

# RANK-1 PERTURBATIONS AND THE LANCZOS METHOD, INVERSE ITERATION, AND KRYLOV SUBSPACES

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## Abstract

The heart of the Lanczos algorithm is the systematic generation of orthonormal bases of invariant subspaces of a perturbed matrix. The perturbations involved are special since they are always rank-1 and are the smallest possible in certain senses. These minimal perturbation properties are extended here to more general cases.

Rank-1 perturbations are also shown to be closely connected to inverse iteration, and thus provide a novel explanation of the global convergence phenomenon of Rayleigh quotient iteration.

Finally, we show that the restriction to a Krylov subspace of a matrix differs from the restriction of its inverse by a rank-1 matrix.

## 1. Introduction

Rank-1 matrices are an integral part of numerical linear algebra; to give just one example, the reduction of a matrix to tridiagonal form by Householder transformations. More recently, rank-1 matrices have found other uses in eigenvalue problems. For example, the divide-and-conquer strategy uses a rank-1 perturbation to decompose a tridiagonal matrix into a decoupled tridiagonal form (see [1]). In [2] a rank-1 perturbation is also used as a starting point for homotopy methods.

To motivate what is to follow, we begin with some elementary facts concerning rank-1 perturbations. These are in the spirit of backward error analysis and show how small a perturbation is required to specify an eigenpair.

*Fact 1* Any scalar  $\mu$  and unit vector  $z$  is an eigenpair of a perturbation of  $A$ :

$$\begin{aligned} [A - (A - \mu)zz^T]z &= \mu z, \\ \|(A - \mu)zz^T\| &= \|(A - \mu)z\|. \end{aligned}$$

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*Fact 2* The above perturbation is the smallest possible: if  $E$  is any other matrix for which  $(A - E)z = \mu z$ , then

$$\|(A - \mu)z\| = \|Ez\| \leq \|E\|.$$

*Fact 3* The norm of the residual  $\|(A - \mu)z\|$  is minimized by the Rayleigh quotient

$$\rho(z) = \langle Az, z \rangle / \langle z, z \rangle,$$

that is,

$$\|(A - \rho(z))z\| \leq \|(A - \mu)z\| \quad \forall \mu.$$

Thus  $D = (A - \rho(z))zz^T$  is the smallest matrix for which the unit vector  $z$  is an eigenvector of  $A - D$ .

## 2. The Lanczos method

In this section, we will see how the preceding results apply to the Lanczos method.

The Lanczos method for approximating outer eigenvalues and eigenvectors of a symmetric  $n \times n$  matrix  $A$  can be motivated in several different ways; for details see [4]. Algebraically, it is the successive formation of  $m \times m$  tridiagonal matrices  $T_m$  defined by

$$T_m = Q_m^T A Q_m. \quad (1)$$

The  $n \times m$  orthonormal matrices  $Q_m$  are built up column by column,

$$Q_{m+1} = [q_1, \dots, q_{m+1}]$$

from vectors defined by

$$\begin{aligned} A Q_m - Q_m T_m &= \hat{q}_{m+1} e_m^T, \\ q_{m+1} &= \hat{q}_{m+1} / \|\hat{q}_{m+1}\|, \\ e_m &= (0, \dots, 1)^T \in \mathcal{R}^m. \end{aligned} \quad (2)$$

The matrix  $T_m$  is the restriction of  $A$  to the subspace spanned by the columns of  $Q_m$  and as such its eigenvalues  $\rho_i$  are approximations to outer eigenvalues of  $A$ . The corresponding eigenvectors,  $y_i$ , of  $T_m$  need to be transformed as  $x_i = Q_m y_i$  to yield approximations to the corresponding eigenvectors of  $A$ .

From now on we will assume  $m$  is fixed and understand that the  $\rho_i$ ,  $x_i$ ,  $y_i$  depend on  $m$ . We shall also assume the  $x_i$  are normalized:  $\|x_i\| = 1$ .

Of course, the Lanczos *algorithm* does not form the matrices explicitly in (1) and (2), but rather is a clever way of recursively generating the vectors  $q_m$  and using

this information to obtain  $T_{m+1}$  by bordering  $T_m$ . For our purposes, though, (2) highlights an essential but often overlooked feature, namely, the rank-1 right hand side. It is this aspect which we focus on in this section.

From (2) we get immediately

$$(A - \rho_i)x_i = \hat{q}_{m+1}e_m^T y_i.$$

That is, each residual is a multiple of  $q_{m+1}$ .

Parallel residuals is essentially a geometric property. The elegance of the Lanczos algorithm turns on the fact that the orthonormal columns of  $Q_m$  span the Krylov subspace

$$\mathcal{K}_m = \text{span}\{q_1, \dots, A^{m-1}q_1\}.$$

The scalars  $\rho_i$  and vectors  $x_i$  are the eigenpairs of the projection of  $A$  onto  $\mathcal{K}_m$  :

$$Q_m Q_m^T A Q_m Q_m^T x_i = \rho_i x_i, \quad i = 1, \dots, m,$$

which is just the Galerkin approximation on  $\mathcal{K}_m$ . Therefore each residual  $r_i = (A - \rho_i)x_i$  is orthogonal to  $\mathcal{K}_m$  (which is the motivation and definition of the Galerkin approximation). At the same time, it is clear that  $r_i \in \mathcal{K}_{m+1}$ . However,  $\mathcal{K}_{m+1}$  is at most one dimension more than  $\mathcal{K}_m$ , and therefore all the residuals are contained in a 1-dimensional subspace. So we conclude as before that the residuals are scalar multiples of a unit vector  $r : r_i = \varepsilon_i r$ . And of course  $r = q_{m+1}$ .

Since the  $x_i$  are mutually orthogonal, we have, similar to Fact 1,

$$\left[ A - \sum r_j x_j^T \right] x_i = \rho_i x_i.$$

On setting  $s = \sum \varepsilon_j x_j^T$ , this is  $[A - r s^T] x_i = \rho_i x_i$  or

$$[A - r s^T] X = X \Delta, \tag{3}$$

$$X = [x_1, \dots, x_m], \quad \Delta = \text{diag}(\rho_1, \dots, \rho_m).$$

In this form, the Lanczos method is interpreted as producing an invariant subspace of a perturbation of  $A$ .

In [1], it is shown how this can be carried one step further by constructing a symmetric rank-1 matrix  $E$  for which the subspace spanned by the columns of  $X$  is an invariant subspace of  $A + E$ . This can then be used as another way of approximating the eigenvalues of  $A$ .

That the perturbation which occurs in the Lanczos method is necessarily rank-1 rests on the fact that all the residuals are parallel. This last property is shared by all rank-1 changes, that is, if  $(A - ab^T)z = \mu z$ , then  $(A - \mu)z$  is a multiple of  $a$ .

It also turns out that the perturbation  $r s^T$  is the smallest possible in the same sense as Fact 2.

**THEOREM 1.** *If  $E$  is such that  $(A - E)x_i = \rho_i x_i$  for each  $i$ , then*

$$\|rs^T\| \leq \|E\|.$$

**PROOF.** Since the  $x_i$  are orthonormal, and  $\|r\| = 1$ , we have

$$\|rs^T\|^2 = \|s\|^2 = \sum \varepsilon_j^2 = \left\| \sum \varepsilon_j^2 r \right\|.$$

But

$$\begin{aligned} \left\| \sum \varepsilon_j^2 r \right\| &= \left\| \sum \varepsilon_j r_j \right\| \\ &= \left\| \sum \varepsilon_j E x_j \right\| \\ &= \left\| E \sum \varepsilon_j x_j \right\| \\ &= \|Es\| \\ &\leq \|E\| \|s\|. \end{aligned}$$

That is,  $\|rs^T\| \leq \|E\|$ .

**ALTERNATIVE PROOF.** First write  $(A - E)X = X\Delta$ , which, with (3), gives  $EX = rs^T X$ . This, together with the fact that  $X$  is orthonormal and  $\|r\| = 1$ , yields,

$$\|s\| = \|rs^T X\| = \|EX\| \leq \|E\| \|X\| = \|E\|.$$

More generally, for any subspace spanned by the columns of an orthonormal  $Z$ , there is a perturbation of  $A$  for which  $\text{span}(Z)$  is an invariant subspace.

**THEOREM 2.** *Suppose  $Z = [z_1, \dots, z_m]$  is orthonormal,  $M$  is  $m \times m$ , and  $R = [AZ - ZM]Z^T$ , then*

$$\begin{aligned} (A - R)Z &= ZM, \\ \|R\| &= \|AZ - ZM\| \end{aligned}$$

and  $\|R\| \leq \|E\|$  for every  $E$  such that  $(A - E)Z = ZM$ .

**PROOF.** Since  $Z$  is orthonormal, it follows that

$$\|R\| = \|(AZ - ZM)Z^T\| = \|AZ - ZM\| = \|RZ\|$$

and

$$\|RZ\| = \|EZ\| \leq \|E\| \|Z\| = \|E\|.$$

Just as the norm of the residual is minimized by the Rayleigh quotient (Fact 3), so it is that  $\|AZ - ZM\|$ , and consequently  $\|R\|$ , is minimized by  $M = Z^T AZ$  (see [4]).

As a final comment, we note that as the Lanczos algorithm proceeds, it is not necessarily the case that the norms of the perturbations  $rs^T$  decrease monotonically.

### 3. Inverse iteration and iterated Galerkin approximations

**3.1. Inverse iteration** As noted before,  $[A - (A - \mu)zz^T]z = \mu z$ , for an arbitrary  $A$ , scalar  $\mu$  and unit vector  $z$ . Now consider the corresponding left eigenvector  $y$  given by

$$y^T[A - (A - \mu)zz^T] = y^T\mu,$$

that is,

$$(A - \mu)^T y = zz^T(A - \mu)^T y.$$

When  $A$  is symmetric, the left eigenvector  $y$  is the inverse iterate of  $z$ , and, in general, if  $\mu$  is near an eigenvalue,  $y$  will be closer than  $z$  to the corresponding eigenvector.

The idea of smallest perturbation occurs in this context too, and is closely connected to the global convergence property of Rayleigh quotient iteration.

**THEOREM 3.** *If  $D$ ,  $\mu$ , and unit vectors  $x$  and  $y$  satisfy*

$$\begin{aligned} (A - D)x &= \mu x, \\ y^T(A - D) &= \mu y^T \text{ and} \\ \|D\| &\leq \|E\| \quad \forall E : (A - E)x = \mu x, \end{aligned}$$

then

$$\|(A - \mu)^T y\| \leq \|(A - \mu)x\|.$$

**PROOF.**

$$\|(A - \mu)^T y\| = \|D^T y\| \leq \|D^T\| \|y\| = \|D\| \|y\| = \|(A - \mu)x\|.$$

The final equality is just Fact 2.

That is, the residual of the left eigenvector  $y$  (with respect to  $A^T$ ) is less than the residual of the right eigenvector (with respect to  $A$ ). Basically this means that  $y$  is a better approximation to a left eigenvector than  $x$  is to a right eigenvector. The implication of this is that iterating with the adjoint is the best way to improve an approximation to a right eigenvector. It also strongly suggests that inverse iteration should really be performed with  $(A - \mu)^{-1}(A - \mu)^T$  when  $A$  is nonsymmetric.

Combining Theorem 3 with Fact 3 results in a novel interpretation of some of the properties of Rayleigh quotient iteration as defined by

$$(A - \rho(x_k))x_{k+1} = \alpha_k x_k, \quad \|x_{k+1}\| = 1.$$

Using the equivalence between the inverse iterate and left eigenvector this is recast as

$$x_{k+1}^T(A - R_k) = \rho(x_k)x_{k+1}^T, \quad R_k = (A - \rho(x_k))x_k x_k^T.$$

It is evident from Theorem 3 that the residuals  $\|(A - \rho(x_k))x_k\|$  decrease monotonically when  $A$  is normal. In fact, this results in the global convergence property of Rayleigh quotient iteration (see [3, 4]).

Since  $\|R_k\| = \|(A - \rho(x_k))x_k\|$ , the vectors  $x_k$  are eigenvectors of successively smaller perturbations of  $A$ , while the scalars  $\rho(x_k)$  are such that the perturbation at each step is the smallest possible given  $x_k$ .

Bearing in mind the remark about iterating with  $(A - \mu)^{-1}(A - \mu)^{-T}$ , we see that the appropriate strategy in Rayleigh quotient iteration is

$$(A - \rho(x_k))^T(A - \rho(x_k))x_{k+1} = \alpha_k x_k,$$

which is just the alternative version of the iteration examined in [3].

Recall that Lanczos, at step  $m$ , produces vectors  $x_i$  with  $(A - rs^T)x_i = \rho(x_i)x_i$ . In view of the above, it is natural to ask if the left eigenvectors  $y_i$  of  $A - rs^T$  corresponding to  $\rho_i$  have smaller residuals than the right (Lanczos) eigenvectors  $x_i$ . In fact this is not necessarily the case; however, we may expect the residuals of  $y_i$  to be less than those of the  $x_i$  in a collective sense. Such a result is contained in the following more general situation.

Suppose  $Z$  is orthonormal,  $R = [AZ - ZM]Z^T$  for some  $M$ , and  $Y^T(A - R) = MY^T$ , then

$$\|Y^T A - MY^T\| = \|Y^T R\| \leq \|R\| \|Y^T\|,$$

that is,

$$\|A^T Y - Y M^T\| \|Y\|^{-1} \leq \|R\|.$$

In the Lanczos case,  $Z = X$ ,  $M = \Delta$ ,  $R = rs^T$  and  $\|R\| = \|s\|$ , but the columns of  $Y$  are not necessarily mutually orthogonal, even though they can be normalized. We therefore cannot conclude that

$$\sum \| (A - \rho(y_i)I)y_i \|^2 \leq \|s\|^2$$

and counterexamples to this inequality are easily found. However, we do have

$$\begin{aligned} \|(A - \rho(y_i))y_i\| &\leq \|(A - \rho(x_i))y_i\| \\ &= \|sr^T y_i\| \\ &= \|s\| |r^T y_i|, \end{aligned}$$

so

$$\sum \| (A - \rho(y_i))y_i \|^2 \leq \|s\|^2 \sum |r^T y_i|^2.$$

Again, it is not necessarily true that  $\sum |r^T y_i|^2 \leq 1$ . But  $\sum |r^T y_i|^2 = \|Vr\|^2 \leq \|V\|^2$ , where  $V = \sum x_i y_i^T$ , so  $\sum \| (A - \rho(y_i))y_i \|^2 \leq \|V\|^2 \|s\|^2$ .

So far it still an open question as to how Theorem 3 can be extended to cover the Lanczos method.

**3.2. Iterated Galerkin approximation** As an aside, we comment on the way left eigenvectors also occur as refinements in Sloan’s iteration of Galerkin approximations examined in [5].

Sloan’s iteration arises in connection with a sequence of Galerkin approximations

$$P_i A x_i = \mu_i x_i, \quad P_i \rightarrow I, \quad (\mu_i, x_i) \rightarrow (\lambda, u).$$

The  $P_i$  are projections onto subspaces in a nested sequence. The idea behind Sloan’s iteration is that the sequence  $\{Ax_i/\mu_i\}$  converges more rapidly than does  $\{x_i\}$ . What interests us here is the fact that  $AP_i A x_i = \mu_i A x_i$ . When  $A$  is symmetric this is just

$$(P_i A)^T A x_i = \mu_i A x_i.$$

In other words, the Sloan iterates are simply the left eigenvectors of the Galerkin approximations  $P_i A$ .

#### 4. Krylov subspaces

For various reasons it is of some interest to examine the Galerkin-Krylov method for the matrix  $A^{-1}$

$$P A^{-1} y_i = \mu_i y_i,$$

where  $P$  is the orthogonal projection onto the Krylov subspace

$$\mathcal{K}_m = \text{span}\{q, Aq, \dots, A^{m-1}q\}.$$

This leads to a surprising result concerning the ‘Gram’ matrices

$$G(A) = [\langle Ax_i, x_j \rangle] = X^T A X \quad \text{and} \quad G(A^{-1}) = [\langle A^{-1}x_i, x_j \rangle] = X^T A^{-1} X,$$

where  $\{x_i\}$  is an orthonormal basis for  $\mathcal{K}_m$  and  $X = [x_1, \dots, x_m]$ .

**THEOREM 4.**  $G(A^{-1}) - G(A)^{-1}$  is rank-1.

**PROOF.** We use the same notation as in Section 1. Without losing any generality, we may take the  $x_i$  to be the unit eigenvectors of  $PA$ .

Now  $G(A) = \text{diag}(\rho_1, \dots, \rho_m) = \Delta$ , by properties of the vectors  $x_i$ , and

$$G(A^{-1}) - G(A)^{-1} = X^T A^{-1} X - \Delta^{-1}.$$

But, recalling the results of Section 1,

$$[A - r s^T] X = X \Delta,$$

that is,

$$\begin{aligned}\Delta^{-1} &= X^T[A - rs^T]^{-1}X \\ &= X^T\left[A^{-1} + \frac{A^{-1}rs^TA^{-1}}{1 - s^TA^{-1}r}\right]X.\end{aligned}$$

Therefore

$$G(A^{-1}) - G(A)^{-1} = \frac{-X^TA^{-1}rs^TA^{-1}X}{1 - s^TA^{-1}r},$$

which has rank 1 as claimed.

This relies crucially on the  $x_i$  being orthonormal vectors which span a Krylov subspace for some  $q$ . Neither condition can be relaxed.

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