(A', B')$  such that there is an embedding  $L' \rightarrow L$ .

PROOF. There is an embedding  $C(A', B') \rightarrow C(A, B)$  if and only if every tournament  $L$  in  $L(A, B)$  satisfies  $L \not\rightarrow C(A', B')$ , that is, there exists  $L' \in L(A', B')$  such that  $L' \rightarrow L$ . ■

Observe now that if the derived sets of  $\mathcal{T}$  are totally ordered, then the intersection of a set of derived sets is again a derived set. But this means that for  $(A, B) \in \mathcal{F}$ ,  $L(A, B)$  is a derived set of  $\mathcal{T}$ . Furthermore every derived set is equal to some  $L(A, B)$ . Hence we have observed: The derived sets of  $\mathcal{T}$  form a total order under  $\prec$  if and only if  $\prec$  is a total preorder of the pairs  $(A, B)$  of  $\mathcal{F}$ . If  $D = H_{\mathcal{T}}$  is indivisible, then the set of derived sets of  $\mathcal{T}$  form a total order under  $\prec$ . This together with Lemma 6 and Theorem 7 establishes the theorem stated in the introduction.

**3. The proof of the sufficiency of the condition.** Assume now that the relation  $\prec$  is a total preorder on  $\mathcal{F}$ . Let  $\sim$  denote the equivalence relation defined on  $\mathcal{F}$  by  $(A, B) \sim (A', B')$  if and only if  $C(A, B) \cong C(A', B')$ . Then  $(\mathcal{F} | \sim, \prec)$  is a linear order. We prove the converse of Theorem 6 under the assumption that this linear order is finite.

THEOREM 7. Let  $(\mathcal{F} | \sim, \prec)$  be a finite total order. Then  $D$  is indivisible.

PROOF. Assume that the vertices of  $D$  are colored red and blue. We must show that one of the two color classes contains an isomorphic copy of  $D$ . By the hypothesis there is an integer  $n \geq 1$  and a function  $\rho: \mathcal{F} \rightarrow \{1, 2, \dots, n\}$  such that  $\rho(A, B) < \rho(A', B')$  if

and only if  $(A, B) \prec (A', B')$  but  $(A', B') \not\prec (A, B)$ . We also assume that the range of  $\rho$  is an initial interval of  $\omega$  with length  $n$ . We call  $\rho$  the *rank function* for  $\mathcal{F}$ . Let  $(A, B) \in \mathcal{F}$  and let  $m = \max(A \cup B)$ . If  $\rho(A - \{m\}, B - \{m\}) < \rho(A, B)$  then  $(A, B)$  will be called *rank-critical*. Let

$$\mathcal{H} = \{(A, B, \alpha) : (A, B) \in \mathcal{F}, \alpha \in D, \max(A \cup B) < \alpha\}.$$

For  $E = (A, B, \alpha) \in \mathcal{H}$ , we let  $\Pi_1(E) = A, \Pi_2(E) = B$  and  $\Pi_3(E) = \alpha$ . We also define  $\rho(E) = \rho(A, B), \mu(E) = \min(A \cup B)$  and  $C(E) = \{y \in D : y > \alpha \text{ and } \Gamma'(y) \cap [\alpha] = A \text{ and } \Gamma''(y) \cap [\alpha] = B\}$ . For  $E' = (A', B', \alpha') \in \mathcal{H}$ , we define  $E \cup E' = (A \cup A', B \cup B', \max\{\alpha, \alpha'\})$ , and write  $E < E'$  whenever  $\alpha < \mu(E')$ . For each  $x \in D$ , let  $\varphi(x)$  denote the following formula:

$$\begin{aligned} &(\exists E_1 \in \mathcal{H})(x = \mu(E_1) \wedge \rho(E_1) = 1) \text{ such that} \\ &(\forall F_1 \in \mathcal{H})(E_1 < F_1 \wedge \rho(E_1 \cup F_1) = 1) \\ &\quad \vdots \\ &(\exists E_i \in \mathcal{H})(F_{i-1} < E_i \wedge \rho(E_1 \cup F_1 \cup E_2 \cup \dots \cup E_i) = i) \text{ such that} \\ &(\forall F_i \in \mathcal{H})(E_i < F_i \wedge \rho(E_1 \cup F_1 \cup E_2 \cup F_2 \cup \dots \cup E_i \cup F_i) = i) \\ &\quad \vdots \\ &(\exists E_n \in \mathcal{H})(F_{n-1} < E_n \wedge \rho(E_1 \cup F_1 \cup \dots \cup E_n) = n) \text{ such that} \\ &(\forall F_n \in \mathcal{H})(E_n < F_n \wedge \rho(E_1 \cup F_1 \cup \dots \cup E_n \cup F_n) = n) \end{aligned}$$

the set  $(C(E_1 \cup F_1 \cup \dots \cup F_n))$  contains infinitely many blue vertices.

Let  $\psi(x)$  denote the formula obtained from  $\varphi(x)$  by interchanging the quantifiers  $\exists, \forall$  and replacing the word ‘blue’ by ‘red’. It is clear that for each  $x \in D$  at least one of  $\varphi(x)$  and  $\psi(x)$  holds.  $x$  is called a *blue generator* if  $\varphi(x)$  holds and otherwise  $x$  will be called a *red generator*. The proof will be divided into two cases.

CASE 1. There are infinitely many blue generators.

We shall construct a sequence  $B_1 < \sigma(1) < B_2 < \sigma(2) < \dots$  of elements of  $D$  such that

- (1)  $B_1, B_2, \dots$  are blue generators;
- (2) the vertices  $\sigma(1), \sigma(2), \dots$  are colored blue;
- (3) the map  $\sigma: D \rightarrow D$  sending  $k$  into  $\sigma(k)$  is an embedding.

Certain sets will be ‘squeezed’ between the elements of the above sequence; these are exactly  $\Pi_1(E_i), \Pi_2(E_i)$  for the triples  $E_i$  obtained from the formulas  $\varphi(B_j)$ . We shall always use sets of the form  $\{\sigma(i_1), \dots, \sigma(i_\ell)\}$  as  $\Pi_1(F_i), \Pi_2(F_i)$  in these formulas. The elements  $\Pi_3(E_i), \Pi_3(F_i)$  will be called *constraints* and will be used to ensure that the constructed sets and elements are disjoint. Let us describe this construction in more detail.

Assume we have just chosen  $B_k$ . We then introduce  $\Pi_1(E_1(B_k)), \Pi_2(E_1(B_k))$  into our construction. We call  $\Pi_3(E_1(B_k))$  the *current constraint*. Any element or set to be subsequently included in the construction must be larger than the current constraint. Before we construct  $\sigma(k)$  we consider all pairs  $(C, H) \in \mathcal{F}$  satisfying

(4)  $\max(C \cup H) < k$ ;

(5) either  $(C \cup \{k\}, H)$  or  $(C, H \cup \{k\})$  is rank-critical and let  $t_{j+1}$  be  $k$ .

Let  $(C, H)$  be such a pair, say  $(C \cup \{k\}, H)$  is rank-critical. Let  $t_1 = \min(C \cup H)$  and denote by  $t_2 < \dots < t_j$  those vertices  $t \in C \cup H$  for which  $(C \cap [t], H \cap [t])$  is rank-critical. Put

(6)  $C_i = \{x \in C : t_i \leq x < t_{i+1}\}$  and

(7)  $H_i = \{x \in H : t_i \leq x < t_{i+1}\}$  for all  $i$  with  $1 \leq i \leq j$ .

We call  $t_1, \dots, t_j$  the *critical vertices* of  $(C, H)$  and  $(C_1, H_1), \dots, (C_j, H_j)$  the partition of  $(C, H)$  corresponding to them. The construction is such that for suitable  $\beta_i \in D$ , the triples  $F_i = (\sigma(C_i), \sigma(H_i), \beta_i)$ ,  $1 \leq i \leq j - 1$ , have been used in the formula  $\varphi(B_{t_i})$  to create triples  $E_i = E_i(B_{t_i}; F_1, \dots, F_{i-1})$ ,  $1 \leq i \leq j - 1$ . The sets  $\Pi_1(E_i), \Pi_2(E_i)$  were included in the construction such that they lie in the interval between  $B_{t_i}$  and  $\sigma(t_i)$ . Now choose  $\beta_j$  larger than the current constraint. Put  $F_j = (\sigma(C_j); \sigma(H_j), \beta_j)$ . From the formula  $\varphi(B_{t_i})$  we obtain  $E_{j+1} = E_{j+1}(B_{t_i}, F_1, \dots, F_j)$ . We then include  $\Pi_1(E_{j+1}), \Pi_2(E_{j+1})$  in the construction and take  $\Pi_3(E_{j+1})$  as the new constraint. In order to be able to apply  $\varphi(B_{t_i})$  we have assumed that

(8)  $\rho(E_1 \cup F_1 \cup \dots \cup E_i \cup F_i) = \rho(E_1 \cup F_1 \cup \dots \cup E_i)$ ,  $1 \leq i \leq j$ ;

(9)  $\rho(E_1 \cup F_1 \cup \dots \cup F_i \cup E_{i+1}) = \rho(E_1 \cup F_1 \cup \dots \cup E_i) + 1$ ,  $1 \leq i \leq j - 1$ .

We shall show later that condition (8) always holds. However, if (9) is not satisfied then some modification is needed. To demonstrate this, let us assume that  $\rho(C \cup \{k\}, H) = \rho(C, H) + r$  where  $r \geq 2$ . The triple  $E_{j+1}$  above satisfies

$$\rho(E_1 \cup F_2 \cup \dots \cup F_j \cup E_{j+1}) = \rho(E_1 \cup F_1 \cup \dots \cup F_j) + 1 = \rho(C, D) + 1.$$

The second equality follows from (8). We want to replace  $E_{j+1}$  by a triple  $E_{j+1}^*$  which satisfies

$$\rho(E_1 \cup F_1 \cup \dots \cup F_j \cup E_{j+1}^*) = \rho(C \cup \{k\}, H).$$

We recursively choose arbitrary triples  $F'_{j+1}, \dots, F'_{j+r-1}$  and apply  $\varphi(B_{t_i})$  to get  $E'_{j+2}, \dots, E'_{j+r}$  where

$$\begin{aligned} E_{j+1} &< F'_{j+1} < E'_{j+2} < \dots < F'_{j+r-1} < E'_{j+r}, \\ \rho(E_1 \cup F_1 \cup \dots \cup E_{j+1} \cup F'_{j+1} \cup E'_{j+2} \cup \dots \cup E'_{j+\ell} \cup F'_{j+\ell}) \\ &= \rho(E_1 \cup F_1 \cup \dots \cup E_{j+1} \cup F'_{j+1} \cup E'_{j+2} \cup \dots \cup E'_{j+\ell}). \end{aligned}$$

That is, after having found  $E'_{j+\ell}$  we choose  $F'_{j+\ell}$  so that  $\Pi_1(F'_{j+\ell}) = \Pi_2(F'_{j+\ell}) = \emptyset$  and  $\Pi_3(F'_{j+\ell}) > \Pi_3(E'_{j+\ell})$ . Then the existential quantifier in line 2( $j + \ell$ ) of the formula  $\varphi(B_{t_i})$  will produce  $E'_{j+\ell+1}$ . We then put  $E_{j+1}^* = E_{j+1} \cup F'_{j+1} \cup \dots \cup F'_{j+r-1} \cup E'_{j+r}$ . We include  $\Pi_1(E_{j+1}^*), \Pi_2(E_{j+1}^*)$  in our construction and consider  $\Pi_3(E_{j+1}^*)$  as the new constraint. We repeat this procedure for each  $(C, H)$  satisfying (4) and (5). We then proceed to construct  $\sigma(k)$ .



If  $\Gamma(k) = (\emptyset, \emptyset)$  then we choose a blue vertex  $y \in D$  such that  $(\Gamma'(y) \cap [\beta], \Gamma''(y) \cap [\beta]) = (\emptyset, \emptyset)$  where  $\beta$  is some element in  $D$  larger than the current constraint. The existence of such a  $y$  follows from Lemma 2. Then we put  $\sigma(k) = y$ .

Assume now that  $\Gamma(k) = (C, H) \neq (\emptyset, \emptyset)$ . Let  $t_1 < \dots < t_j$  be the critical points of  $(C, H)$  and  $(C_i, H_i)$ ,  $1 \leq i \leq j$ , be the corresponding partition of  $(C, H)$ . It follows from the above construction that the triples  $F_i = (\sigma(C_i), \sigma(H_i), \beta_i)$ ,  $1 \leq i \leq j - 1$ , have been used in  $\varphi(B_{t_i})$  to induce triples  $E_i = E_i(B_{t_i}, F_1, \dots, F_{i-1})$ ,  $1 \leq i \leq j$ . Assume that

$$\rho(E_1 \cup F_1 \cup \dots \cup E_j \cup F_j) = \rho(E_1 \cup F_1 \cup \dots \cup E_j).$$

Then, choose an element  $\beta_j \in D$  larger than the current constraint. Put  $F_j = (C_j, H_j, \beta_j)$ . The last line of  $\varphi(B_{t_i})$  says we can choose a blue element  $y \in D$ ,  $y > \beta_j$  such that  $y \in C(E_1 \cup F_1 \cup \dots \cup F_j)$ . Observe that this implies that

$$(\Gamma'(y) \cap \{\sigma(1), \dots, \sigma(k-1)\}, \Gamma''(y) \cap \{\sigma(1), \dots, \sigma(k-1)\}) = (\sigma(C), \sigma(H)).$$

Then we put  $\sigma(k) = y$  and continue the construction by choosing a blue generator  $B_{k+1} > \sigma(k)$ . Thus we have to prove that

$$\rho(E_1 \cup F_1 \cup \dots \cup E_i \cup F_i) = \rho(E_1 \cup F_1 \cup \dots \cup E_i), \quad 1 \leq i \leq j,$$

or, equivalently,  $C(E_1 \cup F_1 \cup \dots \cup E_i)$  can be embedded into  $C(E_1 \cup F_1 \cup \dots \cup E_i \cup F_i)$ . We prove this by induction on  $i$ . Arguing by contradiction, we assume that there is a decomposable tournament  $(M, N) \in \mathcal{T}(D)$  such that  $L \rightarrow C(E_1 \cup F_1 \cup \dots \cup E_i)$  but  $L \not\rightarrow C(E_1 \cup F_1 \cup \dots \cup E_i \cup F_i)$ , and such that

$$(M, N) \not\rightarrow (\Pi_1(E_1 \cup F_1 \cup \dots \cup E_i), \Pi_2(E_1 \cup F_1 \cup \dots \cup E_i)),$$

and

$$(M, N) \rightarrow (\Pi_1(E_1 \cup F_1 \cup \dots \cup E_i \cup F_i), \Pi_2(E_1 \cup F_1 \cup \dots \cup F_i)).$$

Let  $\{x_1, \dots, x_n\}$  be the image of  $(M, N)$  under this last embedding where  $x_1 < \dots < x_n$ . Thus we have  $x_n \in \Pi_1(F_i) \cup \Pi_2(F_i)$ . By induction,

$$\rho(E_1 \cup F_1 \cup \dots \cup E_i) = \rho(C_1 \cup \dots \cup C_i, H_1 \cup \dots \cup H_i),$$

which implies that  $\{x_1, \dots, x_n\} \not\subseteq \sigma(C_1 \cup \dots \cup C_i \cup H_1 \cup \dots \cup H_i)$ . It follows that some  $x_h$  belongs to

$$\Pi_1(E_1 \cup E_2 \cup \dots \cup E_i) \cup \Pi_2(E_1 \cup E_2 \cup \dots \cup E_i).$$

Let  $j$  be the maximum of such  $h$  and assume say, that  $x_j \in \Pi_1(E_\ell) \cup \Pi_2(E_\ell)$ . Since  $(M, N)$  is a tournament, each of  $x_{j+1}, \dots, x_n$  is connected by an arc to  $x_j$ . According to the above description of the construction we have  $\overrightarrow{x_h x_j}, \overrightarrow{y_h x_j} \in \vec{E}(D)$  for every  $j < h \leq n$  and every

$$\begin{aligned} x &\in \Pi_1(E_1 \cup F_1 \cup \dots \cup F_{\ell-1} \cup E_\ell), \\ y &\in \Pi_2(E_1 \cup F_1 \cup \dots \cup F_{\ell-1} \cup E_\ell). \end{aligned}$$

It also follows that each of  $x_{j+1}, \dots, x_n$  is connected by an arc to  $\sigma(t_\ell)$  where the direction of this arc depends only on  $\sigma(t_\ell)$ . We note that  $\sigma(t_\ell)$  is different from  $x_1, \dots, x_n$  since there is no arc between  $\sigma(t_\ell)$  and  $E_\ell$ . (An arc from  $\sigma(w)$  to  $E_\ell$  means there is an arc from  $w$  to  $t_\ell$ .) Let us assume that  $t_\ell \in C_\ell$  (the case  $t_\ell \in H_\ell$  is similar). Thus we have

$$\{x_{j+1}, \dots, x_n\} \subset C(\sigma(C_1 \cup \dots \cup C_{\ell-1} \cup \{t_\ell\}), \sigma(H_1 \cup \dots \cup H_{\ell-1})),$$

and

$$\{x_{j+1}, \dots, x_n\} \subset C(E_1 \cup F_1 \cup \dots \cup E_\ell).$$

By induction, we have

$$\begin{aligned} \rho(E_1 \cup F_1 \cup \dots \cup E_\ell) &= \rho(C_1 \cup \dots \cup C_\ell, H_1 \cup \dots \cup H_\ell) \\ &= \rho(C_1 \cup \dots \cup C_{\ell-1} \cup \{t_\ell\}, H_1 \cup \dots \cup H_{\ell-1}) \end{aligned}$$

since  $t_\ell$  was a critical vertex. However, this is a contradiction, since

$$L \cup \{x_{j+1}, \dots, x_n\} \rightarrow C(C_1 \cup \dots \cup C_{\ell-1} \cup \{t_\ell\}, H_1 \cup \dots \cup H_{\ell-1})$$

but

$$L \cup \{x_{j+1}, \dots, x_n\} \not\rightarrow C(E_1 \cup F_1 \cup \dots \cup E_\ell).$$

This completes the proof for the first case.

CASE 2. There are infinitely many red generators.

In this case we construct a sequence  $R_1 < \sigma(1) < R_2 < \sigma(2) < \dots$  where  $R_1, R_2, \dots$  are red generators and  $\sigma(1), \sigma(2), \dots$  are red vertices forming an isomorphic copy of  $D$ . Here we use sets of the form  $\{\sigma(i_1), \dots, \sigma(i_j)\}$  as  $\Pi_1(E_i), \Pi_2(E_i)$  in the formulas  $\psi(R_j)$  to create the  $F'_i$ 's. The details of this construction are essentially the same as in the previous case and indeed can be obtained by systematically replacing symbols in the proof of Case 1 by appropriate other symbols. Essentially in the same way as formula  $\psi(X)$  can be obtained from  $\varphi(X)$  by formal negation and then replacing the phrase all but finitely many by the phrase infinitely many. The details will therefore be omitted.

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*Department of Mathematics*  
*University of Calgary*  
*Calgary, Alberta*  
*T2N 1N4*

*Department of Mathematics*  
*University of Calgary*  
*Calgary, Alberta*  
*T2N 1N4*